

A short recap on products of monoids

Recall the def of monoid homomorphism
btwn two monoids M, N :

a function $f: M \rightarrow N$ such that

$$f(m \cdot m') = f(m) \cdot f(m')$$

\uparrow operation in M \uparrow op in N

Definition ISOMORPHISM of monoids

$$\begin{array}{ccc} M & \xrightarrow{f} & N \\ & \xrightarrow{g} & M \\ fg = id; gf = id \end{array}$$

+ we have seen a few examples.

+ Defined the product of monoids as

$$M \times N = \{(m, n) \mid m \in M, n \in N\}$$

with operation pointwise

identity $(1_M, 1_N)$.

⚠ A central idea of category theory.
is that one can recognize $M \times N$
thru a UNIVERSAL PROPERTY.

A monoid P is $\sim M \times N$ IF AND ONLY IF

P satisfies a certain property that
will characterize it "uniquely"

(Similarly, "being a top" characterizes "unique")
bottom

Start observing that for $M \times N$ the following property is true *

Given a monoid A and monoid homomorphisms $u: A \rightarrow M$
 $v: A \rightarrow N$

there exists a unique monoid homomorphism $\langle u, v \rangle: A \rightarrow M \times N$

with the property that

$$u = (A \xrightarrow{\langle u, v \rangle} M \times N \xrightarrow{\text{proj}_M} M)$$

$$v = (A \xrightarrow{\langle u, v \rangle} M \times N \xrightarrow{\text{proj}_N} N) \gg$$

- Define $\langle u, v \rangle$
- Show that if there is another $\langle u, v' \rangle'$ with the same property then $\langle u, v \rangle = \langle u, v' \rangle'$

CLAIM If P satisfies the same property
then there is a unique [bijective
homomorphism] $P \xrightarrow{\sim} M \times N$.

So, P is "essentially equal" to $M \times N$

Start by studying what happens when $A = M \times N$ in the previous property *

Diagrammatic notation:

$$\begin{array}{ccc}
 & A & \\
 u \swarrow & ! \downarrow & \searrow v \\
 M & \xleftarrow{\text{proj}_M} & M \times N \xrightarrow{\text{proj}_N} N
 \end{array}$$

& if P satisfies
the same
property $P \cong M \times N$

Proof: Since P satisfies the U.P.

$$\exists ! M \times N \xrightarrow{k} P \text{ such that}$$

$$\begin{array}{ccc}
 & M \times N & \\
 \text{proj}_M \swarrow & ! \downarrow k & \searrow \text{proj}_N \\
 M & \xleftarrow{P_M} & P \xrightarrow{P_N} N
 \end{array}$$

Since $M \times N$ satisfies the U.P.

$$\exists ! P \longrightarrow M \times N$$

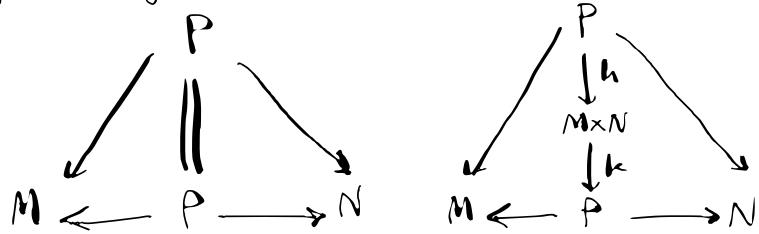
$$\begin{array}{ccc}
 & P & \\
 P_M \swarrow & ! \downarrow h & \searrow P_N \\
 M & \xleftarrow{\text{proj}_M} & M \times N \xrightarrow{\text{proj}_N} N
 \end{array}$$

Uniqueness implies that $P \rightarrow M \times N \rightarrow P = \text{id}$

$$M \times N \rightarrow P \rightarrow M \times N = \text{id}$$

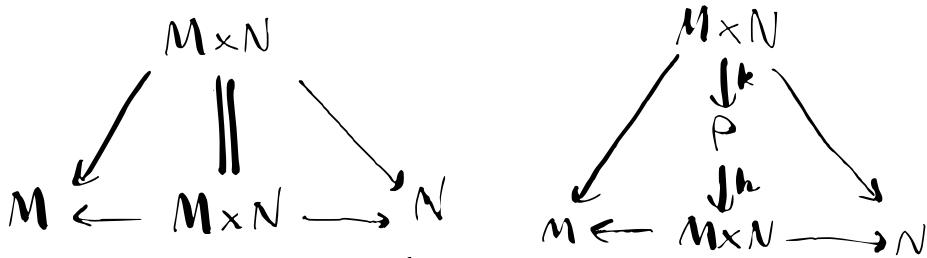
(because if $A = P$, id does the job
if $A = M \times N$, id does the job)-

More precisely: if $A = P$



both solve the problem

If $A = M \times N$



both solve the problem.

But then, $kh = id$ & $hk = id$

Morale of this story

Category theory is a language apt to characterize structures and operations on structures (like Cartesian products) "synthetically", which means without looking inside the carriers of the structures, but relying only on universal properties of arrow diagrams.

FREE MONOIDS

Definition of the monoid of lists

Let A be a set, consider the set

$$\text{List}(A) = \{(a_1 - a_n) \mid a_i \in A, n \geq 0\}$$

with the convention that if $n=0$

$()$ is the empty list, with no elements

Constructors for $\text{List}(A)$:

List A :

$[] : \text{List } A$

$\text{cons}(a, as) = a :: as$

$\text{cons} : A \rightarrow \text{List } A \rightarrow \text{List } A$

-- $\text{cons}(a, as)$ generates the list

-- having a as head, as as tail

Monoid operation on $\text{List}(A)$: concat

$(\text{++}) : \text{List } A \rightarrow \text{List } A \rightarrow \text{List } A$

$[] ++ ys = ys$

$(x :: xs) ++ ys = x :: (xs ++ ys)$

This is a monoid operation

• — associative

• — $[]$ is identity on both sides

Proof by induction

The monoid $\text{List}(A)$ is clearly non-commutative

The monoid $\text{List}A$ has no nontrivial invertible elements

(In a monoid M , $x \in M$ is INVERTIBLE if $\exists y$ such that $xy = 1$, $yx = 1$)

$$as + bs = [] \Rightarrow as = bs = []$$

Proposition Let $A = \{s\}$ a set with a single symbol.

Then $\text{List}(A)$ is the monoid of natural numbers

Proof Define N as the set

$$\{0, s0, ss0, sss0, ssss0, \dots\}$$

List A , (to h), is the set

$$\{[], (s), (ss), (sss), (ssss), \dots\}$$

And these two sets are evidently "the same" up to relabeling their elements -