

CATEGORY THEORY ITI9200
EXERCISES



V. Brauner, *L'objet qui rêve II*, 1938.

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CHAPTER 0

A word of introduction

As a general rule, solving an exercise in pure Mathematics, the ability to think should be rewarded more than correctness; this means that good ideas leading to wrong answers are more valuable than bad ideas yielding the correct answer.

Category theory follows an even stronger claim: it is based on the belief that the right answer is useless when found by means of an unenlightening train of thought; you should grow accustomed to this philosophy.

Don't expect all questions to be straightforward. On the contrary, some exercises are meant to be difficult and right above your level; other exercises are meant to force you learn new things in the process.



FIGURE 1. You are supposed to get your hands dirty.

Mathematics, structurally

1. Orders and relations

Preliminaries. Given sets A, B , we will say that two functions $f, g : A \rightarrow B$ **coincide** if they assume the same values elementwise; in symbols,

$$f \equiv g \iff \forall a \in A, f(a) = g(a). \quad (1.1)$$

This is usually called the **extensionality principle** for functions.

Let X be a set. One of the axioms of set theory asserts that the following collection

$$P(X) := \{U \mid U \subseteq X\} \quad (1.2)$$

is a set. It is the set whose elements are exactly all **subsets** of X .¹

We can build a correspondence between $P(X)$ and another set: the set whose elements are all functions $f : X \rightarrow \{0, 1\}$. In order to define a function

$$c_\bullet : P(X) \longrightarrow 2^X \quad (1.3)$$

just send $U \subseteq X$ to the function $c_U : X \rightarrow \{0, 1\}$ sending x to 1 if and only if $x \in U$ (so, since there is no other choice, $c_U(x) = 0$ if and only if $x \notin U$). The function c_U is called the **characteristic function** of the subset $U \subseteq X$.

¹Recall that a subset U of X is a set with the property that all elements of U are also elements of X : in symbols, $U \subseteq X$ means the formula

$$x \in U \Rightarrow x \in X.$$

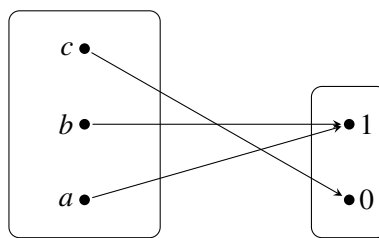


FIGURE 1. A function $\{a, b, c\} \rightarrow \{0, 1\}$: the one sending $a \mapsto 1$, $b \mapsto 1$, $c \mapsto 0$.

PROPOSITION 1.1. The correspondence c_\bullet is a function, and it is a bijection, because

- it is injective: when the functions c_U, c_V are the same functions (which means: for every $x \in A$, $c_U(x) = c_V(x)$), then $x \in U$ if and only if $x \in V$, and thus the set U is equal to V ;
- it is surjective, because given any function $f : X \rightarrow 2$, f is the characteristic function of the set

$$f^{\leftarrow}(1) := \{x \in X \mid f(x) = 1\}. \quad (1.4)$$

EXERCISE OR.1 \square Fill in the details in the proof of [Proposition 1.1](#) above. Draw a picture of all functions $\{a, b, c\} \rightarrow \{0, 1\}$ and verify that they are $8 = 2^3$, as expected.

EXERCISE OR.2 \square Do the same exercise, but backwards: induce from a series of specific examples. If the function $f : \{a, b, c\} \rightarrow \{0, 1\}$ of [Figure 1](#) is represented as the triple (110), to denote the fact that the image of a under f is 1, the image of b under f is 1, and the image of c is 0, explain (in words) what's the meaning of the following sentence:

Here's a list of all functions $\{a, b, c\} \rightarrow \{0, 1\}$:

(000) (100) (010) (001) (110) (101) (011) (111).

What can you infer about a similar statement regarding the set of functions $\{a, b, c, d\} \rightarrow \{0, 1\}$? What can you infer about a similar statement regarding the set of functions $\{a, b, c, d, e\} \rightarrow \{0, 1\}$? What can you infer about a similar statement regarding the set of functions $\{a_1, \dots, a_n\} \rightarrow \{0, 1\}$?

The weak point of the above 'inductive' approach to discover that the set of all subsets of A is as big as the set of all functions $A \rightarrow \{0, 1\}$ is that it cannot be generalised to the case when A is infinite: if A is finite, say with n elements, then PA has 2^n elements; but if A is infinite, what does 2^{infinite} mean –provided it even means anything? **Cardinal arithmetic** is the part of set theory that makes sense of this statement, and others.

Informally speaking, one of the major achievements of Cantor's set theory is to acknowledge the existence of *more than one* infinite set; and in fact, of something more: given *any* infinite, there is a way to build one infinite that is strictly bigger. (Compare this statement with the following fact of life: given a number n , there is a way to build a strictly larger number.)

DEFINITION 1.2. Given a set A , a function $p : A \rightarrow \{0, 1\}$ is called a **predicate** or a **proposition**; in mathematical discourse, a predicate is a statement you make about an object A ; a proposition is a statement that you make regarding an object, and that you are interested in deeming true or false.

The difference between the two concepts is tenuous, and in fact, you model them mathematically using the same concept: a function $A \rightarrow \{0, 1\}$.

Following [Proposition 1.1](#), the predicate/proposition $p : A \rightarrow \{0, 1\}$ defines a unique subset of A : the subset of elements of A that make the proposition true.

Following standard practice (it is for example very common in programming) we blur the distinction between the lines of the following table:

| Booleans | answers | truth values |
|----------|---------|--------------|
| 0 | no | false |
| 1 | yes | true |

TABLE 1. Truth values, booleans, binary answers to questions.

DEFINITION 1.3. A **bijection** between sets X, Y is a function $f : X \rightarrow Y$ that is injective and surjective:

- (injective): if for some $x, x' \in X$ we have $f(x) = f(x')$, then $x = x'$;
- (surjective): for every $y \in Y$, there exists at least one $x \in X$ such that $f(x) = y$.

EXERCISE OR.3 \square Prove that a function $f : X \rightarrow Y$ is bijective if and only if there exists a function $g : Y \rightarrow X$ such that for every $y \in Y$, $f(g(y)) = y$ and for every $x \in X$, $g(f(x)) = x$. Such a g is called **inverse** of f ; so, f is a bijection if and only if it has an inverse (we also say that f ‘is invertible’).

NOTATION 1.4. To denote that there exists *some* unnamed bijection between two sets A, B we often write $A \cong B$. Other synonyms for ‘there exists a bijective function $f : A \rightarrow B$ ’ are:

- ‘the set A can be identified with B ’;
- ‘the set A , or equivalently the set B ’, and also
- ‘the set A , also called the set B ’.

(This remark is a half-joke to convey the idea that if $A \cong B$ then the two sets ‘behave the same way’: each property of A is also enjoyed by B because it can be ‘transported’ along a bijection $f : A \rightarrow B$, and every property of B can be transported along its inverse $f^{-1} : B \rightarrow A$.)

EXERCISE OR.4 \square Show that if f has an inverse g as above, g is unique. So, we are allowed to call g *the* inverse of f , when it exists.

EXERCISE OR.5 \square If f is not injective, it cannot be invertible; what is the problem, exactly? What is the obstruction to define the inverse of f ? If f is not surjective, it cannot be invertible; again, where is the problem exactly?

DEFINITION 1.5. A set A is called **infinite** (à la Dedekind) if there exist a *proper* subset $U \subset A$ and a bijection $f : U \rightarrow A$.²

EXERCISE OR.6 □ Welcome to the magic, counterintuitive world of infinite sets! Prove that $\mathbf{N} = \{0, 1, 2, \dots\}$ is infinite à la Dedekind (this was first observed by Galileo). Prove that \mathbf{Z} is infinite à la Dedekind. Write down an explicit bijection between \mathbf{N} and \mathbf{Z} ; prove that there exists a bijection between \mathbf{N} and $\mathbf{N} \times \mathbf{N}$. Prove that there exists a bijection between the set \mathbf{N} and the set $\{\text{DIS}, \text{DAT}\} \times \mathbf{N}$, defined as

$$\{\text{DIS}, \text{DAT}\} \times \mathbf{N} := \{(\text{DIS}, 0), (\text{DIS}, 1), (\text{DIS}, 2), \dots, (\text{DIS}, n), \dots, \\ (\text{DAT}, 0), (\text{DAT}, 1), (\text{DAT}, 2), \dots, (\text{DAT}, n), \dots\} \quad (1.5)$$

EXERCISE OR.7 □ (Very hard, but try to chew this diamond) Let A be an infinite set; prove that there exist a subset E of A and a bijection $f : \mathbf{N} \rightarrow E$. So, \mathbf{N} is the **smallest** infinite set, because every other infinite set contains a copy of it.

DEFINITION 1.6. Given sets X, Y , a **relation between X and Y** is a subset of the product $X \times Y$; a **relation on X** is a relation between X and itself. The set $\text{Rel}(X, Y)$ of all relations between X and Y then is just the set $P(X \times Y) \cong 2^{X \times Y}$ of all subsets of $X \times Y$.

EXERCISE OR.8 □ Show that there exists a bijection σ between $P(X \times Y)$ and $P(Y \times X)$, induced by a bijection $X \times Y \cong Y \times X$.

Given a relation R between X and Y , the **opposite relation** R^{op} is the image of $R \in P(X \times Y)$ under the bijection σ . The relation R and its opposite carry the exact same amount of information, although formally $R : X \rightarrow Y$ and $R^{\text{op}} : Y \rightarrow X$. This remark is meant to formalise the idea that the existence of a relation R between sets X, Y is not a ‘directed’ property (compare this with the fact that a function has instead a specified *domain* and *codomain*).

EXERCISE OR.9 □ Count how many relations there are on a finite set $X = \{x_1, \dots, x_n\}$.

From the very definition of the set $\text{Rel}(X, Y)$ it follows that there exists a natural notion of partial order between relations, defined by $R \leq S$ if and only if $R \subseteq S$; thus the intersection (resp., union) of an arbitrary number of relations between X and Y is still a relation.

EXERCISE OR.10 □ Does the poset $(\text{Rel}(X, Y), \leq)$ has a top element, a bottom element? Define the **complementary** relation of a given $R \in P(X \times Y)$.

EXERCISE OR.11 □ On the relation between relations and functions. Show that a function $f : X \rightarrow Y$ is precisely a relation $R \subseteq X \times Y$ with the property that each ‘ x -section’ set $R_x := \{y \in Y \mid (x, y) \in R\}$ is a singleton $\{y\} =: \{y(x)\}$.

²Note that this definition says when a set is infinite: as counterintuitive as it may seem (because finite sets are ‘evidently there’ in everyday life, whereas no one can see an infinite one), in Mathematics we explicitly define infinite sets, and we just say that a set is finite if it is not infinite.

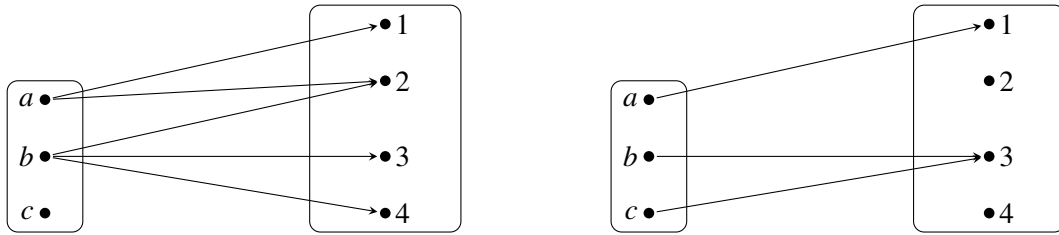


FIGURE 2. On the left, a relation $A = \{a, b, c\} \rightarrow \{1, 2, 3, 4\} = B$, as a correspondence between potatoes; on the right, a function $A \rightarrow B$.

The set $Y^X \subseteq \text{Rel}(X, Y)$ is the set of all *functional* relations, i.e. the set of all relations that are functions.

Let $R : X \rightarrow Y$ be a functional relation. under which condition the relation R^{op} is a function?

EXERCISE OR.12 \square Let (P, \leq) be a partially ordered set; the **Hasse diagram** of P is the directed graph built in the following way:

- there is a vertex for each element of P ;
- there is an edge $q \rightarrow p$ connecting p (below) and q (below) if $p \leq q$ and there is no $x \neq p, q$ such that $p \leq x \leq q$.

Draw the Hasse diagram of the following posets:

- $P = \{a, b, c, d\}$ where $a \leq b, a \leq c, b \leq d, c \leq d$;
- $P = 2^A$ where $A = \{0, 1\}$, $P = 2^B$ where $B = \{0, 1, 2\}$, $P = 2^C$ where $C = \{0, 1, 2, 3\}$; do you see a pattern? Generalize.
- P is the set of divisors of 60, ordered by the relation $a \leq b$ if and only if $b = k \cdot a$ for some $k \in \mathbf{N}$.

DEFINITION 1.7. An **algebraic lattice** (X, \wedge, \vee) is a set X equipped with binary operations \wedge, \vee enjoying the following properties: for all $a, b, c \in X$,

- (commutative) $a \wedge b = b \wedge a$ e $a \vee b = b \vee a$;
- (associative) $a \wedge (b \wedge c) = (a \wedge b) \wedge c$ and $a \vee (b \vee c) = (a \vee b) \vee c$;
- (absorption laws) $a \vee (a \wedge b) = a$ and $a \wedge (a \vee b) = a$.

EXERCISE OR.13 \square Prove that from the absorption laws it follows that both \wedge and \vee are **idempotent** operations: for all $a \in X$, one has

$$a \wedge a = a \quad a \vee a = a. \quad (1.6)$$

EXERCISE OR.14 \square If (X, \wedge, \vee) is an algebraic lattice, we can define a partial order relation on X by saying that $a \leq b$ iff $a \wedge b = a$, or equivalently $a \vee b = b$; prove that for every $a, b, c \in X$ the following inequalities hold



FIGURE 3. A delightfully devilish exercise on order theory.

$$\text{D1) } (a \wedge b) \vee (a \wedge c) \leq a \wedge (b \vee c);$$

$$\text{D2) } a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c).$$

DEFINITION 1.8. An algebraic lattice (X, \wedge, \vee) is called **distributive** if the converse inequality in D1 holds.

EXERCISE OR.15 \square Prove that the following conditions are equivalent for an algebraic lattice (X, \wedge, \vee) :

- X is distributive;
- for every $a, b, c \in X$, $a \vee (b \wedge c) \geq (a \vee b) \wedge (a \vee c)$;
- for every $a, b, c \in X$, $(a \wedge b) \vee (a \wedge c) \vee (b \wedge c) = (a \vee b) \wedge (a \vee c) \wedge (b \vee c)$.

EXERCISE OR.16 \square Let (X, \wedge, \vee) be a distributive lattice with a top and a bottom element; given $a, b, t \in X$, show that there exists at most one $x_{b,t} \in X$ such that $a \wedge x_{b,t} = b$ and $a \vee x_{b,t} = t$. Define the **complement** $\neg a$ of $a \in X$ in a distributive lattice to be $x_{\perp, \top}$. Prove or disprove that $\neg(a \wedge b) = \neg a \vee \neg b$ and $\neg(a \vee b) = \neg a \wedge \neg b$.

EXERCISE OR.17 \square Prove that a lattice (X, \wedge, \vee) is distributive if and only if the following equation is true for every $x, y, z \in X$:

$$(x \wedge y) \vee (y \wedge z) \vee (z \wedge x) = (x \vee y) \wedge (y \vee z) \wedge (z \vee x). \quad (1.7)$$

EXERCISE OR.18 \square (One implication of this exercise is easy; the other is devilishly difficult. See Figure 3.) Prove that a lattice (X, \wedge, \vee) is distributive if and only if the *cancellation properties* hold: given $y, z \in X$, if there exists $x \in X$ such that $x \wedge z = y \wedge z$ and $x \vee z = y \vee z$, then $y = z$.

EXERCISE OR.19 \square Let (X, \wedge, \vee) be a lattice with a top element \top ; assume X is totally ordered by the partial order relation associated to the lattice structure, $x \leq y$ if and only if $x \wedge y = x$, if and only if $x \vee y = y$. Define the binary operation $X \times X \rightarrow X : (x, y) \mapsto x/y$ by saying that $x/y = \top$ if $x \leq y$ and b otherwise. Prove that for every $x, y, z \in X$ we have

$$x \wedge y \leq z \quad \text{if and only if} \quad x \leq y/z. \quad (1.8)$$

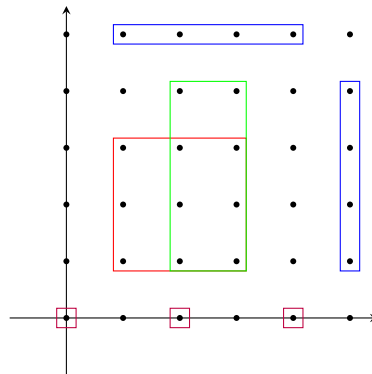


FIGURE 4. The **graph** of a relation depicts the subset $R \subseteq X \times X$ in the ‘plane’ $X \times X$. Here you see the graph of a few relations on the set \mathbf{N} , in different colours.

A relation R on a set X is called

- **reflexive** if for every $x \in X$ we have $(x, x) \in R$;
- **symmetric** if for every $x, y \in X$ we have that, if $(x, y) \in R$, then $(y, x) \in R$;
- **transitive** if for every $x, y, z \in X$ we have that, if $(x, y) \in R$ and $(y, z) \in R$, then $(x, z) \in R$.

EXERCISE OR.20 \square Describe the structure of the set $\text{rRel}(X)$ of reflexive relations on X , of the set $\text{sRel}(X)$ of symmetric relations on X , and of the set $\text{tRel}(X)$ of transitive relations: is the intersection (resp., union) of an arbitrary number of elements in $\text{rRel}(X)$, $\text{sRel}(X)$, $\text{tRel}(X)$, still an element of $\text{rRel}(X)$, $\text{sRel}(X)$, $\text{tRel}(X)$?

EXERCISE OR.21 \square On the graphical representation of a relation. Let X be a set; a relation R on X can be depicted as in Figure 4, as a subset of the Cartesian product $X \times X$. This allows for a graphical representation of properties of R . Show that a relation R is reflexive if and only if it contains the diagonal. Show that a relation R is symmetric if and only if it is symmetric with respect to the diagonal. Find a similar graphical interpretation for the transitive property.

EXERCISE OR.22 \square Find a relation R on a set X that is

- reflexive and symmetric, but not transitive;
- symmetric and transitive, but not reflexive;
- reflexive and transitive, but not symmetric;
- reflexive, but neither symmetric nor transitive;
- symmetric, but neither reflexive nor transitive;
- transitive, but neither reflexive nor symmetric;
- not symmetric, not transitive, not reflexive.

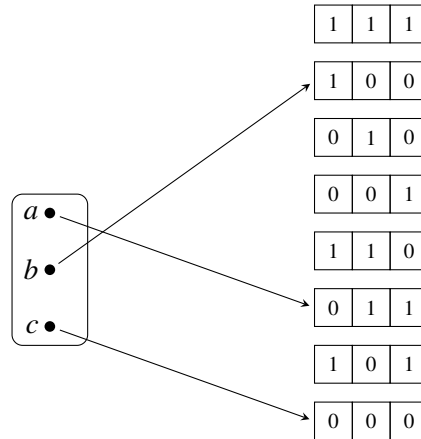


FIGURE 5. The relation $\{a, b, c\} \rightarrow \{1, 2, 3\}$ depicted as a function $\{a, b, c\} \rightarrow \{0, 1\}^3$ sending $a \mapsto \{2, 3\}$, $b \mapsto \{1\}$, $c \mapsto \emptyset$.

(You can choose different sets X for each item of the list.)

EXERCISE OR.23 \square Does the poset of reflexive relations on a set X admit a top element? A bottom element? Same question with the poset of symmetric relations; same question with the poset of transitive relations.

EXERCISE OR.24 \square A generalization of the order on a powerset. Let U be a set called **universe**; consider the set of **multisets** in U , i.e. sequences $\langle x_1, \dots, x_n \rangle$ of **possibly repeated** elements of U , irregardless of order.³ The set of all multisets in a given U is denoted $\mathcal{M}(U)$.

- Prove that the set $\mathcal{M}(U)$ can be identified with the set of all functions $U \rightarrow \mathbf{N}$.

Let $A \in \mathcal{M}(U)$ be a multiset. Define the **counting function** $\epsilon_A : U \rightarrow \mathbf{N}$ of A , mapping each element of U to the number of times that element occurs in A . We can use counting functions to define a ‘generalised inclusion’ relation \leq for multisets. For $A, B \in \mathcal{M}(U)$, we write $A \leq B$ whenever for all $x \in U$, $\epsilon_A(x) \leq \epsilon_B(x)$.

- Prove or disprove that \leq is a partial order relation on $\mathcal{M}(U)$.

Define the following operations on $\mathcal{M}(U)$:

- the **union** $A \vee B$, with associated counting function

$$x \in U \mapsto \max\{\epsilon_A(x), \epsilon_B(x)\}$$

- the **intersection** $A \wedge B$, with associated counting function

$$x \in U \mapsto \min\{\epsilon_A(x), \epsilon_B(x)\}$$

³This means that the multiset $\langle 1, 1, 2 \rangle$ is considered equal to the multiset $\langle 1, 2, 1 \rangle$, but not to the multiset $\langle 1, 2 \rangle$. Compare what happens, instead, with sets.

- the **sum** $A \oplus B$, with associated counting function

$$x \in U \mapsto \epsilon_A(x) + \epsilon_B(x)$$

- the **difference** $A \ominus B$, with associated counting function

$$x \in U \mapsto \epsilon_A(x) - \epsilon_B(x)$$

if this number is nonnegative, and 0 otherwise (more formally, the counting function sends $x \in U$ to $\max\{0, \epsilon_A(x) - \epsilon_B(x)\}$).

EXERCISE OR.25 \square Prove that $(\mathcal{M}(U), \wedge, \vee)$ is a distributive lattice. Prove the **inclusion-exclusion** principle for multisets:

$$A \vee B = (A \oplus B) \ominus (A \wedge B). \quad (1.9)$$

EXERCISE OR.26 \square Count how many reflexive relations exist on a set with 7 elements. Same question, but with symmetric relations. Same question, but with transitive relations.

EXERCISE OR.27 \square Let P be the set of all sets $[n] = \{1, \dots, n\}$, for $0 \leq n \leq 7$, with the convention that $[0]$ is the empty set. Define the relation $[i] \leq [j]$ on P if there is a function $[i] \rightarrow [j]$. Study the order-theoretic properties of the pair (P, \leq) . Is (P, \leq) a poset? If not, what fails? Does P have a bottom element? A top element? Suprema for each subsets?...

EXERCISE OR.28 \square A relation R on X is an **equivalence relation** if it is reflexive, symmetric and transitive; count how many equivalence relations are there on a 4-elements set; count how many equivalence relations are there on a set with 17 elements (but do not ask me to do it at the exercise sessions!).

Denote $\text{eRel}(X)$ the set of equivalence relations on X . Does the poset $(\text{eRel}(X), \subseteq)$ have a top element? A bottom element?

EXERCISE OR.29 \square Which of these relations are equivalence relations on their respective domains?

- The relation R_1 defined on the set \mathbf{Z} of integers, by the rule $(x, y) \in R_1$ if and only if $x - y$ is a multiple of 7.
- The relation R_2 defined on the set $\mathbf{N} = \{0, 1, 2, \dots\}$ of natural numbers, by the rule $(x, y) \in R_2$ if and only if the sum $x + y$ is a prime number.
- The relation R'_2 defined on the set $\mathbf{N} = \{0, 1, 2, \dots\}$ of natural numbers, by the rule $(x, y) \in R'_2$ if and only if the product xy is a prime number.
- The relation R_3 defined on the set A^A of functions $f : A \rightarrow A$, by the rule $(f, g) \in R_3$ if and only if $f \circ g = g \circ f$ (\circ denotes function composition).
- The relation R_4 defined on the set \mathbf{R} of real numbers, by the rule $(x, y) \in R_4$ if and only if the difference $x - y$ is an integer.
- The relation R_5 defined on the set $\mathbf{Z} \times \mathbf{Z}$ of pairs of integers, by the rule $(x_1, x_2), (y_1, y_2) \in R_5$ if and only if $x_1 y_2 = x_2 y_1$.

- The relation R'_5 defined on the set $\mathbf{Z} \times \mathbf{Z}^\times$ of pairs of integers *where the second component is not zero*, by the rule $(x_1, x_2), (y_1, y_2) \in R'_5$ if and only if $x_1 y_2 = x_2 y_1$.

EXERCISE OR.30 \square Given a relation R on X , the equivalence relation **generated** by R is the intersection of all equivalence relations on X containing R ; we denote it \bar{R} . Show that \bar{R} coincides with the subset of $X \times X$ defined ‘**recursively**’ as follows:

- $(x, x) \in \bar{R}$ for each $x \in X$;
- $(y, x) \in \bar{R}$ for each $(x, y) \in \bar{R}$;
- $(x, y) \in \bar{R}$ each time there are $z_0, z_1, \dots, z_n, z_{n+1} \in X$ such that $z_0 = x, z_{n+1} = y$ and $(z_i, z_{i+1}) \in R$ for each $i = 0, \dots, n$.

A **chain** in a poset (P, \leq) consists of a subset $C \subseteq P$ that is totally ordered as a subset of P . An **upper bound** for a subset $C \subseteq P$ of a poset is an element $m \in P$ such that $x \leq m$ for every $x \in C$. A **maximal element** in (P, \leq) is an element $t \in P$ such that, if $t \leq x$ for some $x \in P$, then $t = x$.

- (**Zorn lemma**) Let (X, \leq) be a nonempty poset where every chain $C \subseteq X$ has an upper bound. Then, X admits a maximal element t .
- (**Axiom of choice**) Let I be any set, and X_i a family of sets indexed by I . Let $X := \bigcup_{i \in I} X_i$. Then, there exists a **choice function** for the family $\{X_i\}$, i.e. a function $f : I \rightarrow X$ with the property that for each $i \in I$, $f(i) \in X_i$.

EXERCISE OR.31 \square Prove that Zorn lemma implies the Axiom of choice; prove (more difficult) that the Axiom of Choice implies Zorn lemma. You will have a hard time proving that the Axiom of choice ‘just holds’, meaning that it can be deduced from other statements than the Zorn lemma, in some possibly convoluted way. Do not try!

EXERCISE OR.32 \square Given an equivalence relation R on a set X , define the **equivalence class** of an element x in X to be the set

$$[x]_R := \{y \in X \mid (x, y) \in R\} \quad (1.10)$$

Show that if $[x]_R \cap [y]_R \neq \emptyset$, then $[x]_R = [y]_R$ (two equivalence classes are either disjoint sets, or they coincide). The set X/R is the set of all equivalence classes of X defined by R :

$$X/R := \{[x]_R \mid x \in X\}. \quad (1.11)$$

Define a function $\pi_{/R} : X \rightarrow X/R : x \mapsto [x]_R$, called **projection to the quotient**.

EXERCISE OR.33 \square Let $f : A \rightarrow B$ be a function between sets A, B ; the equivalence relation **induced by f** is the equivalence relation defined by $(a, a') \in R_f$ if and only if $fa = fa'$ (this means: a, a' have the same image under f). Describe the equivalence relations induced by the following functions:

- $f_1 : A \rightarrow B$ sending each element of A in a fixed $b_0 \in B$.
- $f_2 : \mathbf{Z} \rightarrow \mathbf{Z} : n \mapsto 7n$ multiplying an integer n by 7.



FIGURE 6. Zorn lemma, 1998

- $f_3 : \mathbf{R} \rightarrow \mathbf{R}$ taking the *floor* of a real number (the floor of $x \in \mathbf{R}$ is the greatest integer k such that $k \leq x$).
- $f_4 : \mathbf{Q} \rightarrow \mathbf{R}$ multiplying $q \in \mathbf{Q}$ for $\sqrt{27}$.
- $f_5 : A^A \rightarrow A^A : \varphi \mapsto \varphi^{\circ 7}$ that composes $\varphi : A \rightarrow A$ with itself seven times.
- $f_6 : \mathbf{N} \rightarrow \mathbf{N}$ sending m into m_0 , where $m_k \dots m_1 m_0$ is the binary expansion of m in base 2.
- $f_7 : \mathbf{N} \rightarrow \mathbf{N}$ sending n into n^2 .

EXERCISE OR.34 \square Let $f : X \rightarrow Y$ be a function between two sets; let R_f be the equivalence relation generated by f . Prove that f induced an injective function $\bar{f} : X/R_f \rightarrow Y$. Who is the image of \bar{f} ? What can you deduce when f is surjective?

EXERCISE OR.35 \square A **partition** \mathcal{E} of a set X consists of a family of pairwise disjoint subsets $E_i \subseteq X$ (this means that if $i \neq j$, $E_i \cap E_j = \emptyset$) such that $\bigcup E_i = X$. Show that every equivalence relation R on X defines a partition $\mathcal{E}(R)$ di X , and conversely, every partition \mathcal{E} of X defines an equivalence relation $R(\mathcal{E})$ on X , in such a way that $\mathcal{E}(R(\mathcal{E})) = \mathcal{E}$ and $R(\mathcal{E}(R)) = R$.

EXERCISE OR.36 \square Describe X/R for the (equivalence) relations R_1, \dots, R'_5 in [Exercise 29](#); for those that are not equivalence relations, describe X/\bar{R} . Describe X/\bar{R}_{f_i} for f_1, \dots, f_7 in [Exercise 33](#).

EXERCISE OR.37 \square Use Zorn lemma to show that every infinite set X admits a partition in subsets X_α such that every X_α is countable. Use Zorn lemma to show that given *any* two sets X, Y , there exists either an injective function $f : X \rightarrow Y$, or an injective function $g : Y \rightarrow X$.

EXERCISE OR.38 \square Define an equivalence relation Γ on the set $\mathbf{Rel}(X)$ of equivalence relations on X , positing that $(R, S) \in \Gamma$ if and only if there exists a bijection between X/R and X/S . How many elements does the quotient $\epsilon\mathbf{Rel}(X)/\Gamma$ have?

EXERCISE OR.39 \square Prove that every relation $R \in \mathbf{Rel}(A, B)$ defines a Galois connection between the set PA of subsets of A and the set PB of subsets of B : this means that there exists a pair of monotone functions

$$R(-) : PA \rightarrow PB \quad (-)^R : PB \rightarrow PA \quad (1.12)$$

such that $V \subseteq R^R U$ if and only if $U \subseteq V^R$, for each $V \in PB$ and $U \in PA$.

EXERCISE OR.40 \square On a graphical representation for relations. Show that a relation $R : X \rightarrow Y$ can be represented equivalently as follows: a pair of functions

$$X \xleftarrow{u} R \xrightarrow{v} Y \quad (1.13)$$

such that the function $R \rightarrow X \times Y$ defined as $r \mapsto (ur, vr) \in X \times Y$ is injective. In such a representation, a relation is denoted (R, u, v) .

EXERCISE OR.41 \square Define the **composition** of two relations $(R, u, v) : X \rightarrow Y$, $(S, w, t) : Y \rightarrow Z$ as

$$R \circ S := \{(x, z) \mid \exists y \in Y. (x, y) \in R, (y, z) \in S\} \subseteq X \times Z. \quad (1.14)$$

EXERCISE OR.42 \square Define a couple of functions h, k ,

$$R \xleftarrow{h} R \circ S \xrightarrow{k} S \quad (1.15)$$

so that $w(k(q)) = v(h(q))$ for every $q \in R \circ S$. This means that the function $wk = w \circ k$ coincides with the function $vh = v \circ h$; a graphical way to represent such a situation is to depict w, k, v, h as edges of a graph, in this case as a square

$$\begin{array}{ccc} R \circ S & \xrightarrow{k} & S \\ h \downarrow & & \downarrow w \\ R & \xrightarrow{v} & Y \end{array} \quad (1.16)$$

and to declare that the square **commutes** when $wk = vh$.

EXERCISE OR.43 \square Prove that given any other commutative square

$$\begin{array}{ccc} E & \xrightarrow{s} & S \\ r \downarrow & & \downarrow w \\ R & \xrightarrow{v} & Y \end{array} \quad (1.17)$$

There is a unique function $(r/s) : E \rightarrow R \circ S$ with the property that $k \circ (r/s) = s$ and $h \circ (r/s) = r$. This is called the **universal property** of $R \circ S$.

EXERCISE OR.44 \square Let $\Delta_X = \{(x, x) \mid x \in X\}$. Prove that for every relation $(R, u, v) : X \rightarrow Y$ one has $\Delta_X \circ R = R$ and $R \circ \Delta_Y = R$. The relation Δ_X on a set X plays the role of **identity** for the composition operation on relations.

Let $(R, u, v) : X \rightarrow Y$, $(S, w, t) : Y \rightarrow Z$ be relations; prove that $(R \circ S)^{\text{op}} = S^{\text{op}} \circ R^{\text{op}}$.

EXERCISE OR.45 \square Let R be a relation on a set X ; define the relation \tilde{R} to be $R \cup R^{\text{op}} \cup \Delta$, where Δ is the diagonal relation, as above. Prove that \tilde{R} is the smallest reflexive and symmetric relation containing R .

EXERCISE OR.46 \square Let X be a set and R a relation on X . Show that the transitive closure of R , i.e. the smallest transitive relation containing R , coincides with the set

$$\bigcup_{n=1}^{\infty} R^{\circ n} := R \cup (R \circ R) \cup (R \circ R \circ R) \cup \dots \quad (1.18)$$

Prove that the equivalence relation generated by R , as defined in [Exercise 30](#), is the transitive closure of \tilde{R} as defined above.

EXERCISE OR.47 \square Let R, S be equivalence relations on a set X , such that $R \circ S = S \circ R$. Prove that $R \circ S$ is an equivalence relation on X , and in fact it is the join $R \vee S$ of $\{R, S\}$ in the poset $(\text{eRel}(X), \subseteq)$.

EXERCISE OR.48 \square Let R, S, T, T' be relations on a set X . For each of the following items, prove it if they are true, or provide a counterexample if they are false.

- If $R \subseteq S$, then $R \circ T \subseteq S \circ T$ and $T \circ R \subseteq T \circ S$;
- $R \circ (T \cap T') = (R \circ T) \cap (R \circ T')$ and $(T \cap T') \circ R = (T \circ R) \cap (T' \circ R)$;
- $R \circ (T \cup T') = (R \circ T) \cup (R \circ T')$ and $(T \cup T') \circ R = (T \circ R) \cup (T' \circ R)$;
- $R \subseteq S$ if and only if $R^{\text{op}} \subseteq S^{\text{op}}$;
- $(T \cup T')^{\text{op}} = T^{\text{op}} \cup (T')^{\text{op}}$ and $(T \cap T')^{\text{op}} = T^{\text{op}} \cap (T')^{\text{op}}$;

EXERCISE OR.49 \square Prove **Szpilrajn extension theorem**: let (P, \leq) be a poset; prove that there exists a relation \leq extending \leq , i.e. such that $x \leq y$ implies $x \leq y$, which is also a total order, i.e. either $x \leq y$ or $y \leq x$.

EXERCISE OR.50 \square Let X, Y be posets; define a relation \leq on the cartesian product $X \times Y$ by saying

$$(x, y) \leq (x', y') \iff x \leq x' \text{ in } X, y \leq y' \text{ in } Y. \quad (1.19)$$

Show that the projection functions $\pi_X : X \times Y \rightarrow X$ and $\pi_Y : X \times Y \rightarrow Y$ are monotone maps, when $X \times Y$ is equipped with this order relation. Show that the diagonal map

$$d_X : X \longrightarrow X \times X \quad (1.20)$$

sending $x \in X$ into $(x, x) \in X \times X$ is monotone.

Now, let (X, \vee, \wedge) be an algebraic lattice. Show that there are Galois connections

$$_ \vee _ : X \times X \rightleftarrows X : d_X \quad d_X : \rightleftarrows X \times X : _ \wedge _ \quad (1.21)$$

if $\vee, \wedge : X \times X \rightarrow X$ are respectively the sup and inf operation on X .

EXERCISE OR.51 \square Let $f : X \rightarrow Y$ be a monotone map between posets; assume f fits in a Galois connection $f \dashv u$, where $u : Y \rightarrow X$. Prove that f preserves the bottom element of X , if it exists; prove that u preserves the top element of Y , if it exists. Prove that $f(\sup S) = \sup f(S)$ for every subset $S \subseteq X$ for which $\sup S$ exists; in particular, $f(x \vee x') = fx \vee fx'$ for every $x, x' \in X$ admitting a sup in X .

EXERCISE OR.52 \square Let (P, \leq) be a poset; a **down-set** in P is a subset $S \subseteq P$ such that if $x \in S$ and $y \leq x$, then $y \in S$. Let DP be the set of all down-sets of P ; for each $x \in P$, define the downset $\downarrow x$ **generated** by x as the set $\{y \in P \mid y \leq x\}$. Show that the map $x \mapsto \downarrow x$ is a monotone, injective map $\downarrow(-) : P \rightarrow DP$, i.e. show that

$$a \leq b \implies \downarrow a \leq \downarrow b. \quad (1.22)$$

Now let (P, \leq) admit suprema for all down-sets; show that $S \mapsto \bigvee S$ defines a monotone map $DP \rightarrow P$, and show that

$$\bigvee S \leq x \iff S \subseteq \downarrow x \quad (1.23)$$

for every $S \in DP$ and $x \in P$.

EXERCISE OR.53 \square Define $\downarrow U := \bigcup_{x \in U} \downarrow x$ for every $U \subseteq P$. Prove that U is a down-set if and only if $\downarrow U = U$.

EXERCISE OR.54 \square In the same notation of [Exercise 52](#), we call a lattice (P, \leq) admitting all suprema **completely distributive** if there exists a monotone map $\Downarrow(-) : P \rightarrow DP$, such that

$$\Downarrow a \subseteq S \iff a \leq \bigvee S \quad (1.24)$$

for all $a \in P$ and $S \in DP$. Show that the only possible definition for $\Downarrow a$ is as the set $\bigcap \{S \in DP \mid a \leq \bigvee S\}$. Define a relation \ll on P as follows: $x \ll a$ if and only if $x \in \Downarrow a$. Prove that for any $a, b, x \in P$ one has

- if $x \ll a$ and $a \leq b$, then $x \ll b$;
- $a \leq \bigvee \{x \in P \mid x \ll a\}$.

EXERCISE OR.55 \square Let (P, \wedge, \vee) be a lattice. Prove that given a set I of indices and a family $\{S_i \mid i \in I\}$ of down-sets, the following equality holds:

$$\bigvee \left(\bigcap_{i \in I} S_i \right) = \bigwedge_{i \in I} s_i \quad (1.25)$$

where $s_i := \bigvee S_i$. Let $\{A_i \mid i \in I\}$ be a family of subsets of P ; prove that the above equality holds if and only if it holds for downsets, i.e.

$$\bigvee \left(\bigcap_{i \in I} \downarrow A_i \right) = \bigwedge_{i \in I} a_i \quad (1.26)$$

where $a_i := \bigvee A_i$.

EXERCISE OR.56 \square On free distributive lattices on a set. Let A be a set; this series of exercises is intended to build the **free distributive lattice** on A , i.e. a lattice $(L[A], \wedge, \vee)$ enjoying the following properties:

- $(L[A], \wedge, \vee)$ is distributive, and there exists a function $\eta : A \rightarrow L[A]$;
- For every other distributive lattice (X, \wedge, \vee) , every function $f : A \rightarrow X$ can be extended to a unique lattice homomorphism $\bar{f} : (L[A], \wedge, \vee) \rightarrow (X, \wedge, \vee)$ such that $\bar{f} \circ \eta = f$.

DEFINITION 1.9. Given a set A , an **antichain**, or a **clutter**, or an **irredundant family** in A is a family of subsets $(E_\alpha \subseteq A \mid \alpha \in I)$ with the property that given $\alpha \neq \beta$ one has $E_\alpha \not\subseteq E_\beta$ (in words: none of the E_α contains one of the E_β as a subset). Define the set $L[A]$ to be the set of finite sets of finite antichains in A , i.e. the set of all

$$\{X_1, \dots, X_n\} \quad (1.27)$$

where each X_i is a finite antichain in A .

EXERCISE OR.57 \square Prove that given two antichains $X = (X_1, \dots, X_n), Y = (Y_1, \dots, Y_m)$, one can obtain an antichain from $X \cup Y$ by ‘removing repetitions’: the set-theoretic union $\{X_1, \dots, X_n, Y_1, \dots, Y_m\}$ of the two antichains is not, in general, an antichain (find a counterexample for two antichains on $\{1, 2, 3, 4, 5\}$), but we can define a ‘one-step reduction’ operation \rightsquigarrow by declaring that

$$\{X_1, \dots, X_n, Y_1, \dots, Y_m\} \rightsquigarrow \{X_1, \dots, X_n, Y_1, \dots, \widehat{Y}_j, \dots, Y_m\} \quad (1.28)$$

if there exists an index i such that either $X_i \subseteq Y_j$ or $X_i \supseteq Y_j$. The relation of reduction \rightsquigarrow^* now is defined to be the transitive closure of \rightsquigarrow ; We denote the binary operation of union $(X, Y) \mapsto X \cup Y$, followed by complete reduction \rightsquigarrow^* as $X \# Y$.

Given two elements $X = (X_1, \dots, X_n)$ and $Y = (Y_1, \dots, Y_m)$ in $L[A]$, define the join of two elements in $L[A]$ as $X \vee Y := X \# Y$ and the meet $X \wedge Y := \{X_i \# Y_j \mid 1 \leq i \leq n, 1 \leq j \leq m\}$.

EXERCISE OR.58 \square Prove that these definitions equip $L[A]$ with the structure of a (distributive) lattice. Prove that $L[A]$ is the free distributive lattice on the set A . [Hint: interpret a generic element

$$\{\{a_{1,1}, \dots, a_{1,i_1}\}, \dots, \{a_{n,1}, \dots, a_{n,i_n}\}\} \quad (1.29)$$

of $L[A]$ as the element $(a_{1,1} \wedge \dots \wedge a_{1,i_1}) \vee \dots \vee (a_{n,1} \wedge \dots \wedge a_{n,i_n})$.]

EXERCISE OR.59 \square On filters. Let (P, \leq) be a poset. A nonempty subset A of P is called **down-directed** if for all $x, y \in A$ there exists a $z \in A$ such that $z \leq x$ and $z \leq y$. A subset A of P is called a **filter** if it is down-directed and up-closed; a filter is called **proper** if $A \neq P$. Show that if P has finite meets, a filter A is a subset of P such that

- for all $x, y \in A, x \wedge y \in A$;
- $\top \in A$;

- if $x \in A$ and $x \leq y$, then $y \in A$.

Let $A \subseteq P$ be any down-directed subset; show that $\uparrow A = \bigcup_{x \in A} \uparrow x$, where $\uparrow x := \{y \in P \mid x \leq y\}$ is a filter, called the filter **generated** by A . In particular, every $\uparrow x$ is a filter, the **principal filter** generated by x . Show that for every $A \subseteq X$, and every function $f : X \rightarrow Y$ one has $f(\uparrow A) = \uparrow(fA)$.

EXERCISE OR.60 \square Let $f : X \rightarrow Y$ be a function; let $\mathfrak{a} \subseteq PX$ be a filter in the powerset of X (a common shorthand for this is to say that \mathfrak{a} is a filter on X). Show that

$$\{B \subseteq Y \mid f^{\leftarrow} B \in \mathfrak{a}\} \quad (1.30)$$

defines a filter in PY , called the **direct image** filter on Y under f .

Let $\mathfrak{b} \subseteq PY$ be a filter in PY ; show that

$$\{A \subseteq X \mid \exists B \in \mathfrak{b} : f^{\leftarrow} B \subseteq A\} \quad (1.31)$$

is a filter in PX , called the **inverse image filter** on X under f .

EXERCISE OR.61 \square Let X be a set, and \mathfrak{A} a filter on the set of filters on X ; define a filter $\sum \mathfrak{A}$ (called the **Kowalski sum** of \mathfrak{A}) on X as follows: $A \subseteq X$ is an element of $\sum \mathfrak{A}$ if and only if the set of filters on X that are also filters on A belongs to \mathfrak{A} . Prove that $\sum \mathfrak{A}$ is indeed a filter on X . Prove that $A \in \sum \mathfrak{A}$ if and only if the set of filters \mathfrak{a} on X that contain A as an element belongs to \mathfrak{A} .

DEFINITION 1.10. An **ultrafilter** \mathfrak{x} on a set X is a maximal element within the set of proper filters on X , ordered by inclusion; i.e. \mathfrak{x} is a proper filter on X such that, if \mathfrak{a} is a proper filter on X containing \mathfrak{x} , then $\mathfrak{x} = \mathfrak{a}$.

EXERCISE OR.62 \square Prove that the following conditions are equivalent for a filter \mathfrak{x} on X .

- \mathfrak{x} is an ultrafilter on X ;
- for all $A, B \subseteq X$, if $A \cup B \in \mathfrak{x}$ then either $A \in \mathfrak{x}$ or $B \in \mathfrak{x}$;
- for every $A \subseteq X$, either $A \in \mathfrak{x}$ or $A^c = X \setminus A \in \mathfrak{x}$.

EXERCISE OR.63 \square Let $f : X \rightarrow Y$ be a function between sets, and let \mathfrak{x} be an ultrafilter on X ; prove that the direct image of \mathfrak{x} under f is again an ultrafilter. As a corollary, when $f : X \hookrightarrow Y$ is the inclusion of a subset, for every ultrafilter \mathfrak{y} on Y such that $Y \in \mathfrak{y}$ the direct image filter $\mathfrak{y}|_X = \{U \cap X \mid U \in \mathfrak{y}\}$ is an ultrafilter on X .

EXERCISE OR.64 \square Let X be a set. Prove that for every element $x \in X$, the principal filter $\uparrow x$ is an ultrafilter on X .

EXERCISE OR.65 \square Prove that if \mathfrak{X} is an ultrafilter on the set of ultrafilters on X , then the Kowalski sum $\sum \mathfrak{X}$ is an ultrafilter on X .

Defining ultrafilters of different kinds requires the Axiom of Choice, in the form of Zorn lemma.

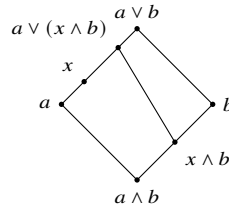


FIGURE 7. A minimal example of non-modular lattice, by construction.

EXERCISE OR.66 \square Use Zorn lemma to prove that given a set X , every proper filter \mathfrak{a} on X is contained in an ultrafilter. As a consequence, prove that for every filter \mathfrak{b} and every filter $\mathfrak{a} \subset \mathfrak{b}$, there is an ultrafilter \mathfrak{x} on X such that $\mathfrak{a} \subseteq \mathfrak{x}$ but $\mathfrak{b} \not\subseteq \mathfrak{x}$. As a consequence, prove that given a filter \mathfrak{a} on a set X , \mathfrak{a} is the intersection of all ultrafilters on X containing \mathfrak{a} .

2. More on ordered sets

In the following, we will always denote (L, \wedge, \vee) an algebraic lattice. Given two elements $a, b \in L$ we define the **interval** $[a, b]$ as the subset $\{y \mid a \leq y \leq b\}$ of L .

EXERCISE PO.1 \square Show that the set $\mathbf{I}(L)$ of intervals in L becomes a lattice, defining the lattice operations \wedge, \vee on $\mathbf{I}(L)$.

EXERCISE PO.2 \square Show that the pair of functions

$${}_-\vee b : [a \wedge b, a] \xleftrightarrow{\quad} [b, a \vee b] : {}_-\wedge a \tag{2.1}$$

defines a Galois connection; we say that a lattice L is **modular** when ${}_-\vee b \dashv {}_-\wedge a$ is an order-isomorphism between $[b, a \vee b]$ and $[a \wedge b, a]$.

EXERCISE PO.3 \square Show that the following conditions are equivalent for a lattice L :

- L is modular;
- for every $a, b, x \in L$, if $a \leq b$ then $(x \vee a) \wedge b \leq (x \wedge b) \vee a$ (and thus the equality holds);
- every interval $[a, b] \subseteq L$ has the following property: every $c \in [a, b]$ has at most one complement c' in $[a, b]$.

EXERCISE PO.4 \square Show that if L is complemented and modular, then every interval in L is also complemented.

EXERCISE PO.5 \square Prove that every distributive lattice (cf. [Definition 1.8](#)) is modular.

DEFINITION 2.1. Let $I, J \in \mathbf{I}(L)$ be two intervals in the lattice L ; we say that I and J are **similar**, and we write $I \asymp J$, if there exist $a, b \in L$ such that one of the intervals is $[a \wedge b, a]$ and the other is $[b, a \vee b]$. This defines a relation ${}_-\asymp {}_-$ on $\mathbf{I}(L)$. Clearly, the ${}_-\asymp {}_-$ relation is not transitive.

EXERCISE PO.6 \square Denote \approx^P the transitive closure of \approx ; is \approx^P an equivalence relation? We say that $I, J \in \mathbf{I}(L)$ are **projective** when $I \approx^P J$. Show that projective intervals are order-isomorphic; is the converse true?

Denote respectively $0_L, 1_L$ the bottom and top element of the lattice L ; two finite intervals with endpoints a, b , i.e. two chains

$$\begin{aligned}\underline{u} &= \{a = u_0 \leq u_1 \leq \cdots \leq u_m = b\} \\ \underline{v} &= \{a = v_0 \leq v_1 \leq \cdots \leq v_n = b\}\end{aligned}$$

of elements of L are called **equivalent** if $m = n$ and there exists a permutation $\pi : \{1, \dots, n\} \rightarrow \{1, \dots, n\}$ such that for every $i = 1, n$ one has $[u_{i-1}, u_i] \approx^P [v_{\pi i-1}, v_{\pi i}]$ (see [Exercise 6](#)). We denote this situation as $\underline{u} \approx \underline{v}$.

EXERCISE PO.7 \square Show that \approx is an equivalence relation.

A **refinement** of a finite interval \underline{u} as above can be obtained by inserting further elements in the chain: more formally, given $\underline{u}, \underline{v}$ as above, we say that the interval \underline{v} *refines* the interval \underline{u} (and we write $\underline{u} \triangleleft \underline{v}$) if $n \geq m$ and $\{u_0, \dots, u_m\} \subseteq \{v_0, \dots, v_n\}$. Prove that refinement relation \triangleleft is a partial order on $\mathbf{I}(L)$.

EXERCISE PO.8 \square Prove the **Schreier refinement lemma**: any two finite chains $\underline{u}, \underline{v}$ between the same pair of elements a, b in a modular lattice L admit equivalent refinements.

More formally: given $\underline{u}, \underline{v}$ as above we can find intervals $\underline{x}, \underline{y}$ such that the following conditions are satisfied:

- $\underline{u} \triangleleft \underline{x}$;
- $\underline{v} \triangleleft \underline{y}$;
- $\underline{x} \approx \underline{y}$.

DEFINITION 2.2. A **composition series** between $a, b \in L$ is a chain

$$a = u_0 \leq u_1 \leq \cdots \leq u_m = b \tag{2.2}$$

which has no refinement, except by introducing repetitions of some of the given elements a_i . The integer m is the length of the chain.

EXERCISE PO.9 \square Prove the **Jordan-Hölder theorem** on composition series: any two composition series between the same pair of elements a, b in a modular lattice are equivalent.

DEFINITION 2.3. A modular lattice L is of **finite length** if there is a composition chain between 0_L and 1_L ; we define the *length* of L to be the length of such a composition chain (by the Jordan-Hölder theorem, this is well-defined).

EXERCISE PO.10 \square Prove that in a modular lattice of finite length, every chain

$$[a, b] = a = u_0 \leq u_1 \leq \cdots \leq u_m = b \tag{2.3}$$

can be refined to a composition series.

DEFINITION 2.4 (Noetherian and Artinian lattices). A lattice L is **Noetherian**, or *satisfies the ascending chain condition* if there is no infinite ascending sequence

$$a_0 < a_1 < a_2 < \dots \quad (2.4)$$

of distinct elements. Dually, L is called **Artinian**, or it *satisfies the descending chain condition* if there is no infinite descending sequence

$$\dots < a_2 < a_1 < a_0 \quad (2.5)$$

of distinct elements.

EXERCISE PO.11 \square Prove that L is Noetherian (resp., Artinian) if and only if every nonempty subset S of L has a maximal (resp., minimal) element. [Assuming Noetherianity, you will need the Zorn lemma to prove the existence of a maximal element.]

EXERCISE PO.12 \square Deduce from the previous exercise that a modular lattice is of finite length if and only if it is both noetherian and Artinian.

EXERCISE PO.13 \square Let a be an element of a modular lattice L . Then L is Noetherian (resp., Artinian) if and only if both intervals $[0, a]$ and $[a, 1]$ are Noetherian (resp., Artinian).

EXERCISE PO.14 \square Prove **Knaster-Tarski** fixpoint theorem: every monotone endofunction $f : L \rightarrow L$ of a complete lattice has at least a fixpoint. In fact, the set of fixpoints of f in L also forms a complete lattice, so that f has a *least* and a *greatest* fixpoint.

A converse of this theorem also holds: if every monotone function $f : L \rightarrow L$ on a lattice L has a fixpoint, then L is a complete lattice.

DEFINITION 2.5. A totally ordered set W is called **well-ordered** if every nonempty subset $S \subseteq W$ admits a least element.

Well-ordered sets serve the purpose to classify order types: every well-ordered set is completely described by a certain distinguished element in its order-isomorphism class, an *ordinal number*. Roughly speaking, an ordinal number describes how you can line up the elements of a set so that the relation \leq is a well-order according to the definition above. Assuming the axiom of choice, it is possible to prove something quite strong: *every* set, no matter what is its internal structure, can be turned into a well-order. However, as it is customary with AoC-dependent theorems, there is no way to make this definition an explicit construction.

EXERCISE PO.15 \square Let (W, \leq) be a well-ordered set, and let $f : W \rightarrow W$ be a monotone endofunction; show that for every $a \in W$, $a \leq f(a)$.

EXERCISE PO.16 \square Show that a well-ordered set is *rigid*: the only order-automorphism of a well-ordered set W is the identity.

Show that the isomorphisms between ordered sets are rigid: if there exists an order-isomorphism $f : V \rightarrow W$ between two well-ordered sets, then f is unique.

EXERCISE PO.17 \square Show that well-ordered sets are *trichotomous*: given two well-ordered sets V, W , exactly one of the following three cases holds:

- V is isomorphic to W ; in this case, we write that the *order type* of V $o(V)$ is the same of the order type $o(W)$ of W .
- V is isomorphic to an initial segment of W ; in this case, we write $o(V) \leq o(W)$.
- W is isomorphic to an initial segment of V ; in this case, we write $o(W) \leq o(V)$.

EXERCISE PO.18 \square Assume that each well-ordered set W is assigned its ordinal $\alpha = o(W)$. Show that the assignment is well-defined and that \leq totally orders the class Ord of all ordinals, and in fact, the pair (Ord, \leq) is a well-order...

... But it's not an ordinal: it lacks the property of being a *set*.

We want to have a more hands-on model of ordinals to work with: 'isomorphism classes of well-ordered sets' is a bit too elusive of a definition.

DEFINITION 2.6. A set X is transitive if each element of X is also a subset of X :

$$x \in X \Rightarrow x \subseteq X. \quad (2.6)$$

Equivalently, X is transitive if $\bigcup X \subseteq X$, or $X \subseteq 2^X$.

A set X is an ordinal if it is transitive and well-ordered by the relation \in (it's an \in -wo set).

Historically, an ordinal is denoted with a lowercase Greek letter

$$\alpha, \beta, \gamma, \dots, \omega \quad (2.7)$$

or with 'decorated' versions thereof: α_0, γ' , etc.

EXERCISE PO.19 \square Prove that

- The two definitions of ordinals do not conflict: given a well-ordered set W , there exists *exactly one* transitive and \in -wo set X with an order isomorphism $W \cong X$;
- $0 := o(\emptyset)$ is an ordinal;
- if α is an ordinal, and $\beta \in \alpha$, then β is an ordinal;
- if α, β are ordinals and $\alpha \subsetneq \beta$ then $\alpha \in \beta$;
- if α, β are ordinals then either $\alpha \subseteq \beta$ or $\beta \subseteq \alpha$.

EXERCISE PO.20 \square Prove that

- for each ordinal α , $\alpha = \{\beta \mid \beta < \alpha\}$;
- if C is a nonempty class of ordinals, then $\bigcap C$ is an ordinal, $\bigcap C \in C$ and $\bigcap C = \inf C$;
- if X is a nonempty *set* of ordinals, then $\bigcup X$ is an ordinal, and $\bigcup X = \sup X$;
- for every ordinal α , the set $\alpha^+ := \alpha \cup \{\alpha\}$ is an ordinal, and $\alpha^+ = \inf\{\beta \mid \beta > \alpha\}$.

DEFINITION 2.7. A **successor ordinal** is an ordinal α such that there exists an ordinal β for which $\alpha = \beta^+$; α is then the *successor* of β , and it's usually written $\alpha = \beta + 1$.

A **limit ordinal** is an ordinal α such that $\alpha = \bigcup \alpha = \sup\{\beta \mid \beta < \alpha\}$.

We consider 0 to be a limit ordinal and define $0 = \sup \emptyset$.

EXERCISE PO.21 \square Prove that α is a limit ordinal if and only if for every β , $\beta < \alpha$ implies $\beta + 1 < \alpha$.

The possibility to build limit ordinals relies on the axiom of infinity of ZF; in particular, one can

EXERCISE PO.22 \square Prove that if a set X is inductive,⁴ then $X \cap \text{Ord}$ is also inductive, and the set $\omega = \bigcap \{X \mid X \text{ is inductive}\}$ is the least nonzero limit ordinal.

Stripped of its order, the ordinal ω is just ‘the set of natural numbers’ $\{0, 1, 2, \dots\}$.

Without the axiom of infinity, the only ordinals to which we have access are $0 = \sup \emptyset$ and the *finite* ones $0^+ = 1 := \{\emptyset\}$, $2 = 1^+ = 0^{++} := \{\emptyset, \{\emptyset\}\}$, ..., $n + 1 := n^+$.

The power of ordinals lies in the fact that one can provide inductive definitions. A ‘definition by transfinite recursion’ usually takes the following form: let \mathcal{K} be a class, then a function $h : \text{Ord} \rightarrow \mathcal{K}$, called a **transfinite sequence**, is uniquely determined by

- the definition of the ‘base’ of the induction: a certain element h_0 of the class \mathcal{K} ;
- the definition of the ‘successor’ step: a way to ‘compute’ $h_{\alpha+1}$ in terms of all $h_0, h_1, \dots, h_\alpha$;
- the definition of the ‘limit’ step: a way to ‘compute’ h_λ , when λ is a limit ordinal, in terms of all $h_\beta, \beta < \lambda$.

The most profitable way to employ such definitions is to define *arithmetic operations* on Ord: sum, product, and exponentiation.

Each such binary operation is specified by a recursion on the second argument, i.e. (for example for addition) there will be

- a base definition of what it means $\beta + 0$;
- a successor definition of what it means $\beta + (\alpha^+)$ in terms of $(\beta + 0, \beta + 1, \dots, \beta + \alpha)$;
- a limit definition of what it means $\beta + (\sup\{\theta \mid \theta < \lambda\})$.

DEFINITION 2.8 (Ordinal sum). The operation of ordinal sum is defined as follows: fix any ordinal β ; then

- (base) $\beta + 0 := \beta$;
- (successor) $\beta + \alpha^+ := (\beta + \alpha)^+$;
- (limit) $\beta + \bigcup \{\alpha \mid \alpha < \lambda\} := \bigcup \{\beta + \alpha \mid \alpha < \lambda\}$.

EXERCISE PO.23 \square Prove that the ordinal sum is associative; prove that it is commutative when restricted to finite ordinals $0, 1, 2, \dots$. Can you also prove that it is *not* commutative,

⁴A set X is *inductive* if for every $x \in X$ also $x^+ = x \cup \{x\} \in X$. Since $\emptyset \in X$ for every X , a set X is inductive if it contains $\{\emptyset\}$, and thus also $\{\emptyset, \{\emptyset\}\}$, and thus also... The axioms of infinity says that there exists at least an inductive set.

because $\omega + 1$ is not order-isomorphic to $1 + \omega$? (Bonus points: prove that in fact $1 + \omega = \omega$ as ordinals.) Is it true that $0 + \alpha = \alpha$ for each ordinal α ?⁵

DEFINITION 2.9 (Ordinal product). The operation of ordinal product is defined as follows: fix any ordinal β ; then

- (base) $\beta \cdot 0 := 0$;
- (successor) $\beta \cdot \alpha^+ := (\beta \cdot \alpha) + \beta$;
- (limit) $\beta \cdot \bigcup\{\alpha \mid \alpha < \lambda\} := \bigcup\{\beta \cdot \alpha \mid \alpha < \lambda\}$.

EXERCISE PO.24 \square Prove that the ordinal product is associative; prove that it is commutative when restricted to finite ordinals $0, 1, 2, \dots$. Can you also prove that it is *not* commutative, because $\omega \cdot 2$ is not order-isomorphic to $2 \cdot \omega$? (Bonus points: prove that in fact $2 \cdot \omega = \omega$ as ordinals). Is it true that $1 \cdot \alpha = \alpha$ for each ordinal α ?

DEFINITION 2.10 (Ordinal exponentiation). The operation of ordinal exponentiation is defined as follows: fix any ordinal β ; then

- (base) $\beta^0 := 1$;
- (successor) $\beta^{(\alpha^+)} := (\beta^\alpha) \cdot \beta$;
- (limit) $\beta^{\bigcup\{\alpha \mid \alpha < \lambda\}} := \bigcup\{\beta^\alpha \mid \alpha < \lambda\}$.

(When it could be potentially confusing to adopt the superscript notation β^λ we might write $\exp(\beta, \lambda)$. So, $\exp(\beta, \alpha^+) = \exp(\beta, \alpha) \cdot \beta$ and $\exp(\beta, \sup\{\alpha \mid \alpha < \lambda\}) = \sup\{\exp(\beta, \alpha) \mid \alpha < \lambda\}$.)

Ordinal exponentiation is *not* associative, it's not commutative, it doesn't have a unit element.

EXERCISE PO.25 \square Prove the three identities

$$\beta^1 = \beta \quad \beta^{\alpha+\gamma} = \beta^\alpha \cdot \beta^\gamma \quad (\beta^\alpha)^\gamma = \beta^{\alpha \cdot \gamma} \quad (2.8)$$

for ordinal exponentiation, valid for all ordinals β, α, γ .

The operations of sum, product and exponentiation of ordinals can be succinctly constructed using just the successor function $\alpha \mapsto \alpha^+$ in a clever way. This is ultimately based on the fact that given the ordinal β , each function $\beta + _$, $\beta \cdot _$, $\exp(\beta, _)$ can be regarded as a transfinite sequence $\text{Ord} \rightarrow \text{Ord}$.

EXERCISE PO.26 \square Let $h : \text{Ord} \rightarrow \text{Ord}$ be a function; define the α -th iterate of h by transfinite recursion as follows:

$$h^0 := \text{id} \quad h^{\alpha+1} = h^\alpha \circ h \quad h^\lambda := \beta \mapsto \sup\{h^\alpha \beta \mid \alpha < \lambda\}. \quad (2.9)$$

Prove that

- $\beta + \alpha = (_ + 1)^\alpha \beta$;

⁵Find an intuitive argument first, and then formalise the fact that $1 + \omega = \sup\{1 + \alpha \mid \alpha < \omega\}$ equals ω . Watch this video [watch?v=YApat9UmUNg](https://www.youtube.com/watch?v=YApat9UmUNg) for inspiration.

- $\beta \cdot \alpha = (_ + \beta)^\alpha 0$;
- $\exp(\beta, \alpha) = (_ \cdot \beta)^\alpha 1$.

The point is that now we can apply the same recursive definition for the iterates of the exp function, building $\omega^\omega, \omega^{\omega^\omega}, \omega^{\omega^{\omega^\omega}} \dots$ more formally, we define a transfinite sequence

$$h_0 := \omega \quad h_{\alpha+1} := \omega^{h_\alpha} \quad h_\lambda := \sup\{h_\beta \mid \beta < \lambda\} \quad (2.10)$$

Now, what is h_ω exactly? It is a tower of ω 's that is ω steps high.

EXERCISE PO.27 \square Prove that h_ω is a fixpoint for the function $x \mapsto \omega^x$, or in other words a solution to the equation

$$\omega^x = x \quad (2.11)$$

in Ord.

The ordinal h_ω , that is more often written ϵ_0 , is then equal to $\sup\{\omega, \omega^\omega, \omega^{\omega^\omega}, \omega^{\omega^{\omega^\omega}}, \dots\}$. The notation suggests that there exists an entire hierarchy of greater ϵ 's, $\epsilon_1, \epsilon_2, \dots$ where

$$\epsilon_1 = \sup\{\epsilon_0 + 1, \omega^{\epsilon_0+1}, \omega^{\omega^{\epsilon_0+1}}, \omega^{\omega^{\omega^{\epsilon_0+1}}}, \dots\} \quad (2.12)$$

and more generally

$$\epsilon_{\alpha+1} = \sup\{\epsilon_\alpha + 1, \omega^{\epsilon_\alpha+1}, \omega^{\omega^{\epsilon_\alpha+1}}, \dots\} \quad (2.13)$$

EXERCISE PO.28 \square Prove that

- $\epsilon_1 = \sup\{0, 1, \epsilon_0, \epsilon_0^{\epsilon_0}, \epsilon_0^{\epsilon_0^{\epsilon_0}}, \dots\}$;
- $\epsilon_{\alpha+1} = \sup\{0, 1, \epsilon_\alpha, \epsilon_\alpha^{\epsilon_\alpha}, \epsilon_\alpha^{\epsilon_\alpha^{\epsilon_\alpha}}, \dots\}$;
- $\epsilon_\omega = \sup\{\epsilon_0, \epsilon_1, \epsilon_2, \dots\}$.

It is now possible to define $\epsilon_{\omega+1}, \epsilon_{\omega+2}, \dots, \epsilon_{\omega+\omega} = \epsilon_{2 \cdot \omega}, \dots, \epsilon_{\omega \cdot \omega} = \epsilon_{\omega^2}, \epsilon_{\omega^3}, \dots, \epsilon_{\omega^\omega}, \dots$ up to ϵ_{ϵ_0} ; what next?

Isn't that obvious? By transfinite recursion, $h_0 = \epsilon_0$, and $h_{\alpha+1} = \epsilon_{h_\alpha}$, whose smallest fixpoint will be an infinitely descending tower of ϵ 's, $\epsilon_{\epsilon_{\dots}}$, with ω nested subscripts.

We could call this ordinal γ_0 , and build $\gamma_1, \gamma_2, \dots, \gamma_{\gamma_0}, \dots, \gamma_{\gamma_{\gamma_0}}, \dots$, up to the smallest fixpoint of the function $h_0 = \gamma_0, h_{\alpha+1} = \gamma_{h_\alpha}$. However, the Greek alphabet will sooner or later fit badly in this picture and turn out to be a very unwieldy choice of notation at some point. A more systematic approach to this notational problem is given by the *Veblen hierarchy* of fixpoint-counting functions, but this is a longer story than it is worth explaining here.

A *normal function* is a class function $f : \text{Ord} \rightarrow \text{Ord}$ such that:

- f is strictly monotone: if $\alpha < \beta$, then $f(\alpha) < f(\beta)$;
- f is *continuous*: for every limit ordinal λ , $f(\lambda) = \sup\{f(\alpha) : \alpha < \lambda\}$.

EXERCISE PO.29 \square

- Prove that the following operations are normal functions: $x \mapsto \alpha + x, x \mapsto \alpha \cdot x, x \mapsto \beta^x$ for every $\alpha, \beta \in \text{Ord}, \beta > 1$.

- Prove that the composition of two normal functions is normal.
- Prove the **Veblen fixpoint lemma** for normal functions: if $f : \text{Ord} \rightarrow \text{Ord}$ is normal, then it has a fixpoint, i.e., for some ordinal α , $f(\alpha) = \alpha$, and in fact the class $\text{Fix}(f) = \{\alpha \in \text{Ord} \mid f(\alpha) = \alpha\}$ is unbounded and *closed*: if $A \subseteq \text{Fix}(f)$ then $\sup A$ exists in $\text{Fix}(f)$.

Given a poset (P, \leq) and a subset $A \subseteq P$, consider the set $\uparrow A = \{x \in P \mid \forall a \in A, a \leq x\}$ of all upper bounds of A , and the set $\downarrow A = \{x \in P \mid \forall a \in A, x \leq a\}$ of all lower bounds of A . The **Isbell envelope** $\text{Isb}(P)$ of P consists of the set all subsets $A \subseteq P$ with the property that $\downarrow(\uparrow A) = A$, ordered by inclusion. An element $A \in \text{Isb}(P)$ will be called *Isbell-closed*.

EXERCISE PO.30 \square Show that $\downarrow(-) \dashv \uparrow(-)$ is a Galois connection $(2^P, \supseteq) \rightarrow (2^P, \subseteq)$; this means that $U \subseteq \downarrow V$ if and only if $V \subseteq \uparrow U$; deduce that the following inequalities hold:

- for every $A \subseteq P$, $A \subseteq \downarrow(\uparrow A)$;
- $\uparrow(\downarrow(\uparrow A)) = \uparrow A$.

EXERCISE PO.31 \square Define a monotone function $\iota : P \rightarrow \text{Isb}(P)$ sending $x \in P$ to the principal ideal $\downarrow x$ (show that this is well-defined, i.e. that $\downarrow x$ is Isbell-closed, and that ι is monotone).

EXERCISE PO.32 \square Show that the Isbell envelope of P is order-isomorphic to the poset of **cuts** in P : a cut in P is a pair (A, B) of subsets of P such that $\uparrow A = B$ and $\downarrow B = A$. (Show that, if (A, B) is a cut, A is Isbell-closed, and conversely if A is Isbell-closed...), and for two cuts $(A, B), (A', B')$ we define $(A, B) \leq (A', B')$ if and only if $A \subseteq A'$. (Show that this is in turn equivalent to $B' \subseteq B$).

EXERCISE PO.33 \square Show that $\text{Isb}(P)$ is a complete lattice equipped with an injective monotone function $P \hookrightarrow \text{Isb}(P)$; show the *universal property* of the Isbell envelope: given an injective monotone function $\eta : P \rightarrow D$, with codomain a complete lattice, there exists a monotone embedding $\bar{\eta} : \text{Isb}(P) \hookrightarrow D$ such that the triangle

$$\begin{array}{ccc}
 P & \xrightarrow{\eta} & D \\
 \searrow \iota & & \nearrow \bar{\eta} \\
 & \text{Isb}(P) &
 \end{array}
 \tag{2.14}$$

is commutative.

EXERCISE PO.34 \square (Difficult.) Assume the poset (P, \leq) has a certain Hasse diagram H_P ; devise a method to find the Hasse diagram of $\text{Isb}(P)$.

DEFINITION 2.11. A **locally finite** poset P is a poset (P, \leq) such that each element $[a, b] \in \mathbf{I}(P)$ of the interval poset of P is a finite set.

EXERCISE PO.35 \square Let P be a locally finite poset. Define the **incidence algebra** JP of a poset P to be the set of all functions $f : \mathbf{I}(P) \rightarrow \mathbf{C}$, assigning to each interval $[a, b] \in \mathbf{I}(P)$ a complex number $f_{[a,b]}$. On this set we define the **convolution** operation as follows:

$$f * g : [a, b] \mapsto \sum_{a \leq x \leq b} f_{[a,x]} g_{[x,b]}. \quad (2.15)$$

Show that the incidence algebra JP of P is a monoid, when it is equipped with the convolution product and when the identity element is the ‘delta’ function $\delta : [a, b] \mapsto 1$ if $a = b$ and 0 otherwise.

EXERCISE PO.36 \square The *zeta function* of an incidence algebra JP is the constant function $\zeta(a, b) = 1$ for every nonempty interval $[a, b]$. Prove that ζ is an invertible element of JP (with respect to convolution); the inverse is the *Möbius function* of P , defined as follows:

$$\mu(x, y) = \begin{cases} 1 & x = y \\ \sum_{x \leq z < y} \mu(x, z) & x < y \\ 0 & \text{otherwise.} \end{cases} \quad (2.16)$$

Prove that indeed $\zeta * \mu = \delta = \mu * \zeta$.

EXERCISE PO.37 \square Find an explicit expression for $\zeta * \zeta$. What about $\zeta * \zeta * \zeta$? Generalise.

Consider the poset (\mathbf{N}, \mid) of positive integers partially ordered by divisibility. The *reduced* incidence algebra consists of functions $f(a, b)$ that are invariant under multiplication, i.e. $f(ka, kb) = f(a, b)$ for all $k \geq 1$.

For a multiplicative invariant function, the value $f(a, b)$ depends only on b/a , so a natural basis consists of *invariant delta functions* δ_n defined by $\delta_n(a, b) = 1$ if $b/a = n$ and 0 otherwise: any invariant function can be written $f = \sum_{n \geq 1} f(1, n) \cdot \delta_n$.

EXERCISE PO.38 \square Show that the convolution of two invariant delta functions is still an invariant delta function.

To every element of the reduced incidence algebra we associate the *Dirichlet series* $\kappa_f := \sum_{n \geq 1} \frac{f(1, n)}{n^s}$.

The zeta function ζ belongs to the reduced incidence algebra and its associated Dirichlet series corresponds to the so-called *Riemann zeta function*

$$\zeta(s) = \sum_{n \geq 1} \frac{1}{n^s}. \quad (2.17)$$

EXERCISE PO.39 \square

- Show that $0 = \zeta(0) = \zeta(-2) = \zeta(-4) = \dots = \zeta(-2k)$ for all $k \geq 0$.
- (difficult) Show that all other zeros of ζ belong to the set $\frac{1}{2} + i\mathbf{R} = \{\frac{1}{2} + it \mid t \in \mathbf{R}\}$.

3. Semigroups, monoids

EXERCISE SM.1 \square Observe that if $a, b \in \mathbf{R}^+$ are strictly positive real numbers, the quotient a/b is still strictly positive. Is the set $(\mathbf{R}^+, /)$ endowed with the operation $(a, b) \mapsto a/b$ a semigroup? Is it commutative?

EXERCISE SM.2 \square Define a binary operation on the set of natural numbers as follows:

$$a \circ b := a + b + ab \quad (3.1)$$

Show that (\mathbf{N}, \circ) is a commutative semigroup.

EXERCISE SM.3 \square Let (S, \cdot) be a monoid, and X any set. Define a binary operation $*$ on the set S^X of all functions $X \rightarrow S$ as follows:

$$(f, g) \mapsto f * g : x \mapsto f(x) \cdot g(x) \quad (3.2)$$

Show that $(S^X, *)$ is a monoid, and that $(S^X, *)$ is commutative if (S, \cdot) is. Is the converse implication true (if $(S^X, *)$ is commutative, (S, \cdot) is commutative)?

EXERCISE SM.4 \square Let S be a *finite* set, and consider the monoid S^S of all functions $f : S \rightarrow S$. Prove that f is invertible if and only if it is an injective function, if and only if it is a surjective function. Find a counterexample when S is not finite.

EXERCISE SM.5 \square Consider the set $S = \mathbf{R}^\times \times \mathbf{N}$, where $\mathbf{R}^\times = \mathbf{R} \setminus \{0\}$; define a binary operation on S as follows:

$$(a, n)(b, m) := (ab^n, nm) \quad (3.3)$$

Show that with this definition S is a semigroup. Is it a monoid? Is it commutative?

EXERCISE SM.6 \square Let A^A be the monoid of endofunctions of a set A ; count how many elements there are in A^A if $A = \{1, 2, 3, 4, 5\}$. Let $f(a) = \min\{a^2, 5\}$, and consider the cyclic monoid $N = \langle f \rangle$ generated by f ; how many elements does N have? Which ones?

EXERCISE SM.7 \square Let $\mathbf{Z} \times \mathbf{Z}$ be the set of pairs of integers (m, n) . Define a monoid operation

$$(p, q)(r, s) := (pr - qs, ps + qr) \quad (3.4)$$

prove that it is indeed a monoid operation, and that the function

$$v : \mathbf{Z} \times \mathbf{Z} \rightarrow \mathbf{N} : (p, q) \mapsto p^2 + q^2 \quad (3.5)$$

is a monoid homomorphism.

The relation of *congruence modulo an integer* is arguably to be the most important kind of equivalence relation in all Mathematics. Let \mathbf{Z} be the set of integers $\{\dots, -2, -1, 0, 1, 2, \dots\}$, and fix an integer $n \in \mathbf{Z}_{\geq 1}$. Define an equivalence relation \equiv_n on \mathbf{Z} as follows: $a \equiv_n b$ if and only if $a - b$ is a multiple of n .

EXERCISE SM.8 \square Show that \equiv_n is an equivalence relation, and in fact a *congruence* on the set of integers, namely that if $a \equiv_n b$ then for every $c \in \mathbf{Z}$ $a + c \equiv_n c + b$ and $c \cdot a \equiv_n c \cdot b$.

EXERCISE SM.9 \square Consider the set $\mathbf{Z}/n := \mathbf{Z}/\equiv_n$ of equivalence classes modulo \equiv_n . Show that \mathbf{Z}/n is a finite set, with exactly n elements $\{[0], [1], \dots, [n-1]\}$. Define a *sum* operation on \mathbf{Z}/n ,

$$[a] + [b] = [a + b] \quad (3.6)$$

and a product operation

$$[a] \cdot [b] = [a \cdot b] \quad (3.7)$$

Show that $+$, \cdot are well-defined, and that $(\mathbf{Z}, +)$, (\mathbf{Z}, \cdot) are commutative monoids. Show that moreover $(\mathbf{Z}, +)$ is an Abelian group.

EXERCISE SM.10 \square Let $f : M \rightarrow N$ a surjective monoid homomorphism. Prove that if M is cyclic, so is N .

EXERCISE SM.11 \square Let $\mathbf{R}^+ = \{x \in \mathbf{R} \mid x \geq 0\}$. Prove that $(\mathbf{R}^+, \cdot, 1)$ is a submonoid of $(\mathbf{R}, \cdot, 1)$. Prove that the function $\log_2 : \mathbf{R}^+ \rightarrow \mathbf{R}$ sending x into $\log_2 x$ is a monoid isomorphism.

EXERCISE SM.12 \square Let $f : M \rightarrow M'$ be a monoid homomorphism. Then,

- prove that if N is a submonoid of M , its image $f(N)$ is a submonoid of M' ;
- prove that if N' is a submonoid of M' , then $f^{-1}(N')$ is a submonoid of M .

EXERCISE SM.13 \square Let M be the monoid $(\mathbf{N}, +, 0) \times (\mathbf{N}, +, 0)$, where

$$(a, b) + (c, d) = (a + c, b + d) \quad (3.8)$$

and where the unit is $(0, 0)$. Let $a, b \in \mathbf{N}$ and define the function

$$f_{a,b} : M \rightarrow \mathbf{N} : (x, y) \mapsto a^x b^y \in M. \quad (3.9)$$

Prove that $f_{a,b}$ is a homomorphism $M \rightarrow (\mathbf{N}, \cdot, 1)$. Prove that if $a > 1$ and $b = a^2$, $f_{a,b}$ is not injective. Prove that if a, b are distinct prime numbers, then f is injective.

EXERCISE SM.14 \square Prove that every monoid is isomorphic to a submonoid of (X^X, \circ) for a suitable set X [hint: take X equal to the monoid M , and for each $a \in M$ define the function $f_a : x \mapsto ax$].

EXERCISE SM.15 \square Let $S = \{-1, 0\}$ regarded as a subset of \mathbf{R} . Describe the submonoid $\langle S \rangle$ of (\mathbf{R}, \cdot) . Prove that there exists a unique endomorphism $\varphi : (\mathbf{R}, \cdot) \rightarrow (\mathbf{R}, \cdot)$ with the property that $\varphi(0) = 0$ and $\varphi(a) = -1$ for every $a < 0$. Prove that the image of such φ is exactly S .

EXERCISE SM.16 \square Let X be a set and consider the monoid (PX, \cup) of subsets of X , where the monoid operation is given by union. Let $S = \{\{a\} \mid a \in X\}$ be the subset of PX whose elements are singletons. Describe the submonoid $\langle S \rangle$ in (PX, \cup) .

EXERCISE SM.17 \square In the same notation as above, describe the cyclic submonoid $\langle a \rangle \leq (PX, \cup)$ for a given $a \in X$.

EXERCISE SM.18 \square Recall that a *group* is a monoid $(M, \cdot, 1_M)$ such that every element is invertible (this means that every element $x \in M$ has an *inverse* x^{-1} , such that $x \cdot x^{-1} = 1_M = x^{-1} \cdot x$). Prove that a semigroup (S, \cdot) is a group if and only if for every $a \in S$ we have the equalities

$$aS := \{as \mid s \in S\} = S = Sa := \{sa \mid s \in S\}. \quad (3.10)$$

The notation aS is a particular case, and shorthand, for the following more general definition: let A, B be subsets of a semigroup S ; denote as $A \cdot B$ or just AB the set $\{a \cdot b \mid a \in A, b \in B\}$; in particular, aB is a shorthand for $\{a\} \cdot B$.

This notation clearly extends to the ‘product’ of n subsets $A_1, \dots, A_n \subseteq S$:

$$A_1 \cdots A_n := \{a_1 \dots a_n \mid a_i \in A_i, 1 \leq i \leq n\}. \quad (3.11)$$

EXERCISE SM.19 \square Let (S, \cdot) be a finite semigroup satisfying both the left and right cancellation laws: if $xy = xz$ then $y = z$, and if $yx = zx$, then $y = z$. Prove that S is a group.

EXERCISE SM.20 \square An idempotent element of a monoid $(M, \cdot, 1_M)$ is an element $e \in M$ such that $e \cdot e = e$. Prove that a finite monoid is a group if and only if it has a unique idempotent element (and that element is the identity 1_M).

EXERCISE SM.21 \square Let (S, \cdot) be a finite semigroup, and $a \in S$; show that there exists a smallest positive integer $n \in \mathbf{N}$, called the *index* of a , such that $a^n = a^{n+d}$ for some $d > 0$. The smallest possible choice of d is called the *period* of a . Show that $s^c = s^{n+rd}$ for every $r \geq 0$; show that $s^p = s^q$ if and only if $p = q < n$, or $p, q \geq n$ and $p \equiv_d q$.

EXERCISE SM.22 \square In the same notation as above, S be a finite semigroup, say of cardinality n . Let S^n be the set of products $s_1 \dots s_n$ of n elements of S . Prove the *pumping lemma*: the set S^n equals the set

$$S \cdot E_S \cdot S = \{s \cdot e \cdot s' \mid s, e, s' \in S, ee = e\} \quad (3.12)$$

(in words: every element of S^n can be written as a product ses' , where $s, e, s' \in S$ and e is an idempotent).

EXERCISE SM.23 \square Let S, T be finite semigroups; let $f : S \rightarrow T$ be a surjective homomorphism; prove that the image of the set of idempotents of S , $E_S = \{e \in S \mid ee = e\}$ under f equals the set E_T of idempotents of T .

EXERCISE SM.24 \square Let (M, \star) be a monoid; define $x \sim y$ if there exists $n \in \mathbf{N}$ such that $x^n = y^n$ in M . Show that \sim is an equivalence relation on M , and that if M is commutative and $x \sim y$, then $a \star x \sim a \star y$ for all $a \in M$. Describe the equivalence class of 1_M .

EXERCISE SM.25 \square Let A be a set, and (A^A, \circ) the monoid of all functions $f : A \rightarrow A$. Fix a subset $B \subseteq A$. Show that the set

$$S(B) := \{f : A \rightarrow A \mid f(B) \subseteq B\} \quad (3.13)$$

is a submonoid of A^A . Prove that the function $\psi : S(B) \rightarrow B^B$ sending $f : A \rightarrow A$ to its restriction to B , $f|_B : B \rightarrow B$ is a surjective monoid homomorphism. Define an equivalence relation on $S(B)$ saying that $f \sim g$ if $f(b) = g(b)$ for every $b \in B$; prove that \sim is precisely the equivalence relation induced by ψ .

EXERCISE SM.26 \square Define an equivalence relation on the set $\mathbf{N} \times \mathbf{N}$ as follows:

$$(p, q) \sim (r, s) \iff p + s = q + r. \quad (3.14)$$

Define the sum on $\mathbf{N} \times \mathbf{N}$ as follows:

$$(p, q) + (r, s) = (p + r, q + s); \quad (3.15)$$

define the product

$$(p, q)(r, s) = (pr + qs, ps + qr). \quad (3.16)$$

Prove that

- \sim is an equivalence relation, and in fact a congruence on $\mathbf{N} \times \mathbf{N}$;
- the set $(\mathbf{N} \times \mathbf{N}, +)$ is a monoid;
- the set $(\mathbf{N} \times \mathbf{N}, \cdot)$ is a monoid;
- the *distributive property* holds for every $(a, b), (p, q), (r, s) \in \mathbf{N} \times \mathbf{N}$:

$$(a, b) \cdot ((p, q) + (r, s)) = (a, b)(p, q) + (a, b)(r, s). \quad (3.17)$$

- \sim is the equivalence relation generated by $f : \mathbf{N} \times \mathbf{N} \rightarrow \mathbf{Z} : (a, b) \mapsto a - b$.

EXERCISE SM.27 \square In the same notation of the previous exercise, show that $f : \mathbf{N} \times \mathbf{N} \rightarrow \mathbf{Z} : (a, b) \mapsto a - b$ induces an isomorphism $\mathbf{Z} \cong (\mathbf{N} \times \mathbf{N})/\sim$.

An important role in monoid theory is played by the notion of *ideal* and the closely related notion of *Green's relations*. In the following, let M be a finite monoid.

A *left ideal* (respectively, *right ideal*) of M is a nonempty subset $I \subseteq M$ such that $MI \subseteq I$ (respectively, $IM \subseteq I$). A *two-sided ideal*, or just *ideal* is a nonempty subset $I \subseteq M$ such that $MIM \subseteq I$. We write $I \trianglelefteq M$ to denote the fact that $I \subseteq M$ and I is an ideal. The set $(\text{Idl}(M), \trianglelefteq)$ is a poset.

EXERCISE SM.28 \square If $I \trianglelefteq M$ is an ideal of a finite monoid M , prove that I contains at least an idempotent element $e = ee$.

EXERCISE SM.29 \square Let M be a finite monoid; then it has a finite number of ideals I_1, \dots, I_r ; prove that the product $I_1 I_2 \dots I_r$ is still an ideal of M , and that it is contained in every other ideal. Consequently, each finite monoid M has a unique *minimal* ideal I_{\min} .

Let $m \in M$ be an element of a finite monoid M . The ideals Mm, mM and MmM are respectively called the principal left, principal right, and principal two-sided ideal generated by m .

EXERCISE SM.30 \square Define

$$I[m] := \{x \in M \mid m \notin MxM\}; \quad (3.18)$$

prove that if $I[m] \neq \emptyset$, then it is an ideal of M ; prove that $I[m]$ is empty if and only if it is contained in I_{\min} .

EXERCISE SM.31 \square Let M be a finite monoid. Prove that the following defines three equivalence relations on M :

- $x \equiv y \pmod{j}$ if and only if $MxM = MyM$;
- $x \equiv y \pmod{l}$ if and only if $Mx = My$;
- $x \equiv y \pmod{r}$ if and only if $xM = yM$.

(Writing ‘mod j ’ is just a slick shorthand: take for example the j relation; we write $x \equiv y \pmod{j}$ or $x \equiv_j y$ to denote that $(x, y) \in j \subseteq M \times M$.) These are called respectively the *two-sided*, *left* and *right* Green relations on M . We say that x, y are *two-sided Green equivalent*, or j -equivalent, if $x \equiv y \pmod{j}$. Similarly, we say that x, y are left Green equivalent, or l -equivalent, if $x \equiv y \pmod{l}$, and right Green equivalent, or r -equivalent, if $x \equiv y \pmod{r}$.

EXERCISE SM.32 \square Let $a, b, c \in M$ be elements of a finite monoid. Prove that if $a \equiv_r b$, then $ca \equiv_r cb$ and if $a \equiv_l b$, then $ac \equiv_l bc$; is it also true that if $a \equiv_r b$, then $ac \equiv_l bc$?

EXERCISE SM.33 \square Let M be a finite monoid. Prove that the relation j is the join of r, l in the poset of equivalence relations on M .

EXERCISE SM.34 \square Let k be a field and $n \geq 1$ an integer; consider the monoid $M_n(k)$ of $n \times n$ matrices with entries in k . Show that two matrices $A, B \in M_n(k)$ are j -equivalent if and only if they have the same rank.

EXERCISE SM.35 \square Let A^A be the monoid of endofunctions of a finite set A ; show that $f, g \in A^A$ are j -equivalent if and only if their images have the same cardinality; $f, g \in A^A$ are l -equivalent if and only if the associated equivalence relations R_f, R_g are equal; $f, g \in A^A$ are r -equivalent if and only if they have the same image.

EXERCISE SM.36 \square A monoid M is called *r -trivial* if $mM = nM$ implies $m = n$, or in other words, if the relation r reduces to the identity Δ_M ; similarly we define a l -trivial monoid, and a j -trivial monoid M (in the latter case, we mean that if $MmM = MnM$, then $m = n$). Prove that a monoid is j -trivial if and only if it is both r -trivial and l -trivial.

EXERCISE SM.37 \square Implement *Light associativity test* in your favourite programming language: given a finite set S , a sufficient condition for a binary operation $\star : S \times S \rightarrow S$ to be associative is that given any element $y \in S$, the following two tables coincide:

- the ‘matrix’ $L_1(y)$ constructed from an enumeration $S = \{x_1, \dots, x_n\}$ whose entry (i, j) is $x_i \star (y \star x_j) \in S$;

- the matrix $L_2(y)$, constructed from the same enumeration, whose entry (i, j) is $(x_i \star y) \star x_j$.

The operation \star is associative if and only if for every $y \in S$, $L_1(y) = L_2(y)$.

DEFINITION 3.1. A **partially ordered monoid** is a monoid $(M, \cdot, 1)$ equipped with a partial order \leq that is *compatible* with the monoid operation; this means that for each $a, b, c \in M$, if $b \leq c$ then $ab \leq ac$ and $ba \leq ca$.

A partially ordered monoid $(M, \leq, \cdot, 1)$ is usually called a *po-monoid*.

EXERCISE SM.38 \square Prove that if $b \leq c$ in a po-monoid and b, c are invertible then $c^{-1} \leq b^{-1}$, so that in a po-group G (i.e. a po-monoid that in addition is a group), the inversion map $(-)^{-1} : G^{\text{op}} \rightarrow G$ is monotone.

EXERCISE SM.39 \square Let $(G, \leq, \cdot, 1)$ be a po-group; the **positive cone** $G^+ \subseteq G$ is the set $\{x \in G \mid 1 \leq x\}$. Surprisingly, given a group and a positive cone, we can reconstruct the order relation on G : a group equipped with a positive cone is a pair (G, H) where G is a group and $H \subseteq G$ is such that

- $1 \in H$;
- if $a, b \in H$, then $ab \in H$;
- if $a \in H$ and $x \in G$, then $x^{-1}ax \in H$;
- if $a, a^{-1} \in G$, then $a = 1$.

Prove that if (G, H) is a group with a positive cone, the relation $a \leq b$ iff $ba^{-1} \in H$ defines a partial order on G that renders $(G, \leq, \cdot, 1)$ a po-group, and H coincides with the positive cone of $(G, \leq, \cdot, 1)$.

EXERCISE SM.40 \square Let P, Q be two posets. Define the **product order** on the Cartesian product $P \times Q$ by $(x, y) \leq (x', y')$ if and only if $x \leq x'$ and $y \leq y'$. Show that the product order $G \times H$ of two po-groups is again a po-group, with the usual group structure on the set $G \times H$.

EXERCISE SM.41 \square Let (P, \leq) be a poset, and let \sim be an equivalence relation on P . One says that \sim is *compatible* with the order relation if $x \leq y$, $x \sim x'$ and $y \sim y'$ imply $x' \leq y'$ or $x' \sim y'$. When this happens the quotient set P/\sim carries a relation $[x] \leq [y]$ if and only if $x \leq y$ or $x \sim y$; prove that this is a partial order. Prove that the projection to the quotient $P \rightarrow P/\sim$ is a monotone map.

EXERCISE SM.42 \square Let P, Q be two posets. Define the **lexicographic order** on the Cartesian product $P \times Q$ by $(x, y) \leq_{\text{lex}} (x', y')$ if and only if either $x < x'$, or $x = x'$ and $y \leq y'$. Denote $P \times_{\text{lex}} Q$ the set $P \times Q$ equipped with the lexicographic order.

Show that the lexicographic order on $\mathbf{Z} \times \mathbf{Z}$ makes it a po-group. Is it true that more generally, the lexicographic product $\mathbf{Z} \times_{\text{lex}} P$ of two po-groups is a po-group?

EXERCISE SM.43 \square Prove that any group G can be seen as a po-group with the trivial order relation $g \leq h$ if and only if $g = h$; prove that if G is a finite group, the trivial order ($x \leq y$ iff $x = y$) is the only possible po-group structure.

EXERCISE SM.44 \square Let (P, \leq) be a poset, and Q^P the set of monotone mappings $P \rightarrow Q$; defines the *standard order* on Q^P by saying that $f \leq g$ when for all $p \in P$ $f(p) \leq g(p)$. Prove that this defines a partial order on Q^P .

We will consider the standard order on Q^P in the particular case $P = Q$ and sometimes restricted to the subset of *invertible* monotone mappings $P \rightarrow P$ (i.e. to the set $\text{Aut}(P)$ of ‘monotone automorphisms’ of P). Prove that P^P is a po-monoid, and $\text{Aut}(P) \subseteq P^P$ is a po-group when both sets are equipped with the standard order.

DEFINITION 3.2. Let (G, \leq) be a po-group. A **G -poset**, or a poset *equipped* with a G -action, is a partially ordered set (P, \leq) endowed with a po-group homomorphism $a : G \rightarrow \text{Aut}(P)$ to the group of order isomorphisms of P with its standard po-group structure.

If a po-group G acts on a poset P , the action of the function $a(g, p)$ is usually denoted as an infix dot $g.p$.

EXERCISE SM.45 \square Show that equivalently, a G -poset is a partially ordered set P together with a group action $G \times P \rightarrow P$ which is a monotone map, where on $G \times P$ one puts the product order.

EXERCISE SM.46 \square Prove that the poset (\mathbf{Z}, \leq) of integers with their usual order is a \mathbf{Z} -poset with the action given by the usual sum of integers. More generally, every po-group is a G -poset, where the action $G \times G \rightarrow G$ is the group operation.

Show that the poset (\mathbf{R}, \leq) of real numbers with their usual order is a \mathbf{Z} -poset for the action given by the sum of real numbers with integers (seen as a subring of real numbers).

EXERCISE SM.47 \square Let G be a po-group acting on a poset P . A **G -fixed point** for a G -poset P is an element $p \in P$ kept fixed by all the elements of G under the G -action. Prove that if P has a top element \top_P or bottom element \perp_P , then they both are G -fixed points. Deduce that one can always extend the action of G on P on a larger poset $P_\diamond = P \cup \{+\infty, -\infty\}$ where $-\infty \leq x \leq +\infty$ for all $x \in P$.

EXERCISE SM.48 \square Given a po-group G and two G -posets P, Q we say that a monotone map $f : P \rightarrow Q$ is G -equivariant if for all $g \in G$ one has $f(g.p) = g.f(p)$. More formally, let G be a po-group, and P, Q be two G -posets, respectively with actions $a_P : G \times P \rightarrow P$ and $a_Q : G \times Q \rightarrow Q$. Then, a monotone map $f : P \rightarrow Q$ is G -equivariant if $a_Q(g, f.p) = f(a_P(g, p))$

EXERCISE SM.49 \square An equivalence relation \sim on a G -poset P is said to be *compatible* with the G -action if $x \sim y$ implies $g \cdot x \sim g \cdot y$ for any g in G . If \sim is compatible both

with the order and with the G -action then the quotient set P/\sim is naturally a G -poset with the G -action $g \cdot [x] = [g \cdot x]$. Moreover the projection to the quotient is a morphism of G -posets.

EXERCISE SM.50 \square Consider the group \mathbf{Z} of integers as a po-group acting on itself by addition. Prove that there exists a bijection

$$P \cong \{\text{equivariant maps } \varphi : (\mathbf{Z}, \leq) \rightarrow (P, \leq)\} \quad (3.19)$$

arguing as follows: the choice of an element x in a \mathbf{Z} -poset P is equivalent to the datum of a \mathbf{Z} -equivariant morphism $\varphi_x : (\mathbf{Z}, \leq) \rightarrow (P, \leq)$. (Bonus point if you feel like it: the element x is a \mathbf{Z} -fixed point if and only if the corresponding morphism φ factors \mathbf{Z} -equivariantly through $(*, \leq)$.)

EXERCISE SM.51 \square Consider the group \mathbf{Z} of integers as a po-group acting on itself by addition. Prove that a \mathbf{Z} -equivariant monotone map $\varphi : (\mathbf{Z}, \leq) \rightarrow (P, \leq)$ is either injective or a constant map; as a consequence, if P is a finite poset, there are no nontrivial monotone \mathbf{Z} -actions.

EXERCISE SM.52 \square Show that every equivariant map $\varphi : P \rightarrow Q$ where Q is bounded (i.e. it admits a top and a bottom element) extends to a unique $\bar{\varphi} : P_\diamond \rightarrow Q$ defining $\bar{\varphi}(\infty) = \top_Q$ and $\bar{\varphi}(-\infty) = \perp_Q$ (see [Exercise 47](#) for the notation P_\diamond).

DEFINITION 3.3. A **quantale** is an algebraic lattice (Q, \wedge, \vee) where every subset $S \subseteq Q$ has both an infimum and a supremum (so, in particular, Q has a top element \top_Q and a bottom element \perp_Q), and equipped with a semigroup structure $* : Q \times Q \rightarrow Q$ such that both maps $x * _$ and $_ * y$ preserve arbitrary joins:

$$x * \left(\bigvee_{i \in I} y_i \right) = \bigvee_{i \in I} (x * y_i) \quad \left(\bigvee_{i \in I} x_i \right) * y = \bigvee_{i \in I} (x_i * y). \quad (3.20)$$

The quantale is *unital* when $*$ has an identity element; *commutative* when $*$ is commutative.

EXERCISE SM.53 \square Show that in a quantale the assignment $(a, b) \mapsto \bigvee \{q \in Q \mid a * q \leq b\}$ defines a binary operation $_ \rightarrow _ : Q \times Q \rightarrow Q$ with the property that

$$a * x \leq b \iff x \leq (a \rightarrow b). \quad (3.21)$$

Similarly, the assignment $(a, b) \mapsto \bigvee \{q \in Q \mid q * a \leq b\}$ defines a binary operation $_ \leftarrow _ : Q \times Q \rightarrow Q$ with the property that

$$x * a \leq b \iff x \leq (a \leftarrow b). \quad (3.22)$$

EXERCISE SM.54 \square Show that a two-element set $\{0, 1\}$ has a unique quantale structure if we define $a * b = a \cdot b$ is 1 if and only if both a, b are 1, and 0 otherwise (define $a \Rightarrow b$ in such a way that (3.21) is true, and prove it's the unique possible choice).

EXERCISE SM.55 \square Let $S = \{0, \epsilon, 1\}$ be a three-element set; prove that the following definitions, packaged in a multiplication table, for $*$ and \Rightarrow equip S with a quantale structure, and show that it is the unique one.

| | | | |
|------------|-----|------------|-----|
| $*$ | 0 | ϵ | 1 |
| 0 | 0 | 0 | 0 |
| ϵ | 0 | ϵ | 1 |
| 1 | 0 | 1 | 1 |

| | | | |
|------------|-----|------------|-----|
| $*$ | 0 | ϵ | 1 |
| 0 | 1 | 0 | 0 |
| ϵ | 1 | ϵ | 0 |
| 1 | 1 | 1 | 1 |

A **Heyting algebra** H is a quantale where the $_ * _$ operation coincides with the binary meet.

EXERCISE SM.56 \square Prove that equivalently, a Heyting algebra consists of a meet-semilattice with bottom element, equipped with a binary operation $_ \rightarrow _$ satisfying

$$a \wedge b \leq c \iff c \leq a \rightarrow b \quad (3.23)$$

EXERCISE SM.57 \square Prove that an Heyting algebra H admits a *pseudo-complement* operation: for every $a \in H$ there exists a maximum element $\complement a$ with the property that $a \wedge \complement a = 0$.

EXERCISE SM.58 \square Prove that if H is a Heyting algebra,

- $\complement a$ is uniquely determined by the property that $b \wedge a = 0$ if and only if $b \leq \complement a$;
- if $x \leq y$, then $\complement y \leq \complement x$;
- the assignment $x \mapsto \complement \complement x$ is monotone, and $x \leq \complement \complement x$;
- for each $x \in H$, $\complement \complement \complement x = \complement x$;
- for each $a, b \in H$, $\complement \complement (a \wedge b) = \complement \complement a \wedge \complement \complement b$;
- if $a, b \in H$ have pseudo-complements $\complement a, \complement b$, then $\complement (a \vee b) = \complement a \wedge \complement b$ and actually more generally the *first de Morgan law* holds:

$$\complement (\bigvee_{i \in I} a_i) = \bigwedge_{i \in I} \complement a_i \quad (3.24)$$

- for every $a \in H$, $\complement (a \vee \complement a) = 0_H$ (the bottom element of H).

DEFINITION 3.4. A **Boole algebra** is a Heyting algebra that is *complemented*, i.e. the pseudo-complement of a also satisfies the equation $a \vee \complement a = 1$.

EXERCISE SM.59 \square Prove that a Heyting algebra is a Boole algebra if and only if for every $a \in H$, $\complement \complement a = a$.

EXERCISE SM.60 \square Prove that in a Boole algebra the *second de Morgan law* holds together with the first:

$$\complement (\bigwedge_{i \in I} a_i) = \bigvee_{i \in I} \complement a_i \quad (3.25)$$

EXERCISE SM.61 \square Let \mathfrak{f} be a proper filter in a Boolean algebra B . Prove that the following statements are equivalent.

- \mathfrak{f} is maximal;

- \mathfrak{f} is prime;
- for every $a \in B$ either a or $\complement a \in \mathfrak{f}$.

(This generalises [Exercise 62](#)).

Every Heyting algebra H contains a maximal Boole algebra:

EXERCISE SM.62 \square Prove that $BH = \{a \in H \mid \complement \complement a = a\} \subseteq H$ is a Boole algebra.

EXERCISE SM.63 \square Prove that the following conditions on a Heyting algebra L are equivalent:

- the second de Morgan law (3.25) holds;
- for each $a \in H$, $\complement \complement a \vee \complement a = 1_H$;
- every element of BH , as defined above, has a complement in H ;
- the identity $\complement \complement (a \vee b) = \complement \complement a \vee \complement \complement b$ holds for all $a, b \in H$;
- BH is a sublattice of H .

EXERCISE SM.64 \square Show that any po-group G (written multiplicatively) admitting all sups and inf is a quantale if we define $a * b$ as the group multiplication ab , and $a \rightarrow b = a^{-1}$, $a \leftarrow b = ba^{-1}$.

4. Linear Algebra, done hard

Recall the definition of ring, ring homomorphism, ideal, field and whatever amount of Linear Algebra you have known from any Abstract Algebra book that you can get your hands on.

EXERCISE LA.1 \square Let R be a unital ring; prove that there exists a unique ring homomorphism $\eta_R : \mathbf{Z} \rightarrow R$, called the *characteristic* homomorphism.

Find the kernel of η_R when $R = \mathbf{Z}$, \mathbf{Q} and when $R = \mathbf{Z}/n\mathbf{Z}$ is the ring of integers modulo n . Prove that $\ker \eta_R$ is an ideal of \mathbf{Z} generated by a single element $q \in \mathbf{Z}$.

EXERCISE LA.2 \square Given a unital ring R , its characteristic is the minimal generator q of $\ker \eta_R$.

Prove that a field has characteristic either zero or a prime number.

A *ring extension* of a commutative ring R is a commutative ring E of which R is a subring. In other words, a ring extension is an injective ring homomorphism $R \rightarrow E$.

EXERCISE LA.3 \square Let F be a field; prove that a ring homomorphism $F \rightarrow E$ is either the constant at zero homomorphism, or a ring extension (in particular, if we insist on a ring homomorphism to preserve the multiplicative unit, all ring homomorphisms from a field are ring extensions).

We denote a field extension E of a field F as $E|F$.

EXERCISE LA.4 \square Prove that an extension $E|F$ makes E a vector space over F : the extension is called *finite* if E has finite dimension over F ; note that being finite for E depends on F , by finding a field F_1 such that $E|F_1$ is finite, and an extension $F_1|F_2$ which is not finite, so that $E|F_2$ is not finite.

EXERCISE LA.5 \square Prove that there exists a polynomial p with real coefficients, of degree at most 3, such that $p(x+1) - p(x) = x^2$. Determine $p(x)$. Use this result to compute the sum $\sum_{k=1}^n k^2$ of the squares of the first n positive integers.

EXERCISE LA.6 \square Consider the vector space $L = \mathbf{R}[X]_{<n}$ whose elements are polynomials of degree at most $n-1$; prove that

- the standard basis: i.e. the set $\{1, x, \dots, x^{n-1}\}$ is a basis of L ; the coordinates of a vector in this basis are its coefficients.
- the Taylor basis: for $a \in \mathbf{R}$, the set $\{1, x-a, (x-a)^2, \dots, (x-a)^{n-1}\}$ is a basis of L ; the coordinates of a vector $f \in L$ in this basis are given by its subsequent derivatives evaluated in a : $\{f(a), f'(a), \frac{f''(a)}{2}, \dots, \frac{f^{(n-1)}(a)}{(n-1)!}\}$. Is this still true if instead of polynomials with real coefficients, we take polynomials with coefficients in a finite field?
- interpolation basis: if $a_1, \dots, a_n \in \mathbf{R}$ are pairwise distinct elements of L , define

$$g_i(x) := \prod_{j \neq i} \frac{x - a_j}{a_i - a_j}. \quad (4.1)$$

Then $\{g_1, \dots, g_n\}$ is a basis of L ; the coordinates of a vector $f \in L$ in this basis are given by the values $\{f(a_1), \dots, f(a_n)\}$.

EXERCISE LA.7 \square Let $W \leq V$ be a subspace inclusion; the *quotient* vector space V/W consists of the vector space of equivalence classes of vectors of V by the relation

$$v \sim_W v' \iff v - v' \in W. \quad (4.2)$$

Prove that \sim_W is an equivalence relation, and in fact a congruence ($v \sim_W v'$ implies $u+v \sim_W u+v'$ for all $u \in V$); prove that the set of \sim_W -equivalence classes becomes a vector space if we define the sum $[v] + [v']$ as $[v+v']$, and the scalar multiplication $a[v]$ as $[av]$.

EXERCISE LA.8 \square Find a geometric interpretation for the following quotient vector spaces:

- $\mathbf{R}^2 / \langle (0, 1) \rangle$;
- $\mathbf{R}^3 / \langle (1, 0, 0), (0, 1, 1) \rangle$.

(Hint: represent the elements of the quotient space V/W , i.e. the \sim_W -equivalence classes, as $[v] = v + W = \{v + w \mid w \in W\}$, i.e. as subsets of V resulting as translations of W by vectors of v ; observe that in no case $v + W$ is a vector subspace of V : why? Show that for each $v \in V$, $[v]$ equals $[v_\perp]$, where v_\perp is a vector perpendicular to all vectors of W , in the sense that the scalar product $v \cdot w$ is zero for every $w \in W$; deduce that V/W can be identified with the set of such perpendicular vectors.)

EXERCISE LA.9 \square Let $f : V \rightarrow W$ be a linear maps between F -vector spaces; the *cokernel* of f is the quotient space $W/\text{im } f$, where two vectors $w, w' \in W$ are identified if and only if their difference lies in the image of f ;

- prove that f is surjective if and only if $\text{coker } f$ is the zero vector space;
- what is the cokernel of the zero map $0 : V \rightarrow W$? What is the cokernel of the inclusion of $\langle v \rangle$ in \mathbf{R}^3 as v varies through the vectors of \mathbf{R}^3 ?

DEFINITION 4.1. Let V be a F -vector space. Define the vector space V^* (called the *dual* of V) as the set $\text{hom}(V, F)$ of all F -linear homomorphisms from V to F , equipped with the obvious vector space structure given by $(\alpha + \beta)(v) = \alpha(v) + \beta(v)$ and $(t \cdot \alpha)(v) = t \cdot \alpha(v)$ for every $t \in F, v \in V, \alpha, \beta : V \rightarrow F$.

EXERCISE LA.10 \square Show that this really defines a structure of F -vector space on $\text{hom}(V, F)$.

EXERCISE LA.11 \square Let V be a finite-dimensional F -vector space. Let $\mathcal{V} = \{v_1, \dots, v_n\}$ be a basis of V , and let $\mathcal{V}^* = \{\alpha_1, \dots, \alpha_n\}$ be the set of vectors in V^* defined as follows: $\alpha_i(v_j) = 1$ if $i = j$ and 0 otherwise. Show that \mathcal{V}^* is a basis of V^* , called the *dual basis* associated to \mathcal{V} .

EXERCISE LA.12 \square Show that to each F -linear map $f : V \rightarrow W$ of vector spaces one can associate the *transposed* map $f^* : W^* \rightarrow V^*$ sending an element $\alpha \in W^*$ to the linear map

$$\alpha \circ f : V \xrightarrow{f} W \xrightarrow{\alpha} F \quad (4.3)$$

EXERCISE LA.13 \square Show that there exists a function

$$V^* \times V \xrightarrow{\circ} F \quad (4.4)$$

called *canonical duality* which enjoys the following properties:

- bilinearity:
- nondegenerate:

Let S be a subset of a vector space V ; we define the *orthogonal* of S in V^* as the subspace of V^* as

$$S^\perp := \{\alpha \in V^* \mid \alpha \circ v = 0, \forall v \in S\} \quad (4.5)$$

Similarly, for a subset T of V^* we define the subspace T^\perp of V as

$$T^\perp := \{v \in V \mid \alpha \circ v = 0, \forall \alpha \in T\} \quad (4.6)$$

EXERCISE LA.14 \square Prove that

- if $A \subseteq B$ then $B^\perp \subseteq A^\perp$;
- for every $S \subseteq V, S^{\perp\perp} = \langle S \rangle$ is the subspace generated by S (similarly for $T \subseteq V^*$); moreover, $S^{\perp\perp\perp} = S^\perp$ (and similarly for $T \subseteq V^*$).
- $(A \cap B)^\perp = A^\perp + B^\perp$ and $(A + B)^\perp = A^\perp \cap B^\perp$ for subspaces $A, B \leq V$.

In particular passing to the orthogonal induces an antiisomorphism between the lattices of subspaces of V and V^* .

EXERCISE LA.15 \square Let $f : V \rightarrow W$ be a linear map; show that $\ker(f^*) = (\operatorname{im} f)^\perp$ and $\operatorname{im}(f^*) = (\ker f)^\perp$. Deduce that the rank of f equals the rank of f^*

EXERCISE LA.16 \square Show that given a subspace $A \leq V$ of a finite dimensional k -vector space, the inclusion $A \hookrightarrow V$ induces a surjective linear map $V^* \rightarrow A^*$, having kernel A^\perp ; deduce that $(V/A)^* \cong A^\perp$ as vector spaces.

EXERCISE LA.17 \square More generally, given subspaces $U \leq W$ of the same V show that $(W/U)^*$ is isomorphic to U^\perp/W^\perp .

With this –but also directly– show that

- $((U + W)/W) \cong W^\perp/(U^\perp \cap W^\perp)$;
- $(W/(U \cap W))^* \cong (U^\perp + W^\perp)/W^\perp$.

EXERCISE LA.18 \square Show that the vector space $\mathbf{Q}[X]$ of polynomials with rational coefficients is not isomorphic to its dual $\mathbf{Q}[X]^*$.

EXERCISE LA.19 \square Sia $g : V \rightarrow W$ una mappa lineare; dimostrare che l'equazione $g(v) = w$ ammette una soluzione, fissato $w \in W$, se e solo se $w \in \ker g^*$, dove $g^* : W^* \rightarrow V^*$ è la trasposta di g .

EXERCISE LA.20 \square Let I be a set and V_i a family of F -vector spaces indexed by I ; the *product* $\prod_{i \in I} V_i$ of the family V_i is the set obtained from the cartesian product of the V_i and equipped with componentwise sum and scalar multiplication.

Show that the canonical projections $\pi_j : \prod_i V_i \rightarrow V_j$ for $j \in I$ is F -linear.

Show that $\prod_i V_i$ enjoys the following property:

Given any other family of F -linear maps $f_i : W \rightarrow V_i$, there exists a unique $f : W \rightarrow \prod_i V_i$ which is F -linear and makes the following diagram commute for every $j \in I$:

$$\begin{array}{ccc}
 W & \xrightarrow{\quad f \quad} & \prod_i V_i \\
 \searrow f_j & & \swarrow \pi_j \\
 & & V_j
 \end{array} \tag{4.7}$$

Usually the notation for such an f is $\prod_i f_i$

EXERCISE LA.21 \square Let I be a set and V_i a family of F -vector spaces indexed by I ; the *coproduct* (or *direct sum*) $\sum_{i \in I} V_i$ of the family V_i is the subset of $\prod_i V_i$, equipped with componentwise sum and scalar multiplication, whose elements are those sequences of vectors $(v_i \mid i \in I, v_i \in V_i)$ such that $v_i \neq 0$ for at most a finite number of indices.

Show that this is indeed a vector subspace of $\prod_i V_i$, and that the inclusions $\iota_j : \sum_i V_i \rightarrow V_j$ are F -linear for every $j \in I$.

Show that $\sum_i V_i$ enjoys the following property:

Given any other family of F -linear maps $f_i : V_i \rightarrow W$, there exists a unique $f : \sum_i V_i \rightarrow W$ which is F -linear and makes the following diagram commute for every $j \in I$:

$$\begin{array}{ccc}
 & V_j & \\
 \iota_j \swarrow & & \searrow f_j \\
 \sum_i V_i & \xrightarrow{f} & W
 \end{array} \tag{4.8}$$

Usually the notation for such an f is $\sum_i f_i$.

EXERCISE LA.22 \square Prove that if I is a finite set and V_i a family of vector spaces indexed by I , then there is an isomorphism $\prod_i V_i \cong \sum_i V_i$ (hint: start from $I = \{1, 2\}$ a set with two elements).

EXERCISE LA.23 \square Let S be a finite set of cardinality n ; consider the vector space $V = \{f : 2^S \rightarrow \mathbf{R}\}$ of all functions from 2^S to \mathbf{R} , and the linear map $\varphi : V \rightarrow V$ given by

$$\varphi(f) : T \mapsto \sum_{Y \supseteq T} fY \tag{4.9}$$

Show that φ is an isomorphism of vector spaces. (Hint: induction on n .)

DEFINITION 4.2. Let $k \geq 1$ be an integer, and $\mathcal{W} = (W_1, \dots, W_k)$ and $\mathcal{W}' = (W'_1, \dots, W'_k)$ be two k -tuples of subspaces of the same vector space V ; we say that $\mathcal{W}, \mathcal{W}'$ are *concordant* if there exists an invertible linear map $\varphi : V \rightarrow V$ such that $\varphi(W_i) = W'_i$ for each $1 \leq i \leq k$.

EXERCISE LA.24 \square Show that two tuples of subspaces $\mathcal{W}, \mathcal{W}'$ as above are concordant if and only if for each $1 \leq k \leq n$ and each choice of indices (i_1, \dots, i_k) there is a linear isomorphism $\varphi_{i_1 \dots i_k} : W_{i_1, \dots, i_k} \rightarrow W'_{i_1, \dots, i_k}$, where we denote $W_{i_1 \dots i_k} = W_{i_1} \cap \dots \cap W_{i_k}$.

EXERCISE LA.25 \square Show that a k -tuple $\mathcal{W} = (W_1, \dots, W_n)$ of subspaces of V is equivalently described by its *nerve*: let again $W_{i_1 \dots i_r} := W_{i_1} \cap \dots \cap W_{i_r}$, and consider the diagram

$$\begin{array}{ccccccc}
 \sum_{i=1}^k W_i & \begin{array}{c} \xrightarrow{\pi_1} \\ \xleftarrow{\iota_1} \\ \xrightarrow{\pi_2} \end{array} & \sum_{i_1 < i_2} W_{i_1 i_2} & \begin{array}{c} \xrightarrow{\pi_1} \\ \xleftarrow{\iota_1} \\ \xleftarrow{\pi_2} \\ \xrightarrow{\pi_3} \end{array} & \sum_{i_1 < i_2 < i_3} W_{i_1 i_2 i_3} & \begin{array}{c} \xrightarrow{\pi_1} \\ \xleftarrow{\iota_3} \\ \dots \\ \xleftarrow{\iota_1} \\ \xrightarrow{\pi_4} \end{array} & \dots & \tag{4.10}
 \end{array}$$

where the maps pointing to the right are induced by projections from the intersection of i elements of the tuple to the intersection of $(i + 1)$ elements of the tuple, and the arrows pointing to the left are induced by the inclusions from an intersection of i elements to an intersection of $(i - 1)$ elements of the tuple.

Given two of such diagrams, filled by the dotted arrows $\varphi_{i_1, \dots, i_r}$ below,

$$\begin{array}{c}
 \sum_{i=1}^k W_i \rightleftarrows \sum_{i_1 < i_2} W_{i_1 i_2} \rightleftarrows \sum_{i_1 < i_2 < i_3} W_{i_1 i_2 i_3} \rightleftarrows \dots \\
 \downarrow \Sigma_i \varphi_i \quad \downarrow \Sigma \varphi_{i_1 i_2} \quad \downarrow \Sigma \varphi_{i_1 i_2 i_3} \\
 \sum_{i=1}^k W'_i \rightleftarrows \sum_{i_1 < i_2} W'_{i_1 i_2} \rightleftarrows \sum_{i_1 < i_2 < i_3} W'_{i_1 i_2 i_3} \rightleftarrows \dots
 \end{array} \quad (4.11)$$

the two k -tuples that they represent are concordant if and only if for each choice of homonymous horizontal arrows pointing in the same direction, the resulting diagram is commutative: in other words, all diagrams

$$\begin{array}{ccc}
 \sum_{i=1}^k W_i \xrightarrow{\pi_1} \sum_{i_1 < i_2} W_{i_1 i_2} & \sum_{i=1}^k W_i \xrightarrow{\pi_2} \sum_{i_1 < i_2} W_{i_1 i_2} & \sum_{i=1}^k W_i \xleftarrow{\iota_1} \sum_{i_1 < i_2} W_{i_1 i_2} \\
 \Sigma_i \varphi_i \downarrow \quad \Sigma \varphi_{i_1 i_2} \downarrow & \Sigma_i \varphi_i \downarrow \quad \Sigma \varphi_{i_1 i_2} \downarrow & \Sigma_i \varphi_i \downarrow \quad \Sigma \varphi_{i_1 i_2} \downarrow \\
 \sum_{i=1}^k W'_i \xrightarrow{\pi'_1} \sum_{i_1 < i_2} W'_{i_1 i_2} & \sum_{i=1}^k W'_i \xrightarrow{\pi'_2} \sum_{i_1 < i_2} W'_{i_1 i_2} & \sum_{i=1}^k W'_i \xleftarrow{\iota_1} \sum_{i_1 < i_2} W'_{i_1 i_2}
 \end{array} \quad (4.12)$$

etc., are commutative.

EXERCISE LA.26 \square Let q be a prime number, and $n \geq 1$ an integer; build a field having exactly q^n elements by considering the roots of the polynomial $X^{q^n} - X$; prove that every finite field extension $\mathbf{Z}/q\mathbf{Z}$ arises in this way, or in other words, prove that if F is a finite field of characteristic q , then it has q^n elements for some $n \geq 1$, and it is *the unique* such field up to isomorphism.

EXERCISE LA.27 \square Let $n \geq 1$ be an integer. Prove that the number of subspaces of the vector space $V = (\mathbf{Z}/q\mathbf{Z})^n$ having dimension $k \leq n$ is the q -binomial coefficient

$$\binom{n}{k}_q := \frac{[n]!_q}{[k]!_q [n-k]!_q} \quad (4.13)$$

where for every real number q (so a fortiori for an integer) the quantity $[r]!_q$ is defined as $(1+q)(1+q+q^2) \dots (1+q+\dots+q^{r-1})$;

EXERCISE LA.28 \square Prove that the set of all invertible linear maps of $V = (\mathbf{Z}/q\mathbf{Z})^n$ onto itself has $q^{\binom{n}{2}} (q-1)^n [n]!_q$ elements.

EXERCISE LA.29 \square Let F be a field, and G the subgroup of $n \times n$ matrices having the property that

There exists *exactly* one 1 in each row and in each column.

Prove that there exists a group isomorphism between G and the symmetric group $S(n)$ of all permutations of an n -element set. Is the subgroup G normal in $GL_n(F)$? Are the matrices in G diagonalizable? Who is the characteristic polynomial of a matrix $\Sigma \in G$?

EXERCISE LA.30 \square Let F be a field; interpret the determinant

$$\det := \sum_{\sigma \in S(n)} (-1)^{|\sigma|} \prod_{i=1}^n a_{i, \sigma i} \quad (4.14)$$

of a generic $n \times n$ matrix with coefficients in F as a polynomial in the n^2 indeterminates $\{a_{ij} \mid 1 \leq i, j \leq n\}$.

Prove that \det is an irreducible polynomial in $F[a_{ij} \mid 1 \leq i, j \leq n]$. (Hint: induction on n , starting from $n = 2$.)

EXERCISE LA.31 \square Consider the F -vector space $F^{V \times W} = \sum_{V \times W} F$, and the subspace generated by the relations

$$\begin{cases} (v_1, w) + (v_2, w) \sim (v_1 + v_2, w), \\ (v, w_1) + (v, w_2) \sim (v, w_1 + w_2), \\ c(v, w) \sim (cv, w), \\ c(v, w) \sim (v, cw). \end{cases} \quad (4.15)$$

in $F^{V \times W}$. Show that this quotient satisfies the universal property of the tensor product $V \otimes W$ of V, W :

To each *bilinear* map $\varphi : V \times W \rightarrow U$ corresponds a unique *linear* map $\bar{\varphi} : V \otimes W \rightarrow U$, bijectively.

Prove that when V, W have finite dimension over the field F , there is an isomorphism $V \otimes W \cong \text{Bil}(V \times W, F)^*$ (on the right hand side: the dual of the vector space of bilinear maps $V \times W \rightarrow F$).

EXERCISE LA.32 \square Prove that $V \otimes W$ has a basis $\mathcal{V} \otimes \mathcal{W} = \{v_i \otimes w_j\}$, when v_i runs over a basis \mathcal{V} of V , and w_j runs over a basis \mathcal{W} of W . In the identification $V \otimes W \cong \text{Bil}(V \times W, F)^*$, to which elements of $\text{Bil}(V \times W, F)^*$ does $v_i \otimes w_j$ correspond?

Prove that $F \otimes F \cong F$; prove that $V \otimes F \cong V$ for each vector space V ; prove that $V \otimes (W \otimes Z) \cong (V \otimes W) \otimes Z$; prove that $V \otimes W \cong W \otimes V$.

EXERCISE LA.33 \square Prove that each pair of linear maps $f : V \rightarrow V', g : W \rightarrow W'$ induces a linear map $f \otimes g : V \otimes W \rightarrow V' \otimes W'$; how is this map defined on the basis $\mathcal{V} \otimes \mathcal{W}$ of $V \otimes W$? Assume V, V', W, W' all have finite dimension and fix bases $\mathcal{V}, \mathcal{W}, \mathcal{V}', \mathcal{W}'$ for all vector spaces in question; let A be the matrix of f , and B be the matrix of g , and find an

expression for the matrix of $f \otimes g$ in terms of A, B ; the matrix $A \otimes B$ is called the *Kronecker product* of A, B .

EXERCISE LA.34 \square Show the following properties of the Kronecker product:

- given matrices A, B, C , prove that if A, B have the same size,

$$(A + B) \otimes C = A \otimes C + B \otimes C \quad C \otimes (A + B) = C \otimes A + C \otimes B;$$

- given matrices A, B, C prove that

$$A \otimes (B \otimes C) = (A \otimes B) \otimes C \tag{4.16}$$

- if A, B, C, D are matrices such that the matrix products AC, BD exist, then $(A \otimes B)(C \otimes D) = AC \otimes BD$; as a corollary, if A, B are both invertible, so is $A \otimes B$ and $(A \otimes B)^{-1} = A^{-1} \otimes B^{-1}$;
- if A, B are matrices, respectively $n \times n$ and $m \times m$, then

$$\det(A \otimes B) = (\det A)^m (\det B)^n \quad \text{trace}(A \otimes B) = \text{trace}(A)\text{trace}(B). \tag{4.17}$$

EXERCISE LA.35 \square The *tensor algebra* over a vector space V is defined as the vector space

$$TV = \sum_{n \geq 0} V^{\otimes n} \tag{4.18}$$

equipped with the *tensor product* operation: two elements $x_1 \otimes \cdots \otimes x_n$ and $y_1 \otimes \cdots \otimes y_m$ can be multiplied into an element

$$x_1 \otimes \cdots \otimes x_n \otimes y_1 \otimes \cdots \otimes y_m \in V^{\otimes(n+m)} \tag{4.19}$$

Prove that the tensor product operation equips TV with a *graded algebra* structure.

Prove that when V is finite dimensional, say with dimension d , TV is isomorphic – as algebra – to the algebra of *noncommutative polynomials* $k\{X_1, \dots, X_d\}$.

EXERCISE LA.36 \square For what has been shown in [Exercise 35](#) TV is a ring;

- Describe the ideal I in (TV, \otimes) generated by the set $S = \{v \otimes w - w \otimes v \mid v, w \in V\}$; prove that the quotient TV/I is isomorphic to the ring of polynomials $k[X_1, \dots, X_d]$; define explicitly the isomorphism.
- Describe the ideal J generated by the set $\{v \otimes w + w \otimes v \mid v, w \in V\}$; prove that J is also generated by the set $\{v \otimes v \mid v \in V\}$;
- Fix a basis $\{e_1, \dots, e_n\}$ of V . Describe the ideal L generated by the set $\{e_i \otimes e_i - 1, e_i \otimes e_j + e_j \otimes e_i \mid e_i, e_j \in V\}$.

The quotient TV/J is the *exterior* (or *Grassmann*) algebra $\wedge(V)$ of V constructed from the tensor algebra. The quotient TV/L is the *Clifford algebra* $Cl(V)$ of V . See [Exercise 37](#) for more information on the Grassmann algebra; see [Definition 4.5](#) for more information on the Clifford algebra.

EXERCISE LA.37 \square Prove that the tensor product operation on TV induces a multiplication

$$_ \wedge _ : \wedge(V) \times \wedge(V) \longrightarrow \wedge(V) \tag{4.20}$$

on the quotient that defines $\wedge(V)$.

EXERCISE LA.38 \square Prove that in fact the ideal J such that $\wedge(V) \cong TV/J$ is a *graded* ideal: there exists a decomposition $J = \bigoplus_{r \geq 0} J_r$ where $J_r := J \cap V^{\otimes r}$ such that if we pose $\wedge_r(V) := V^{\otimes r}/J_r$ there exists a decomposition

$$\wedge(V) \cong \bigoplus_{r \geq 0} \wedge_r(V) \tag{4.21}$$

Note that $\wedge_0(V) \cong F$ (the base field) and $\wedge_1(V) \cong V$, so that there exists a canonical F -linear map $j : V \hookrightarrow \wedge(V)$.

EXERCISE LA.39 \square Prove that if V has dimension d and $r > d$, then $\wedge_r(V) \cong (0)$ (the zero vector space). In fact, prove that a basis of $\wedge_r(V)$ can be found, once a basis $\mathcal{V} = \{e_1, \dots, e_d\}$ of V has been fixed, taking the elements

$$e_{i_1} \wedge \dots \wedge e_{i_r} \tag{4.22}$$

where $1 \leq i_1 < \dots < i_r \leq d$. From this, $\dim_F \wedge_r(V) = \binom{d}{r}$.

EXERCISE LA.40 \square Prove the universal property of $\wedge(V)$:

Given a F -algebra (A, \cdot) and a F -linear map $f : V \rightarrow A$ such that for each $v \in V$, $f v \cdot f v = 0$ in A , there exists a unique extension $\bar{f} : \wedge(V) \rightarrow A$ in

$$\begin{array}{ccc} & V & \\ j \swarrow & & \searrow f \\ \wedge(V) & \xrightarrow{\quad \bar{f} \quad} & A \end{array} \tag{4.23}$$

that makes the diagram commute.

EXERCISE LA.41 \square As a corollary of the above universal property, given any linear map $f : V \rightarrow W$ of vector spaces, there exists a unique $\bar{f} = \wedge f : \wedge(V) \rightarrow \wedge(W)$ such that the diagram

$$\begin{array}{ccc} V & \xrightarrow{f} & W \\ j_V \downarrow & & \downarrow j_W \\ \wedge(V) & \xrightarrow{\wedge f} & \wedge(W) \end{array} \tag{4.24}$$

is commutative.

EXERCISE LA.42 \square Show that

- $\wedge f$ is uniquely determined by the request that $\wedge f(v_1 \wedge \cdots \wedge v_r) = f v_1 \wedge \cdots \wedge f v_r$ for every $r \geq 1$ and $v_1, \dots, v_r \in V$;
- $\wedge(g \circ f) = \wedge g \circ \wedge f$ for;
- $\wedge(\text{id}_V) = \text{id}_{\wedge V}$.

EXERCISE LA.43 \square Let $x \in \wedge_r(V)$, $y \in \wedge_s(V)$; then $y \wedge x = (-1)^{rs} x \wedge y$; in particular, the exterior product is *anticommutative*, i.e. for any vectors v, w one has $v \wedge w = -w \wedge v$.

EXERCISE LA.44 \square Let $\{W_i \mid i \in I\}$ be a family of F -vector spaces; prove the isomorphism

$$V \otimes \left(\bigoplus_{i \in I} W_i \right) \cong \bigoplus_{i \in I} V \otimes W_i \quad (4.25)$$

and deduce that

$$\wedge(V) \otimes \wedge(W) \cong \bigoplus_{r,s \geq 0} \wedge_r(V) \otimes \wedge_s(W) \quad (4.26)$$

a relation from which one can prove the *exponential property* for $\wedge(_)$:

$$\wedge(V \oplus W) \cong \wedge(V) \otimes \wedge(W) \quad (4.27)$$

Use the exponential property to find the dimension of $\wedge(V)$ by induction on $d = \dim_F V$.

EXERCISE LA.45 \square Let $f : V \rightarrow W$ be a linear map of F -vector spaces; consider the induced map

$$\wedge f : \wedge(V) \longrightarrow \wedge(W) \quad (4.28)$$

between the exterior algebras.

Find a matrix expression for $\wedge f$ once bases $\mathcal{V} = \{v_1, \dots, v_n\}$ and $\mathcal{W} = \{w_1, \dots, w_m\}$ of V, W respectively have been fixed.

DEFINITION 4.3. A *super vector space* is a vector space V over a field k that can be decomposed as a direct sum $V = V_0 \oplus V_1$; the elements of the form $(v, 0)$ for $v \in V_0$ are called *even*, while the elements $(0, v)$ for $v \in V_1$ are called *odd*. A homomorphism of super vector spaces V, W is a linear map $f : V \rightarrow W$ that preserves the parity of elements (equivalently: such that $f(V_0) \subseteq W_0, f(V_1) \subseteq W_1$).

EXERCISE LA.46 \square Define the *dimension* of a super vector space V to be $\dim V = \dim V_0 - \dim V_1$. Define the direct sum and intersection of super vector subspaces (define them first!), or prove that it doesn't always exist. Does the Grassmann formula still hold true?

EXERCISE LA.47 \square How to define the cartesian product of super vector spaces V, W ? How to define their direct sum? Is it still true that finite products and finite coproducts coincide, $\prod_{i=0}^n V_i \cong \sum_{i=0}^n V_i$, as it happens for vector spaces? [Hint: the \prod and \sum constructions should coincide with the usual ones for 'purely even' super vector spaces, i.e. it should be true that

$$(V, 0) \times (W, 0) \cong (V \times W, 0) \quad (V, 0) \oplus (W, 0) \cong (V \oplus W, 0). \quad (4.29)$$

Given this, find a definition for the 'purely odd' case and for the mixed case.]

DEFINITION 4.4. An *exact sequence* of linear maps is a string of vector spaces and composable linear maps

$$\mathbf{v} : V_0 \xrightarrow{f_0} V_1 \xrightarrow{f_1} \cdots \xrightarrow{f_n} V_{n+1} \quad (4.30)$$

such that $\ker f_{i+1} = \operatorname{im} f_i$ for every $i = 0, \dots, n$.

Note that this defining property implies in particular that $f_{i+1} \circ f_i = 0$; when this weaker condition is satisfied we say that the sequence of spaces and linear maps forms a *chain complex* (so every exact sequence is a chain complex, but the converse does not hold). If \mathbf{v} is a complex, the quotient spaces $H^i(\mathbf{v}) = \ker f_{i+1} / \operatorname{im} f_i$ are an intrinsic invariant of the complex and of great theoretical interest.

EXERCISE LA.48 \square Let

$$0 \rightarrow U \rightarrow V \rightarrow W \rightarrow 0 \quad (4.31)$$

be an exact sequence; show that $\dim U - \dim V + \dim W = 0$; more generally, let

$$\mathbf{v} : 0 \rightarrow V_1 \rightarrow V_2 \rightarrow \cdots \rightarrow V_n \rightarrow 0 \quad (4.32)$$

be an exact sequence; show that the alternating sum of dimensions $\sum (-1)^i \dim V_i$ equals 0.

This leads to the following definition, valid for every complex: let

$$\mathbf{v} : 0 \rightarrow V_1 \rightarrow V_2 \rightarrow \cdots \rightarrow V_n \rightarrow 0 \quad (4.33)$$

be a complex of vector spaces; we define the following objects:

- the i -th *cohomology group* $H^i(\mathbf{v})$ is defined as the quotient $\ker f_{i+1} / \operatorname{im} f_i$, and the total cohomology of \mathbf{v} is defined as $\sum H^i(\mathbf{v})$;
- the i -th *Betti number* $b_j(\mathbf{v})$ of \mathbf{v} is defined as $\dim H^i(\mathbf{v})$;
- the *Euler characteristic* of \mathbf{v} is defined as the signed sum $\sum (-1)^j b_j(\mathbf{v})$.

EXERCISE LA.49 \square Now, let

$$\mathbf{v} : 0 \rightarrow V_1 \rightarrow V_2 \rightarrow \cdots \rightarrow V_n \rightarrow 0 \quad (4.34)$$

be a complex and let $d_j = \dim V_j$; show that the following conditions are equivalent:

- the complex \mathbf{v} is exact;
- each $H^i(\mathbf{v})$ is zero;
- the Euler characteristic of the complex, $\sum (-1)^j d_j$, is zero.

Let R be a (commutative, unital) ring; an R -*module* consists of the exact same thing a vector space is, with the only difference that the ‘scalar multiplication’ operation now takes values in R , a ring that is not necessarily a field. This seemingly innocuous difference generates a lot of differences between the theory of modules and the theory of vector spaces.

EXERCISE LA.50 \square What is a module over the ring $\mathbf{R}[X]$ of polynomials with real coefficients?

EXERCISE LA.51 \square Let I be a set; a *family* of R -modules M_i is a set $\{M_i \mid i \in I\}$ of modules over R ; the *direct sum* of a family of R -modules is a module $\sum_{i \in I} M_i$ such that the following properties are satisfied:

- There exists a family of R -linear maps $\{\iota_j : M_j \rightarrow \sum_{i \in I} M_i \mid j \in I\}$ such that
- given any other family of R -linear maps $\{f_j : M_j \rightarrow X \mid j \in I\}$, there exists a unique $\tilde{f} : \sum_{i \in I} M_i \rightarrow X$ such that $\tilde{f} \circ \iota_j = f_j$.

Prove that whenever another object $\sum_{i \in I}^* M_i$ satisfies the same properties, then there exists a unique isomorphism $\sum_{i \in I}^* M_i \cong \sum_{i \in I} M_i$.

EXERCISE LA.52 \square The *direct product* of a family of R -modules M_i is the module $\prod_{i \in I} M_i$ such that

- There exists a family of R -linear maps $\{\pi_j : \prod_{i \in I} M_i \rightarrow M_j \mid j \in I\}$ such that
- given any other family of R -linear maps $\{f_j : X \rightarrow M_j \mid j \in I\}$, there exists a unique $\tilde{f} : X \rightarrow \prod_{i \in I} M_i$ such that $\pi_j \circ \tilde{f} = f_j$.

Prove that whenever another object $\prod_{i \in I}^* M_i$ satisfies the same properties, then there exists a unique isomorphism $\prod_{i \in I}^* M_i \cong \prod_{i \in I} M_i$.

EXERCISE LA.53 \square Prove that when I is a finite set, $\prod_{i \in I} A_i \cong \sum_{i \in I} A_i$.

EXERCISE LA.54 \square Is it true or false that $\prod_{i \in I} \sum_{j \in J} A_{ij} \cong \sum_{j \in J} \prod_{i \in I} A_{ij}$, for every family of R -modules $\{A_{ij} \mid (i, j) \in I \times J\}$?

EXERCISE LA.55 \square Prove that given an R -module M and a set I , one has $M \cong \sum_{i \in I} M_i$ if and only if one can find homomorphisms $\mu^j : M_j \rightarrow M$ and $\rho_j : M \rightarrow M_j$ for each $j \in I$ satisfying the following properties:

- $\rho_i \circ \mu^i = \text{id}_{M_i}$ for each $i \in I$;
- $\rho_i \circ \mu^j = 0$ for each $i, j \in I, i \neq j$;
- $\rho_i(x) = 0$ for each $x \in M$, for almost all indices $i \in I$;⁶
- $\sum_{i \in I} \mu^i \circ \rho_i = \text{id}_M$.

EXERCISE LA.56 \square An R -module M is called *free* if it is of the form $R^{(\Gamma)} = \sum_{\gamma \in \Gamma} R$, for a set Γ .

Find an R -module that is not free, when $R = \mathbf{Z}$; show that a \mathbf{Z} -module is free when ‘it has a basis’ in the sense of vector spaces (but be careful, it is possible to find rings where free modules do not have a well-defined dimension).

EXERCISE LA.57 \square Prove that if M is a free module over an infinite set of generators, then $M \cong M \oplus M$; deduce that there is an isomorphism between the abelian groups $\text{End}(M)$ and $\text{End}(M) \times \text{End}(M)$. Is this isomorphism also a *ring* isomorphism?

⁶A notation with which it’s better to familiarise soon: *almost all* elements of a set means ‘all, but possibly a finite number’.

EXERCISE LA.58 \square Let

$$\begin{array}{ccccccc}
 & & L & \xrightarrow{u_1} & M & \xrightarrow{u_2} & N \xrightarrow{u_3} 0 \\
 & & \downarrow f & & \downarrow g & & \downarrow h \\
 0 & \xrightarrow{d_1} & L' & \xrightarrow{d_2} & M' & \xrightarrow{d_3} & N'
 \end{array} \tag{4.35}$$

be a diagram of F -linear maps between vector spaces, with the property that the rows are exact sequences (cf. Definition 4.4). Show that there exists an exact sequence

$$\ker f \rightarrow \ker g \rightarrow \ker h \rightarrow \operatorname{coker} f \rightarrow \operatorname{coker} g \rightarrow \operatorname{coker} h \tag{4.36}$$

where $\operatorname{coker} f$ is the ‘cokernel’ of f , i.e. the quotient space $L'/\operatorname{im} f$, and similarly for $\operatorname{coker} g$ and $\operatorname{coker} h$.

EXERCISE LA.59 \square The tensor product $V \otimes_{\mathbb{Z}} W$ of two super vector spaces $V = (V_0, V_1), W = (W_0, W_1)$ over the same field F is the usual vector space $V \otimes W$ equipped with the $\mathbb{Z}/2\mathbb{Z}$ -graduation

$$(V \otimes W)_l = \sum_{i+j=l \pmod 2} V_i \otimes W_j$$

. What is the universal property of this object?

EXERCISE LA.60 \square Show that there exists a super vector space J acting as ‘the square root of -1 ’, in the sense that J is not the tensor unit and $J \otimes_{\mathbb{Z}} J \cong F$.

A *super algebra* over a field F is a super vector space V equipped with an F -bilinear multiplication operation: this means that there is an operation $V \times V \rightarrow V : (a, b) \mapsto a \cdot b$, such that $(a + b) \cdot c = a \cdot c + b \cdot c$ and $a \cdot (b + c) = a \cdot b + a \cdot c$, and $(\alpha a) \cdot (\beta b) = \alpha\beta(a \cdot b)$ for each $a, b, c \in V$ and $\alpha \in F$.

DEFINITION 4.5. Given a vector $v \in \mathbb{R}^2$ define a binary operation, the *Clifford norm*, as follows:

$$(a_1 e_1 + a_2 e_2) \bullet (a_1 e_1 + a_2 e_2) = a_1^2 e_1 \bullet e_1 + a_2^2 e_2 \bullet e_2 + a_1 a_2 (e_1 \bullet e_2 + e_2 \bullet e_1)$$

where (a_1, a_2) are the coordinates of v in the standard basis of \mathbb{R}^2 . Now, if on this expression we impose the relation $v \bullet v = v \cdot v \in \mathbb{R}$ (where $v \cdot v$ is the scalar product of vectors), from the above equation we get that

$$e_i \bullet e_i = 1 \qquad e_1 \bullet e_2 = -e_2 \bullet e_1.$$

Define the *Clifford algebra* $Cl(2, \mathbb{R})$ as the set of elements of the form $a + b_1 e_1 + b_2 e_2 + \underline{c} e_{12}$, where $a \in \mathbb{R}$ is the *scalar part* of a Clifford vector x , $\vec{b} = (b_1, b_2) \in \mathbb{R}^2$ the *vector part* of x , and $\underline{c} \in \mathbb{R}$ its *bivector part* (where for the sake of brevity we write $e_{12} = e_1 \bullet e_2$); each of these three parts is a different *homogeneous component* of $x \in Cl(2, \mathbb{R})$.

EXERCISE LA.61 \square The set $Cl(2, \mathbf{R})$ is a vector space, where the vector space operations are done componentwise (with the sum in \mathbf{R} in the scalar and bivector part, and with the sum in \mathbf{R}^2 in the vector part). Prove that this is in fact a vector space.

EXERCISE LA.62 \square Find an explicit formula for the Clifford product of two elements of $Cl(2, \mathbf{R})$:

$$(x + \vec{y} + \underline{z}) \bullet (x' + \vec{y}' + \underline{z}') = \dots \quad (4.37)$$

EXERCISE LA.63 \square Prove that the Clifford product of two homogeneous elements of vector type, $v = (v_1, v_2)$, $w = (w_1, w_2)$ sn't homogeneous any more, and in fact it decomposes as a nontrivial scalar part plus a nontrivial bivectorial part: the Clifford product of two vectors in $Cl(2, \mathbf{R})$ is

$$\vec{v} \bullet \vec{w} = \vec{v} \cdot \vec{w} + (\vec{v} \wedge \vec{w})e_{12}$$

where $\vec{v} \cdot \vec{w} = v_1w_1 + v_2w_2$ is the dot product of vectors and $\vec{v} \wedge \vec{w} = v_1w_2 - v_2w_1$ their cross product.

EXERCISE LA.64 \square As a consequence of the previous exercise, prove the relations

- $\vec{v} \cdot \vec{w} = \frac{1}{2}(\vec{v} \bullet \vec{w} + \vec{w} \bullet \vec{v})$;
- $\vec{v} \wedge \vec{w} = \frac{1}{2}(\vec{v} \bullet \vec{w} - \vec{w} \bullet \vec{v})$;
- $\vec{a} \parallel \vec{b} \iff \vec{a} \wedge \vec{b} = 0 \iff \vec{a} \bullet \vec{b} = \vec{a} \cdot \vec{b}$;
- $\vec{a} \perp \vec{b} \iff \vec{a} \cdot \vec{b} = 0 \iff \vec{a} \bullet \vec{b} = \vec{a} \wedge \vec{b}$.

EXERCISE LA.65 \square Prove that the assignment sending $1 \mapsto \mathbf{I}_2$ (the 2×2 identity matrix), $e_1 \mapsto \begin{pmatrix} 1 & \\ & -1 \end{pmatrix}$, $e_2 \mapsto \begin{pmatrix} & 1 \\ 1 & \end{pmatrix}$, $e_{12} \mapsto \begin{pmatrix} & 1 \\ -1 & \end{pmatrix}$ defines an algebra isomorphism θ between $Cl(2, \mathbf{R})$ and $M_2(\mathbf{R})$ (the algebra of 2×2 matrices with real coefficients); does this isomorphism depend on the choice we made for a basis of the two spaces?

EXERCISE LA.66 \square Translate the matrix operations on $M_2(\mathbf{R})$ into operations in $Cl(2, \mathbf{R})$ along the isomorphism of [Exercise 65](#):

- *transposition* of a matrix corresponds to changing the sign of the bivector part in the associated Clifford vector \mathbf{u} : if \mathbf{u} and the matrix A correspond each other under θ , then A^t corresponds to the Clifford vector $\widetilde{\mathbf{u}} = u_0 + u_1e_1 + u_2e_2 - u_{12}e_{12}$;
- *inversion* of a matrix corresponds to taking the *Clifford conjugate* of \mathbf{u} : if \mathbf{u} and the matrix A correspond each other under θ , and A is invertible, then A^{-1} corresponds to the Clifford vector $\overline{\mathbf{u}} = u_0 - u_1e_1 - u_2e_2 - u_{12}e_{12}$.

EXERCISE LA.67 \square Prove that $(\widetilde{-}), (\overline{-})$ are involutive algebra antiautomorphisms: $\widetilde{\widetilde{\mathbf{u}}} = \mathbf{u}$, $\overline{\overline{\mathbf{u}}} = \mathbf{u}$, $\widetilde{\overline{\mathbf{u}}} = \overline{\widetilde{\mathbf{u}}}$, $\overline{\widetilde{\mathbf{u}}} = \widetilde{\overline{\mathbf{u}}}$. How do $\widetilde{\mathbf{u}}, \overline{\mathbf{u}}$ relate to each other?

EXERCISE LA.68 \square Show that $Cl(2, \mathbf{R})$ contains a subalgebra isomorphic to the field \mathbf{C} of complex numbers, represented as the set of matrices

$$\left\{ \begin{pmatrix} a & -b \\ b & a \end{pmatrix} \mid a, b \in \mathbf{R} \right\} \cong \{a + be_{12} \mid a, b \in \mathbf{R}\}$$

with real entries; in fact, e_{12} behaves like the imaginary unit, in that there exists a decomposition $Cl(2, \mathbf{R}) \cong (\mathbf{R} \oplus e_{12}\mathbf{R}) \oplus (e_1\mathbf{R} \oplus e_2\mathbf{R})$ (prove it!).

Prove that this decomposition makes $Cl(2, \mathbf{R})$ a super algebra.

CHAPTER 2

Category theory

1. Categories, functors, naturality

EXERCISE CF.1 \square Verify that the following are examples of categories:

- **Set:** objects are set, morphisms are functions between sets, composition is composition of functions, the identity of a given set A is the identity function $A \rightarrow A : a \mapsto a$.
- **Alg(T):** A *signature* is a family $T = (n_i)_{i \in I}$ of natural numbers n_i , indexed by a set I . Let $T = (n_i)$ be a signature; define a category $\text{Alg}(T)$ as follows: A T -algebra is a pair $(X, (t_i)_{i \in I})$ consisting of a set X and a family of functions $t_i : X^{n_i} \rightarrow X$, called n_i -ary operations on X . A T -homomorphism

$$f : (X, (t_i)_{i \in I}) \longrightarrow (Y, (s_i)_{i \in I}) \quad (1.1)$$

is a function $f : X \rightarrow Y$ for which the diagram

$$\begin{array}{ccc} X^{n_i} & \xrightarrow{f^{n_i}} & Y^{n_i} \\ t_i \downarrow & & \downarrow s_i \\ X & \xrightarrow{f} & Y \end{array} \quad (1.2)$$

is commutative, meaning that for every operation $t_i : X^{n_i} \rightarrow X$, $s_i : Y^{n_i} \rightarrow Y$ and every n_i -tuple of elements x_1, \dots, x_{n_i} one has

$$f(t_i(x_1, \dots, x_{n_i})) = s_i(fx_1, \dots, fx_{n_i}) \quad (1.3)$$

Make precise the fact that as a corollary, all classes of algebraic structures are categories.

- **Rel:** objects are sets; given sets A, B , the set of morphisms $A \rightarrow B$ is the powerset of $A \times B$, $\text{Rel}(A, B) := 2^{A \times B}$. Given a relation $R \in 2^{A \times B}$ and a relation $S \in 2^{B \times C}$, define the composition $S \circ R \in 2^{A \times C}$ as the subset

$$(a, c) \in S \circ R \iff \exists b \in B, (a, b) \in R, (b, c) \in S. \quad (1.4)$$

Prove that this defines an associative composition operation.¹ Given this composition law, there is only one possible choice for the identity relation $I \subset A \times A$: find it.

- **Pred**: the category of *predicates*, having objects the pairs (A, X) where A is a subset of X , and where morphisms $(A, X) \rightarrow (B, Y)$ are the functions $f : X \rightarrow Y$ such that $f(a) \in B$ for each $a \in A$;
- **Set***: the category obtained from the previous one, restricting to the objects (A, X) for which A is a set with a single element. Make precise the statement that this is the category of ‘pointed sets’ and ‘basepoint-preserving functions’.
- **Par**: the category having objects sets, and morphisms $A \rightarrow B$ the functions $f : A \rightarrow B$, possibly defined only on a subset D of A (the ‘domain’).² How does this request affect the composition operation (the usual function composition) and the choice of identity morphisms? Make precise the statement that the categories **Par** and **Set*** ‘look very much alike’.
- **Σ -Seq**: the category of (sequential) Σ -*acceptors*, where Σ is a finite set of input symbols, $\Sigma = \{\sigma_1, \dots, \sigma_n\}$. An object of Σ -Seq is a quadruple (Q, δ, q_0, F) , where Q is a finite set of states, $\delta : \Sigma \times Q \rightarrow Q$ is a transition map, $q_0 \in Q$ is the initial state, and $F \subseteq Q$ is the set of final states.

A morphism $f : (Q, \delta, q_0, F) \rightarrow (Q', \delta', q'_0, F')$ (called a simulation) between Σ -acceptors is a function $f : Q \rightarrow Q'$ that preserves

 - transitions, i.e., $\delta'(\sigma, f(q)) = f(\delta(\sigma, q))$,
 - the initial state, i.e., $f(q_0) = q'_0$, and
 - the final states, i.e., $f[F] \subseteq F'$.
- **Aut**: the category of *automata*, with objects all (deterministic, sequential, Moore) automata. Objects are sextuples $(Q, \Sigma, Y, \delta, q_0, y)$, where Q is the set of states, Σ and Y are the sets of input symbols and output symbols, respectively, $\delta : \Sigma \times Q \rightarrow Q$ is the transition map, $q_0 \in Q$ is the initial state, and $y : Q \rightarrow Y$ is the output map. Morphisms from an automaton $(Q, \Sigma, Y, \delta, q_0, y)$ to an automaton $(Q', \Sigma', Y', \delta', q'_0, y')$ are triples (f_Q, f_Σ, f_Y) of functions $f_Q : Q \rightarrow Q'$, $f_\Sigma : \Sigma \rightarrow \Sigma'$, and $f_Y : Y \rightarrow Y'$ satisfying the following conditions:
 - preservation of transition: $\delta'(f_\Sigma(\sigma), f_Q(q)) = f_Q(\delta(\sigma, q))$,
 - preservation of outputs: $f_Y(y(q)) = y'(f_Q(q))$,
 - preservation of initial state: $f_Q(q_0) = q'_0$.

¹If C is a class of sets, we say that a family of functions $\{f_{XYZ} : X \times Y \rightarrow Z\}$ indexed by the elements $X, Y, Z \in C$ is *associative* if

$$f_{WZU}(w, f_{XYZ}(x, y)) = f_{ZYU}(f_{WXZ}(w, x), y)$$

for every tuple X, Y, Z, U, W and elements for which this is meaningful. In the present case this evidently translates into the familiar associativity of composition $u \circ (v \circ w) = (u \circ v) \circ w$.

²For example, the function $f : \mathbf{R} \rightarrow \mathbf{R} : x \mapsto \frac{x+1}{x-1}$ is a morphism in **Par**, but not in **Set**.

- turn these ideas into precise statements: (1) every poset (P, \leq) gives rise to a category $c[P]$ with objects the elements of P , and where there is a unique morphism $x \rightarrow y$ if and only if $x \leq y$ in P ; (2) every monoid M gives rise to a category \mathbf{BM} having a single object \bullet , and where $\mathbf{BM}(\bullet, \bullet) = M$.
- the category \mathbf{Vec}_K where the set of objects is the set of natural numbers $\{0, 1, 2, \dots\}$ and the set of morphisms $n \rightarrow m$ is the set of $m \times n$ matrices with entries in the field K .
- Let \mathbf{Fld} be the category of fields; starting from it we can define a category \mathbf{Vec} (note the absence of subscript) containing *literally all* vector spaces in the following way: an object of \mathbf{Vec} is a pair (K, V) , where K is a field and V a vector space on K . A morphism $(K, V) \rightarrow (L, W)$ is a pair $u : K \rightarrow L$ and $f : V \rightarrow W$ such that u is a homomorphism of rings, and $f : V \rightarrow W$ a function such that $f(v+v') = fv + fv'$ and $f(av) = u(a)f(v)$ for every $a \in K, v \in V$. Define identities and compositions ‘in the obvious way’ and prove that the resulting structure is indeed a category.
- Define a category C with a single object \bullet , and where the set of morphisms $\bullet \rightarrow \bullet$ is specified in BNF as

$$t ::= x_0 \mid c \mid f t \mid g t \quad (1.5)$$

where x_0 is one given variable, c is a constant and f, g are two *different* given function symbols. Composition is defined as substitution $t[t'/x_0]$ (where t' replaces x_0 in t is defined recursively).

If you know what to do with it, you are allowed to use a proof-assistant that checks the axioms of category.³

- From the category of sets, remove all functions $A \rightarrow B$ when $A \neq B$ are different sets; is the result still a category C ?

EXERCISE CF.2 \square Let C be a category. Show that ‘being isomorphic’ is an equivalence relation on the set of objects of a category.

EXERCISE CF.3 \square Let C be a category. Show that ‘there exists a morphism between’ is a preorder relation on the set of objects of a category.

EXERCISE CF.4 \square Deduce from the previous exercise that there is a poset C^P associated to every category C ; the poset C^P is called the *posetal reflection* of C . Show that every functor $C \rightarrow P$, where (P, \leq) is a poset, defines a unique monotone function $C^P \rightarrow P$.

EXERCISE CF.5 \square Let A be a set; show that there exist

- the ‘minimal’ category A^δ on A , where the objects are the elements of A , and only identity morphisms exist;

³If you are very brave: a monoid is exactly a category with a single object; this means that C above is isomorphic to a certain monoid M , whose elements are the terms $t = x_0 \mid c \mid f t \mid g t$ and whose monoid operation is defined by substitution. Describe the monoid M .

- the ‘maximal’ category A^X on A , where the objects are the elements of A , and there exists *exactly one* morphism between any two objects.

Show that any two objects of A^X are isomorphic.⁴

EXERCISE CF.6 □ Detail the construction that gives the ‘minimal’ and ‘maximal’ category on a set A (i.e., ‘regard a set A as a discrete category’, and ‘regard a set A as a category with *exactly one* morphism between *any* two elements’).

EXERCISE CF.7 □ A procedure to build a category is to ‘force a bunch of monoids to live together’: take a family of monoids M_i indexed by a set I , and define a category $\biguplus M_i$ having objects the elements $i \in I$ and morphisms specified by

$$\text{hom}(i, j) = \begin{cases} M_i & i = j \\ \emptyset & i \neq j \end{cases} \quad (1.6)$$

Prove that this is indeed a category.

EXERCISE CF.8 □ Recall that a (*directed*) *graph* \mathcal{G} consists of a pair of sets G_0, G_1 equipped with functions

$$s, t : G_1 \rightarrow G_0 \quad (1.7)$$

sending each *edge* (element of G_1) to a pair of *vertices* (elements of G_0); a directed graph \mathcal{G} gives rise to a *free category*, obtained as follows:

- the set of objects of $F\mathcal{G}$ is the set of vertices G_0 of \mathcal{G} ;
- the set of morphisms $v \rightarrow w$ between two vertices $v, w \in G_0$ is the set of all tuples

$$(v, \vec{x}, w) = v \rightarrow x_1 \rightarrow x_2 \rightarrow \cdots \rightarrow x_n \rightarrow w \quad (1.8)$$

with the convention that if $v = w$ and $n = 0$ the tuple is empty and equal to an element $()_v \in F\mathcal{G}(v, v)$.

The composition operation in $F\mathcal{G}$ is defined as

$$(u, \vec{y}, w) \circ (v, \vec{x}, u) = (v, \vec{x} \uplus_u \vec{y}, w) \quad (1.9)$$

where $\vec{x} \uplus_u \vec{y} = x_1 \rightarrow \cdots \rightarrow x_n \rightarrow u \rightarrow y_1 \rightarrow \cdots \rightarrow y_m$ if $\vec{x} = (x_1 \rightarrow \cdots \rightarrow x_n)$ and $\vec{y} = (y_1 \rightarrow \cdots \rightarrow y_m)$. Show that this is in fact a category (for each $v \in G_0$, the element $()_v$ is the identity arrow in $F\mathcal{G}$ for the composition defined above; composition is associative; etc).

EXERCISE CF.9 □ Define a **verbose category** to be a tuple $\mathcal{A} = (\mathcal{O}, \mathcal{M}, \text{dom}, \text{cod}, \circ)$ consisting of

- a class \mathcal{O} , called the class of \mathcal{A} -objects,
- a class \mathcal{M} , called the class of \mathcal{A} -morphisms,

⁴The notation A^X stands for the *chaotic* (Gr. $\chi\acute{\alpha}\omicron\varsigma$) category on A .

- functions $\text{dom} : \mathcal{M} \rightarrow \mathcal{O}$ and $\text{cod} : \mathcal{M} \rightarrow \mathcal{O}$, assigning to each morphism its domain and codomain, and
- a function \circ from $D = \{(f, g) \mid f, g \in \mathcal{M} \text{ and } \text{dom}(f) = \text{cod}(g)\}$ to \mathcal{M} with $\circ(f, g)$ written $f \circ g$,

subject to the following conditions:

- If $(f, g) \in D$, then $\text{dom}(f \circ g) = \text{dom}(g)$ and $\text{cod}(f \circ g) = \text{cod}(f)$.
- If (f, g) and (h, f) belong to D , then $h \circ (f \circ g) = (h \circ f) \circ g$.
- For each $A \in \mathcal{O}$ there exists a morphism e such that $\text{dom}(e) = A = \text{cod}(e)$ and
 - $f \circ e = f$ whenever $(f, e) \in D$, and
 - $e \circ g = g$ whenever $(e, g) \in D$.
- For any $(A, B) \in \mathcal{O} \times \mathcal{O}$, the class $\{f \in \mathcal{M} \mid \text{dom}(f) = A, \text{cod}(f) = B\}$ is a set.

Compare the definition of verbose category with that of category and determine in which sense these definitions can be considered ‘equivalent’.

EXERCISE CF.10 \square Consider the category C/A of arrows with a common codomain A , and where morphisms are given by commutative triangles

$$\begin{array}{ccc}
 X & \xrightarrow{f} & Y \\
 & \searrow h & \swarrow k \\
 & & A
 \end{array}
 \tag{1.10}$$

(more formally, C/A has as objects all the arrows $h : X \rightarrow A$, and $C/A(h, k)$ for $h : X \rightarrow A$ and $k : Y \rightarrow A$ consists of the subset of $C(X, Y)$ made of those $f : X \rightarrow Y$ such that $kf = h$.)

Does C/A has an initial object? Does it have a terminal object? Try to write down the definition of product of two objects h, k in C/A ; try to write down the definition of coproduct of two object h, k in C/A . A product of h, k in C/A is called the *pullback* or *fibered product* of h, k in C .

EXERCISE CF.11 \square Describe as precisely as possible the category C/A for every example of category in [Exercise 1](#).

EXERCISE CF.12 \square Let C be a category; define the *arrow category* C^{\rightarrow} having

- objects all the morphisms $u : X \rightarrow Y$ of C ;
- morphisms $\left[\begin{array}{c} X \\ u \downarrow \\ Y \end{array} \right] \rightarrow \left[\begin{array}{c} A \\ v \downarrow \\ B \end{array} \right]$ the commutative squares

$$\begin{array}{ccc}
 X & \xrightarrow{f} & A \\
 u \downarrow & & \downarrow v \\
 Y & \xrightarrow{g} & B
 \end{array}
 \tag{1.11}$$

Define identities and composition in the obvious way.

Does C^\rightarrow have an initial object? Does it have a terminal object? Write down the definition of product and of coproduct in C^\rightarrow .

In what sense, if any, $(C^{\text{op}})^\rightarrow$ is equivalent to $(C^\rightarrow)^{\text{op}}$?

EXERCISE CF.13 \square Verify whether the following are examples of functors: for those where only the correspondence on objects is given, try to define the one on morphisms or show that there is none making them a functor.

- the empty functor $F_\emptyset : \emptyset \rightarrow C$, where \emptyset is the empty category, and C is any other category; show that there is only one such functor.
- the identity functor $1 : C \rightarrow C$ of a given category C , acting as the identity function both on objects and on morphisms.
- the constant functor $c_X : \mathcal{A} \rightarrow C$, sending all objects $A \in \mathcal{A}_o$ to $X \in C_o$, and all morphisms $A \rightarrow A'$ to the identity morphism of X .
- the functor $d_X : \bullet \rightarrow C$ from the singleton category, ‘choosing the object X ’ and its identity. (Explain what this means formally.)
- Given a category C and a morphism $f \in C(C_0, C_1)$, the functor $m_f : \{0 \rightarrow 1\} \rightarrow C$ from the category with two objects and a single nonidentity arrow, sending 0 to C_0 , 1 to C_1 , and $0 \rightarrow 1$ to $f : C_0 \rightarrow C_1$.
- Given any category C and object $A \in C_o$, the functor $C(A, -) : C \rightarrow \text{Set}$ sending $X \in C_o$ to the set of morphisms $u : A \rightarrow X$, and each morphism $g : X \rightarrow Y$ to the function $C(A, X) \rightarrow C(A, Y)$ sending $u : A \rightarrow X$ to $g \circ u : A \rightarrow Y$.
- Given any category C and object $A \in C_o$, the functor $C(-, A) : C^{\text{op}} \rightarrow \text{Set}$ sending $X \in C_o$ to the set of morphisms $u : X \rightarrow A$, and each morphism $g : X \rightarrow Y$ to the function $C(Y, A) \rightarrow C(X, A)$ sending $u : Y \rightarrow A$ to $u \circ g : X \rightarrow A$.
- Define a functor $\text{Set} \rightarrow \text{Mon}$ sending a set A to the set A^* of all finite lists with entries in A ; how does a function $f : A \rightarrow B$ induces a monoid homomorphism $A^* \rightarrow B^*$?
- Let G be a graph; send G to the set of its *connected components*, i.e. the set G_0 of its vertices modulo the equivalence relation generated by the source and target function $G_1 \rightarrow G_0 \times G_0$: $a, b \in G_0$ are equivalent if there is an edge $a \rightarrow b$ in G . Show that this is a functor $\text{Graph} \rightarrow \text{Set}$.
- Try to define a functor $\text{Grp} \rightarrow \text{Ab}$ sending a group to its abelianization.
- Try to define a functor $\text{Grp} \rightarrow \text{Ab}$ sending a group to its center.⁵
- Try to define a functor sending a ring to its group of invertible elements.
- There is a functor $\text{Ring} \rightarrow \text{Grp}$ sending a ring R to the group of $n \times n$ invertible matrices with entries in R . Define its correspondence on morphism and show that it is a functor.

⁵The *center* of a group G is the set $\{g \in G \mid \forall x \in G, gx = xg\}$.

- Send a ring R to the set of all its prime ideals. Is this a functor $\text{Ring} \rightarrow \text{Set}$? Covariant or contravariant?
- Send a topological space to the set of its clopen (=both open and closed) subsets; is this a functor? Covariant or contravariant?
- Send a group G to its *group algebra* $\mathbf{Z}[G]$, the set of formal sums of integers indexed by elements of G , $\sum_{g \in G} n_g$; define a ring operation on $\mathbf{Z}[G]$ and show that this is the object part of a functor $\text{Grp} \rightarrow \text{Ring}$. Recall that a few items ago you defined a functor sending a ring R to its group of invertible elements R^\times ; construct a bijection between the set of group homomorphisms $G \rightarrow R^\times$ and the set of ring homomorphisms $\mathbf{Z}[G] \rightarrow R$.
- Send a set X to the set of polynomials with integer coefficients $\mathbf{Z}[x \mid x \in X]$; does this define the object part of a functor $\text{Set} \rightarrow \text{Ring}$?
- Send a ring R to the ring $R[t]$ of polynomials in a single indeterminate t ; does this define the object part of a functor $\text{Ring} \rightarrow \text{Ring}$?
- Send a group to the poset of all its subgroups; is this the object part of a functor $\text{Grp} \rightarrow \text{Pos}$?
- Let M be a monoid regarded as a category with a single object. What is a functor $M \rightarrow \text{Set}$? What is a functor $M^{\text{op}} \rightarrow \text{Set}$?
- Let (P, \leq) be a poset regarded as a category. What is a functor $(P, \leq)^{\text{op}} \rightarrow \text{Set}$?
- In the notation above, let P be the poset of open subsets of the real line \mathbf{R} (with respect to the usual ‘Euclidean’ topology). Consider the correspondence

$$U \longmapsto C^0(U) \tag{1.12}$$

sending an open subset $U \subseteq \mathbf{R}$ to the set of all continuous functions $f : U \rightarrow \mathbf{R}$. Show that this defines a functor $P^{\text{op}} \rightarrow \text{Set}$, and in particular that every inclusion $U \subseteq V$ of open sets gives rise to a *restriction* operation $C^0V \rightarrow C^0U$ sending a function $f : V \rightarrow \mathbf{R}$ to its ‘restriction’ $f|_U : U \rightarrow \mathbf{R}$. Show that the following two properties are satisfied, given any $U \in P$ and any covering $\{V_i \mid i \in I\}$ of U :⁶

- if $f, g \in C^0(U)$ are such that $f|_i = g|_i$ for all $i \in I$, then $f = g$ ($f|_i = f|_{V_i}$ for short).
- if $f_i \in C^0(V_i)$ is a family of functions such that $f_i|_{V_i \cap V_j} = f_j|_{V_i \cap V_j}$ for every $i, j \in I$, then there exists a function $f \in C^0(U)$ such that $f|_i = f_i$ for every $i \in I$.

EXERCISE CF.14 \square Define a correspondence $G : C \rightarrow \text{Mon}$ (the category of monoids), sending a set A to the monoid A^\star of finite lists of elements of A , i.e. to the set of all finite lists (a_1, \dots, a_n) where $n \geq 0$ and $a_i \in A$ for each $i = 1, \dots, n$, and a function $f : A \rightarrow A$ to the function

$$A^\star \rightarrow A^\star : (a_1, \dots, a_n) \mapsto (f(a_1), \dots, f(a_n)) \tag{1.13}$$

⁶A *covering* of $U \in P$ is a family $\{V_i \mid i \in I\}$ of elements of P such that $\bigcup V_i = U$.

Is G a functor $C \rightarrow \text{Mon}$?

EXERCISE CF.15 \square Given two functors $F : C \rightarrow X$ and $G : D \rightarrow X$ define the *comma category* (F/G) having

- objects the triples (C, D, h) where $h : FC \rightarrow GD$ is a morphism in X ;
- morphisms $(C, D, h) \rightarrow (C', D', k)$ the pairs $u : C \rightarrow C', k : D \rightarrow D'$ such that the square

$$\begin{array}{ccc} FC & \xrightarrow{Fu} & FC' \\ h \downarrow & & \downarrow k \\ GD & \xrightarrow{Gv} & GD' \end{array} \quad (1.14)$$

is commutative.

Verify that it is a category; does (F/G) have an initial object? Does it have a terminal object? Fix an object X of the codomain of a given functor $F : C \rightarrow \mathcal{D}$; then, define the category (F/X) as the comma between F and the functor $d_X : \bullet \rightarrow \mathcal{D}$ ‘selecting’ X .

EXERCISE CF.16 \square Describe as precisely as possible the comma category (F/X) for each functor in [Exercise 13](#).

EXERCISE CF.17 \square A *simple polynomial* is a functor $F : \text{Set} \rightarrow \text{Set}$ that is defined from the following inductive rules:

- SP1) the identity functor $X \mapsto X$ is a simple polynomial;
- SP2) every constant functor $X \mapsto A$ is a simple polynomial;
- SP3) the product $F \times G : X \mapsto FX \times GX$ of two simple polynomials is simple;
- SP4) the coproduct $\coprod_{i \in I} F_i : X \mapsto \coprod_{i \in I} F_i X$ of an arbitrary number of simple polynomials is simple.

An example of a polynomial functor is $X \mapsto A \times X^3 + B \times X^2 + X + 1$, where \times denotes cartesian product, and $+$ denotes coproduct; another example is a ‘formal series functor’ $X \mapsto \coprod_{i \in I} A_i \times X^{n_i}$ where n_i are natural numbers and $(A_i \mid i \in I)$ is an arbitrary family of sets.

An *arity function* consists of a set I equipped with a function $a : I \rightarrow \mathbb{N}$; the inverse image $a^{-1}n$ is the set of elements in I having ‘arity’ n .⁷ Every arity function $a : I \rightarrow \mathbb{N}$ defines an *arity functor* as

$$F_a : X \mapsto \prod_{i \in I} X^{a(i)} = \{(i, \underline{x}) \mid i \in I, \underline{x} \in X^{a(i)}\} \quad (1.15)$$

Show that the class of simple polynomials coincides with the class of arity functors (F_a is ‘clearly’ a simple polynomial: how does one define an arity associated to a given simple polynomial?)

⁷The word ‘arity’ is a back-formation from the Latin adjectival numeral suffix *-arius*, used to form adjectives from nouns or numerals.

EXERCISE CF.18 \square Verify that the following are examples of natural transformations:

-
-
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EXERCISE CF.19 \square Natural transformations are very common, but destroying the naturality of a transformation is very easy: find a natural transformation between two functors $F, G : C \rightarrow D$, change the value of one of its components, and prove that the result is not natural any more.

EXERCISE CF.20 \square Let $P : \text{Set} \rightarrow \text{Set}$ be the correspondence that sends a set A to the power set PA of A , the set of all subsets $U \subseteq A$, and a function $f : A \rightarrow B$ to the function $Pf : PA \rightarrow PB$, that sends a subset $U \subseteq A$ to the *image*

$$f_*U := \{fu \mid u \in U\} \quad (1.16)$$

Similarly, let $d : \text{Set} \rightarrow \text{Set}$ be the correspondence that sends A to PA , but a function $f : A \rightarrow B$ to the function $PB \rightarrow PA$, that sends a subset $V \subseteq B$ to the *inverse image*

$$f^*V := \{a \in A \mid fa \in V\} \quad (1.17)$$

- Show that both P, d are functors (d is contravariant, i.e. $d(f \circ g) = dg \circ df$); show that given subsets $U \in PA, V \in PB$ one has

$$f_*U \subseteq V \iff U \subseteq f^*V. \quad (1.18)$$

- What is a natural transformation $f^*f_* \Rightarrow 1_{PA}$, regarding PA as a category? Show that for each $U \in PA$, one has $U \subseteq f^*f_*U$: is this a natural transformation $f^*f_* \Rightarrow 1_{PA}$?

2. Co/limits

EXERCISE CL.1 \square Let C be a category and $F : I \rightarrow C$ be a constant functor, say $FI = C$ for every $I \in I$ and $Ff = 1_C$ for every $f : I \rightarrow I'$. Is it true that, when it exists, the limit of F is also C ? If not, find a counterexample (easy) and a general formula to express $\lim F$ (harder). Dualise to the case of colimits.

EXERCISE CL.2 \square Let the category Dyn be defined by having

- objects the triples (X, f, x_0) where X is a set, $f : X \rightarrow X$ an endofunction, and $x_0 \in X$ an element;

- a morphism $(X, f, x_0) \rightarrow (Y, g, y_0)$ is a function $u : X \rightarrow Y$ with the property that $u(x_0) = y_0$ and that the square

$$\begin{array}{ccc}
 X & \xrightarrow{f} & X \\
 u \downarrow & & \downarrow u \\
 Y & \xrightarrow{g} & Y
 \end{array} \tag{2.1}$$

is commutative.

Explain in what sense the initial object of this category is the set \mathbf{N} of natural numbers. Prove that \mathbf{N} is a monoid *using the universal property only*.

EXERCISE CL.3 \square Define the following categories, show the category axioms, and unwind the definition of what is a terminal object in each of them.

- Let S be a set and $\{A_s \mid s \in S\}$ a collection of sets indexed by S . Define the category $\Pi(A_s \mid s \in S)$ as follows: an object consists of a pair $(Z, \mathbf{f} = \{f_s \mid s \in S\})$ where Z is a set and $\mathbf{f} = \{f_s : Z \rightarrow A_s\}$ is a family of functions indexed by S ; a morphism $(Z, \mathbf{f}) \rightarrow (W, \mathbf{g})$ consists of a function $u : Z \rightarrow W$ such that $g_s \circ u = f_s$ for every $s \in S$:

$$\begin{array}{ccc}
 & A_s & \\
 f_s \nearrow & & \nwarrow g_s \\
 Z & \xrightarrow{u} & W
 \end{array} \tag{2.2}$$

Investigate in particular the edge cases: what if $S = \emptyset$? What if S is a singleton? What if S has two elements $\{a, b\}$?

- Let S, X, Y be sets and $\{f_s : X \rightarrow Y\}$ a collection of functions with the same domain and codomain, indexed by S . Define the category $\Gamma(f_s \mid s \in S)$ as follows: an object consists of a pair $(Z, u : Z \rightarrow X)$ with the property that $f_s \circ u = f_t \circ u$ for every $s, t \in S$; a morphism $(Z, u) \rightarrow (W, v)$ consists of a function $h : Z \rightarrow W$ with the property that $v \circ h = u$:

$$\begin{array}{ccc}
 Z & \xrightarrow{u} & X \xrightarrow{f_s} Y \\
 & \searrow h & \vdots \\
 & & W \xrightarrow{v} X \xrightarrow{f_t} Y
 \end{array} \tag{2.3}$$

Investigate in particular the edge cases: what if $S = \emptyset$? What if S is a singleton? What if S has two elements $\{a, b\}$?

- Let S be a set, and $\{f_s : X_s \rightarrow Y\}$ a family of functions with the same codomain Y , indexed by S . Define the category $\Lambda(f_s \mid s \in S)$ as follows: an object consists of a pair $(Z, \{u_s : Z \rightarrow X_s\})$ where Z is a set and $\mathbf{u} = \{u_s : Z \rightarrow X_s\}$

is a family of functions indexed by S , with the property that the composition $f_s \circ u_s : Z \rightarrow X_s \rightarrow Y$ is independent from the index $s \in S$; a morphism $(Z, \mathbf{u}) \rightarrow (W, \mathbf{v})$ consists of a function $t : Z \rightarrow W$ such that $v_s \circ t = u_s$:

$$\begin{array}{ccc}
 W & \xrightarrow{v_s} & X_s \\
 \downarrow v_{s'} & \searrow & \downarrow f_s \\
 & & \vdots \\
 & & \downarrow f_s \\
 X_{s'} & \xrightarrow{f_{s'}} & Y
 \end{array} \tag{2.4}$$

EXERCISE CL.4 \square Fix a set A . Consider the functor S_A sending a set X to the set $1+(A \times X)$, whose elements are of two kinds: either the single element $\perp \in 1$, or an element $(a, x) \in A \times X$.

An S_A -**algebra** consists of a pair (X, σ) where X is a set and $\sigma : S_A X \rightarrow X$ is a function. A morphism of S_A -algebras $(X, \sigma) \rightarrow (Y, \tau)$ consists of a function $u : X \rightarrow Y$ such that the square

$$\begin{array}{ccc}
 1 + (A \times X) & \xrightarrow{\sigma} & X \\
 1+A \times u \downarrow & & \downarrow u \\
 1 + (A \times Y) & \xrightarrow{\tau} & Y
 \end{array} \tag{2.5}$$

is commutative. Show that this defines a category $\text{Alg}(S_A)$.

Describe the initial object A^* of $\text{Alg}(S_A)$.

EXERCISE CL.5 \square Consider again the functor S_A defined above; an S_A -**coalgebra** consists of a pair (U, r) where U is a set and $r : U \rightarrow S_A U$ is a function. A morphism of S_A -coalgebras $(U, r) \rightarrow (V, t)$ consists of a function $h : U \rightarrow V$ such that the square

$$\begin{array}{ccc}
 U & \xrightarrow{r} & S_A U \\
 h \downarrow & & \downarrow S_A h \\
 V & \xrightarrow{t} & S_A V
 \end{array} \tag{2.6}$$

is commutative.

- Show that this defines a category $\text{coAlg}(S_A)$;
- describe the terminal object \hat{A} of $\text{coAlg}(S_A)$;
- is there a relation between A^* and \hat{A} ? (For example, can one be identified with a subset of the other?)

EXERCISE CL.6 \square Let Ab be the category of abelian groups; find explicit descriptions for

- the product of two objects A, B ; from this, derive an explicit description for $A_1 \times \cdots \times A_n$ for every $n \geq 2$;

- the coproduct of two objects A, B ; from this, derive an explicit description for $A_1 + \dots + A_n$ for every $n \geq 2$. In particular, prove that there is a natural isomorphism

$$A \times B \cong A + B \quad (2.7)$$

i.e. an isomorphism of functors $_ \times _ \cong _ + _$; the construction $A \times B \cong A + B$, in this context, is denoted $A \oplus B$ and called the **biproduct** of A, B ;

- the equalizer $E(f, g)$ of a pair of homomorphisms $f, g : A \rightarrow B$; in particular, find an explicit description when $g = 0$ is the zero map; the equaliser of $(f, 0)$ is called the **kernel** of f ;
- the pullback $A \times_C B$ of a pair of maps $A \xrightarrow{f} C \xleftarrow{g} B$; in particular, find an explicit description for the equalizer of the pair of maps

$$\mathbf{Z} \xrightarrow{-\cdot m} \mathbf{Z} \xleftarrow{-\cdot n} \mathbf{Z} \quad (2.8)$$

EXERCISE CL.7 \square Prove that the pullback of $A \xrightarrow{f} C \xleftarrow{g} B$ is canonically isomorphic to the equaliser of the pair of maps

$$\begin{array}{ccc} & & C \\ & & \downarrow \Delta \\ A \oplus B & \xrightarrow{f \oplus g} & C \oplus C \end{array} \quad (2.9)$$

where Δ is the diagonal map $x \mapsto (x, x)$.

Prove that the equaliser of $f, g : A \rightarrow B$ is canonically isomorphic to the kernel of the map $\begin{bmatrix} f \\ g \end{bmatrix} : A \oplus A \rightarrow B : a \mapsto f(a) - g(a)$.

A **semiadditive category** is a category \mathcal{C} with a zero object, finite products and finite coproducts, such that the canonical map

$$A + B \longrightarrow A \times B \quad (2.10)$$

is the component at (A, B) of a natural isomorphism of functors $_ \times _ \cong _ + _$.

The category of abelian groups is semiadditive; more generally, the category of R -modules over a ring R is semiadditive (cf. [Exercise 51](#), [52](#), [53](#)).

EXERCISE CL.8 \square Show that in a semiadditive category \mathcal{C} , every set $\mathcal{C}(X, Y)$ becomes a commutative monoid under the operation

$$f + g : X \xrightarrow{\Delta} X \oplus X \xrightarrow{f \oplus g} Y \oplus Y \xrightarrow{\nabla} Y. \quad (2.11)$$

Who is the identity element $0 : X \rightarrow Y$?

EXERCISE CL.9 \square **Important Pullbacks and Pushouts in various categories**

EXERCISE CL.10 \square Recall the definition of the category Dyn of (unpointed) **dynamical systems**:

- Objects are pairs (X, s) where $s : X \rightarrow X$ is a function on the set X ;
- Morphisms $(X, s) \rightarrow (Y, t)$ are functions $f : X \rightarrow Y$ such that the diagram

$$\begin{array}{ccc}
 X & \xrightarrow{s} & X \\
 f \downarrow & & \downarrow f \\
 Y & \xrightarrow{t} & Y
 \end{array} \tag{2.12}$$

is commutative.

A dynamical system (X, s) is called **reversible** if $s : X \rightarrow X$ is an invertible function. A morphism between two reversible dynamical systems is just a morphism of dynamical systems.

Show that the inclusion functor $\text{RevDyn} \hookrightarrow \text{Dyn}$ admits a left adjoint, i.e. that for every morphism

$$f : (X, s) \longrightarrow (A, \sigma) \tag{2.13}$$

where (A, σ) is a reversible dynamical system, there exist

- A reversible dynamical system (\bar{X}, \bar{s}) with a map $u : (X, s) \rightarrow (\bar{X}, \bar{s})$ of dynamical systems;
- a *unique* $\bar{f} : \bar{X} \rightarrow A$ which is a morphism of dynamical systems, with the property that $\bar{f} \circ u = f$:

$$\begin{array}{ccc}
 X & \xrightarrow{f} & A \\
 u \downarrow & \nearrow \bar{f} & \\
 \bar{X} & &
 \end{array} \tag{2.14}$$

thus realising the isomorphism

$$\text{Dyn}(X, A) \cong \text{RevDyn}(\bar{X}, A). \tag{2.15}$$

EXERCISE CL.11 \square Define the following category:

$$\begin{array}{ccc}
 & & 1 \\
 & & \parallel \\
 & & \downarrow \\
 0 & \rightrightarrows & 2
 \end{array} \tag{2.16}$$

where the parallel arrows are called $\iota_0, \iota_1 : 1 \rightrightarrows 2$ and $j_0, j_1 : 0 \rightrightarrows 2$.

The **joint coequaliser** for two pairs of functions $i_0, i_1 : X_1 \rightrightarrows X_2$ and $j_0, j_1 : X_0 \rightrightarrows X_2$ consists of a colimit for a diagram $X : \mathcal{J} \rightarrow \text{Set}$ of shape \mathcal{J} ; this means that there is a

diagram

$$\begin{array}{ccc}
 & & X_1 \\
 & & \downarrow i_1 \\
 & & \downarrow i_0 \\
 X_0 & \xrightarrow{j_0} & X_2 \\
 & \xrightarrow{j_1} & \\
 & &
 \end{array}
 \quad (2.17)$$

where $X(t_0) = i_0$, $X(t_1) = i_1$, etc., and a morphism $t : X_2 \rightarrow C$ such that $ti_0 = ti_1$ and $tj_0 = tj_1$, and such that t is initial with respect to this property, i.e. for every other $x : X_2 \rightarrow Z$ such that $xi_0 = xi_1$ and $xj_0 = xj_1$ one has $x = \bar{x}t$ for a unique $\bar{x} : C \rightarrow X$.

Show that the joint coequaliser of (i_0, i_1) , (j_0, j_1) can be obtained as follows: start from the diagram (2.17) above, and consider the diagram

$$\begin{array}{ccccc}
 & & X_1 & & \\
 & & \downarrow i_1 & & \\
 & & \downarrow i_0 & & \\
 X_0 & \xrightarrow{j_0} & X_2 & \xrightarrow{h} & U \\
 & \xrightarrow{j_1} & \downarrow k & & \downarrow \\
 & & V & \xrightarrow{p} & P
 \end{array}
 \quad (2.18)$$

where $h : X_2 \rightarrow U$ is the coequaliser of (j_0, j_1) , $k : X_2 \rightarrow V$ is the coequaliser of (i_0, i_1) and P is the pushout of (k, h) .

EXERCISE CL.12 \square Tensor products

EXERCISE CL.13 \square An object C of a category \mathcal{C} is called **tiny** if the functor $\mathcal{C}(C, -)$ preserves all colimits.

- Prove that the singleton $*$ of **Set** is a tiny object; is the coproduct of two tiny objects still tiny? Is the two-element set $* \amalg *$ tiny in **Set**? Prove that \mathbf{Z} is a tiny object in the category of abelian groups. Is the group \mathbf{Z}^n still tiny?
- Prove that a functor $P \in [\mathcal{C}^{\text{op}}, \mathbf{Set}]$ is a tiny object if and only if it is a retract of a representable functor.

EXERCISE CL.14 \square Direct and inverse images

EXERCISE CL.15 \square Transfinite constructions

3. Adjoints

EXERCISE AD.1 \square Let \mathcal{A} be a small category. Show that there is a pair of adjoint functors

$$O : [\mathcal{A}^{\text{op}}, \mathbf{Set}] \rightleftarrows [\mathcal{A}, \mathbf{Set}]^{\text{op}} : S \quad (3.1)$$

where O sends a functor $P : \mathcal{A}^{\text{op}} \rightarrow \mathbf{Set}$ to the functor sending $A \in \mathcal{A}$ to $[\mathcal{A}^{\text{op}}, \mathbf{Set}](P, \mathcal{A}(-, A))$, and S sends a functor $Q : \mathcal{A} \rightarrow \mathbf{Set}$ to the functor sending $A \in \mathcal{A}$ to $[\mathcal{A}, \mathbf{Set}](Q, \mathcal{A}(A, -))$.

EXERCISE AD.2 \square A **directed graph** is a covariant functor on the category $\Gamma = \{E \begin{smallmatrix} \xrightarrow{s} \\ \xrightarrow{t} \end{smallmatrix} V\}$. Show that the forgetful functor $U : \text{dGph} \rightarrow \text{Set}$, i.e. the functor sending a directed graph to its set of vertices, has both a left and a right adjoint.

Show that the left adjoint is given by the functor that sends a set X to the directed graph having X as set of vertices, and no edges. How is the right adjoint defined?

EXERCISE AD.3 \square The category of adjunctions; mating

EXERCISE AD.4 \square Prove that two functors $F : C \rightleftarrows D : G$ are adjoints, with F left adjoint to G , if and only if the two comma categories $(F/1_D)$ and $(1_C/G)$ are ‘equivalent over $C \times D$ ’, namely there is an equivalence of categories $U : (F/1_D) \cong (1_C/G) : V$ with the property that the diagram

$$\begin{array}{ccc} (F/1) & \xrightarrow{\quad} & (1/G) \\ & \searrow X & \swarrow Y \\ & C \times D & \end{array} \quad (3.2)$$

is commutative (choosing either U or its inverse V as horizontal arrow).

Here, $X : (F/1) \rightarrow C \times D$ is the functor that sends an object $(C, D, FC \rightarrow D)$ to the pair (C, D) , and similarly Y sends an object $(C, D, C \rightarrow GD)$ to the pair (C, D) .

EXERCISE AD.5 \square **Strings of adjoints of arbitrary length**

EXERCISE AD.6 \square **Geometric morphisms in simple cases**

EXERCISE AD.7 \square Let (\mathbf{Z}, \leq) be the totally ordered set of integers, regarded as a category, and $f : \mathbf{Z} \rightarrow \mathbf{Z}$ a monotone function, regarded as an endofunctor. Show that the following conditions are equivalent:

- c1) f has a left adjoint f_L ;
- c2) f has a right adjoint f_R ;
- c3) the image $f(\mathbf{Z})$ of f is unbounded from below and from above.

(hint: show that f has a right adjoint if and only if the following condition holds:

- D1) each set $S_m = \{n \mid fn \leq m\}$ is nonempty and bounded from above; thus $f_R(m) := \max S_m$.

Show that this, in turn, is equivalent to the third condition above. Dualise for left adjoints.)

Let $f : \mathbf{Z} \rightarrow \mathbf{Z}$ be the map sending an integer k to $2k$, so that the image of f is $2\mathbf{Z}$; what are the left (and the right) adjoints f_L, f_R of f ?

Describe the monads obtained from the adjunction $f_L \dashv f$ and from the adjunction $f \dashv f_R$.

EXERCISE AD.8 \square Let $i : H \leq G$ be the inclusion of a subgroup regarded as a functor between one-object categories; if k is a field, the category of k -linear representations of

G can be identified with the functor category $[G, \text{Vect}_k]$, and similarly for H . (Make this statement precise.)

The scope of this exercise is to show that the functor

$$[H, \text{Vect}] \rightarrow [G, \text{Vect}] \quad (3.3)$$

induced by precomposing with i (=restricting the action of G to the subgroup H) has both a left and a right adjoint, i.e. that every k -linear representation of H can be extended in a maximal and minimal way to a representation of the whole G .

- Show that there is an equaliser diagram

$$\text{hom}_{k[H]}(k[G], V) \longrightarrow \text{hom}_k(k[G], V) \rightrightarrows \prod_{h \in H} \text{hom}(k[G], V) \quad (3.4)$$

where $k[G]$ is the group k -algebra of 63, and the two parallel maps are obtained as follows: $l_h : \text{hom}(k[G], V) \rightarrow \text{hom}(k[G], V)$ is obtained sending $f : k[G] \rightarrow V$ to $f(h \cdot _)$, and $r_h : \text{hom}(k[G], V) \rightarrow \text{hom}(k[G], V)$ sending f to $h \cdot f(_)$.

- Show that there is a coequaliser diagram

$$\prod_{h \in H} k[G] \otimes_k V \rightrightarrows k[G] \otimes_k V \longrightarrow k[G] \otimes_{k[H]} V \quad (3.5)$$

obtained quotienting by the relation prescribing $(h \cdot \alpha) \otimes v - \alpha \otimes (h \cdot v)$.

EXERCISE AD.9 \square Change of scalars; adjunctions between module categories

EXERCISE AD.10 \square A notable result in the theory of adjoint functors is the *adjoint functor theorem*, establishing sufficient condition for the existence of a left adjoint to a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ that preserves all limits: a functor F that preserves all limits has a left adjoint if and only if it satisfies a certain condition known as the *solution set condition*; a key result towards the proof of this equivalence is the following **initial object lemma**:

Let \mathcal{C} be a category admitting all small limits; then, \mathcal{C} has an initial object if and only if it has a **weakly initial family**, i.e. a set of objects $\{W_i \mid i \in I\}$ with the property that for every $X \in \mathcal{C}$ there exists at least (but possibly many) arrow $W_{i(X)} \rightarrow X$.

Prove the initial object lemma, following this guide:

- If \mathcal{C} has an initial object, it obviously has a weakly initial family;
- Conversely, build the product $W = \prod_{i \in I} W_i$ of all the elements of a weakly initial family.
- Consider the joint equaliser

$$K \xrightarrow{k} W \begin{array}{c} \rightrightarrows \\ \vdots \\ \rightrightarrows \end{array} W \quad (3.6)$$

of all endomorphisms of W (this means that k has the property that $ku = kv$ for every pair $u, v : W \rightarrow W$, and it is terminal with this property);

- K is a weakly initial object: why? Show that K is an initial object: assume $f, g : K \rightarrow X$ are parallel arrows out of K ; show that $f = g$ (hint: $f = g$ if and only if their equaliser is isomorphic to K).

EXERCISE AD.11 \square Closure of Pos with pointwise order; non-closure with lexicographic order.