

Limits and colimits – Part II

Category theory and its applications – ITI9200

<https://compose.ioc.ee>

15 March 2021

Categories having (co)limits

Category with \mathcal{S} -limits

Let \mathcal{C} be a category. We say that \mathcal{C} **has \mathcal{S} -limits** if, for all categories $\mathcal{I} \in \mathcal{S}$ and functors $F: \mathcal{I} \rightarrow \mathcal{C}$, there is a limit cone over F in \mathcal{C} .

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Theorem

Let \mathcal{C} be a category. Suppose that \mathcal{C}

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- There is a dual version:
small coproducts + coequalisers = small colimits

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- Such a cone is a family of morphisms $b_i: C \rightarrow F(i)$
satisfying the equations

$$F(h) \circ b_i = b_j, \quad i, j \in \text{Ob}(\mathcal{I}), \quad h \in \text{Hom}_{\mathcal{I}}(i, j).$$

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- If we have small products, and \mathcal{I} is small, a family of morphisms with codomain $F(i)$ corresponds to a generalised element of $\prod_{i \in \text{Ob}(\mathcal{I})} F(i)$.

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- If we also have equalisers, we can restrict $\prod_{i \in \text{Ob}(\mathcal{I})} F(i)$ to generalised elements that satisfy equation (1).

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Proof

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Since \mathcal{C} has small products, we can form the products

$$\prod_{i \in \text{Ob}(\mathcal{I})} F(i), \quad \prod_{\substack{i, j \in \text{Ob}(\mathcal{I}), \\ h \in \text{Hom}_{\mathcal{C}}(i, j)}} F(j)$$

the first indexed by all objects, the second by all *morphisms* of \mathcal{I} ,
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For each morphism $h: i \rightarrow j$, we have a pair of morphisms

$$F(h) \circ \pi_i, \quad \pi_j: \prod_{i \in \text{Ob}(\mathcal{I})} F(i) \rightarrow F(j);$$

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this corresponds to the equation $F(h) \circ b_i = b_j$.

Limits = products + equalisers

Proof, cont.d

The families

$$\left(F(h) \circ \pi_i \right)_{\substack{i,j \in \text{Ob}(\mathcal{I}), \\ h \in \text{Hom}_{\mathcal{I}}(i,j)}}, \quad \left(\pi_j \right)_{\substack{i,j \in \text{Ob}(\mathcal{I}), \\ h \in \text{Hom}_{\mathcal{I}}(i,j)}}$$

produce a parallel pair of morphisms

$$\langle F(h) \circ \pi_i \rangle, \quad \langle \pi_j \rangle: \prod_{i \in \text{Ob}(\mathcal{I})} F(i) \rightarrow \prod_{\substack{i,j \in \text{Ob}(\mathcal{I}), \\ h \in \text{Hom}_{\mathcal{I}}(i,j)}} F(j);$$

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Let $e: \lim F \rightarrow \prod_{i \in \text{Ob}(\mathcal{I})} F(i)$ be the equaliser of this parallel pair.

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Proof, cont.d

We claim that $\lim F$ together with $\pi \circ e := (\pi_i \circ e)_{i \in \text{Ob}(\mathcal{I})}$ is a limit cone over F .

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$$F(h) \circ \pi_i \circ e = \pi'_h \circ \langle F(h) \circ \pi_i \rangle \circ e = \pi'_h \circ \langle \pi_j \rangle \circ e = \pi_j \circ e$$

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- 2 Given another cone b over F , the family $(b_i)_{i \in \text{Ob}(\mathcal{I})}$ induces

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We need to show that

$$\langle F(h) \circ \pi_i \rangle \circ \langle b_i \rangle = \langle \pi_j \rangle \circ \langle b_i \rangle$$

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Because b is a cone over F , for all morphisms $h: i \rightarrow j$, we have

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which implies $\langle F(h) \circ \pi_i \rangle \circ \langle b_i \rangle = \langle \pi_j \rangle \circ \langle b_i \rangle$.

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Because e is the equaliser of $\langle F(h) \circ \pi_i \rangle$ and $\langle \pi_j \rangle$, we obtain a unique morphism

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$$(\pi_i \circ e) \circ k = b_i, \quad i \in \text{Ob}(\mathcal{I}).$$

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We have constructed $k: \text{tip}(b) \rightarrow \lim F$ such that

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This is a solution to the system

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and we can prove that it is unique from the uniqueness properties of products and equalisers.

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This proves that $\pi \circ e$ is a limit cone over F . □

Complete and cocomplete categories

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- When \mathcal{C} is small, the functor category $[\mathcal{C}, \mathbf{Set}]$ has all small limits and colimits.
- A poset P has all limits (resp. colimits) if and only if it has greatest lower bounds (resp. least upper bounds) of arbitrary collections of elements.

Functors and limit cones

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Then Gb with components

$$Gb_i: G(\text{tip}(b)) \rightarrow GF(i), \quad i \in \text{Ob}(\mathcal{I}),$$

is a cone over GF .

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The functor G **preserves limits of** F if, whenever f is a limit cone over F , Gf is a limit cone over GF .

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Preservation of (co)limits is also very important in the theory of **adjunctions** (next lecture!)

Doubly indexed limits

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Then $i \mapsto \lim_{\mathcal{J}} F(i, -)$ extends to a functor

$$\lim_{\mathcal{J}} F: \mathcal{I} \rightarrow \mathcal{C}.$$

Doubly indexed limits

Sketch of proof

Let $h: i \rightarrow i'$ be a morphism in \mathcal{I} ,
let f be the chosen limit cone over $F(i, -)$,
and let f' be the chosen limit cone over $F(i', -)$.

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For each $j \in \text{Ob}(\mathcal{J})$, we have morphisms

$$\begin{aligned} f_j: \lim_{\mathcal{J}} F(i, -) &\rightarrow F(i, j), \\ F(h, j): F(i, j) &\rightarrow F(i', j), \end{aligned}$$

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whose composites $(F(h, j) \circ f_j)_{j \in \text{Ob}(\mathcal{J})}$ form
a cone over $F(i', -)$ with tip $\lim_{\mathcal{J}} F(i, -)$.

Doubly indexed limits

Sketch of proof, cont.d

By the universal property of the cone f' , whose tip is $\lim_{\mathcal{J}} F(i', -)$, we obtain a unique morphism

$$\lim_{\mathcal{J}} F(i, -) \rightarrow \lim_{\mathcal{J}} F(i', -),$$

and we define $\lim_{\mathcal{J}} F(h, -)$ to be this morphism.

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We can show that this assignment defines a functor (i.e. is compatible with identities and composition). □

Commutativity of (co)limits

If we are in the conditions of the previous proposition, we can ask if $\lim_{\mathcal{J}} F: \mathcal{I} \rightarrow \mathcal{C}$ has a limit.

Theorem (commutativity of limits)

Suppose that

$$\lim_{\mathcal{J}} F: \mathcal{I} \rightarrow \mathcal{C}, \quad \lim_{\mathcal{I}} F: \mathcal{J} \rightarrow \mathcal{C}$$

are defined and have limit cones with tips

$$\lim_{\mathcal{I}} \lim_{\mathcal{J}} F, \quad \lim_{\mathcal{J}} \lim_{\mathcal{I}} F.$$

Then $\lim_{\mathcal{I}} \lim_{\mathcal{J}} F$ and $\lim_{\mathcal{J}} \lim_{\mathcal{I}} F$ are canonically isomorphic, and are tips of limit cones over F .

Commutativity of (co)limits

Dually, we have isomorphisms

$$\operatorname{colim}_{\mathcal{I}} \operatorname{colim}_{\mathcal{J}} F \simeq \operatorname{colim}_{\mathcal{J}} \operatorname{colim}_{\mathcal{I}} F$$

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However **limits** seldom commute with **colimits**...

Non-commutativity of limits with colimits

Let $F: \mathcal{I} \times \mathcal{J} \rightarrow \mathcal{C}$ be a functor, and suppose

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...but it is not, in general, an isomorphism.

Non-commutativity of limits with colimits

A simple counterexample:

Let $F: 2 \times 2 \rightarrow \mathbf{Set}$ be defined by

$$F(i, j) := 1 \equiv 1_{i,j}$$

for all $i, j \in 2$.

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$$\operatorname{colim}_{\mathcal{I}} \lim_{\mathcal{J}} F = (1_{1,1} \times 1_{1,2}) + (1_{2,1} \times 1_{2,2}) \simeq 2,$$

$$\lim_{\mathcal{J}} \operatorname{colim}_{\mathcal{I}} F = (1_{1,1} + 1_{2,1}) \times (1_{1,2} + 1_{2,2}) \simeq 4.$$

Filtered colimits

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A category \mathcal{C} is *filtered* if, for all finite categories \mathcal{I} and functors $F: \mathcal{I} \rightarrow \mathcal{C}$, there exists a cone under F .

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- A category with finite colimits is filtered.
- A poset is filtered when every finite collection of elements has an upper bound.

Filtered colimits

A **filtered colimit** is a colimit of a functor whose domain is a filtered category.

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Let \mathcal{I} be a filtered category, \mathcal{J} a finite category, and $F: \mathcal{I} \times \mathcal{J} \rightarrow \mathcal{C}$ a functor.

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A colimit of such a functor is called a **sequential colimit**. For example, an **increasing** sequence of sets

$$X_0 \subseteq X_1 \subseteq \dots \subseteq X_n \subseteq \dots$$

defines a functor $\mathbb{N} \rightarrow \mathbf{Set}$, whose colimit is

$$\bigcup_{n \in \mathbb{N}} X_n$$

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It follows that **sequential colimits commute with finite limits**. We recover, for example, the fact that

$$\left(\bigcup_{n \in \mathbb{N}} X_n \right) \times Y \simeq \bigcup_{n \in \mathbb{N}} (X_n \times Y)$$

in **Set**.

Representable property of limits

When $F: \mathcal{I} \rightarrow \mathcal{C}$ is a functor from a small category, we have

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Theorem

The following are equivalent:

- 1 F has a limit;
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Interpretation: **generalised elements** of the limits, i.e.

$x \in \mathrm{Hom}_{\mathcal{C}}(C, \lim F)$, are the same as

elements $x \in \lim \mathrm{Hom}_{\mathcal{C}}(C, F-)$ of the limit of sets of elements

Representable property of limits

A little teaser of future lectures:
the **Yoneda lemma** will allow us to identify

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So $[\mathcal{C}^{\text{op}}, \mathbf{Set}]$ is a kind of “complete extension” of \mathcal{C} .